

## TD 6

**Exercise 1** (Ultimately Periodic Words). An *ultimately periodic word* over  $\Sigma$  is a word of form  $u \cdot v^\omega$  with  $u$  in  $\Sigma^*$  and  $v$  in  $\Sigma^+$ .

Prove that any nonempty recognizable language in  $\text{Rec}(\Sigma^\omega)$  contains an ultimately periodic word.

**Exercise 2** (Muller Automata). A nondeterministic Muller automaton is a tuple  $\mathcal{A} = (Q, \Sigma, I, T, \mathcal{F})$ , where  $Q, \Sigma, I, T$  are as for Büchi automata and  $\mathcal{F} \subseteq 2^Q$  is the acceptance condition. For a run  $\sigma$  of  $\mathcal{A}$ , denote by  $\text{Inf}(\sigma)$  the set of states which are visited infinitely often. A run  $\sigma$  is accepting if  $\text{Inf}(\sigma) \in \mathcal{F}$ .

1. Give a *deterministic* Muller automaton for the language  $(a + b)^* a^\omega$ .
2. Show that for any Muller automaton  $\mathcal{A}$ ,  $L(\mathcal{A})$  is  $\omega$ -regular.
3. Show that any  $\omega$ -regular language is accepted by some (nondeterministic) Muller automaton.

Remark: in fact, any  $\omega$ -regular language can be recognized by some *deterministic* Muller automaton.

**Exercise 3** (Synchronous Büchi Transducers). Give unambiguous synchronous Büchi transducers for the following formulæ:

1.  $\text{SF } q$
2.  $\text{SG } q$
3.  $\text{G}(p \rightarrow \text{F } q)$