

An Experimental Ambiguity Detection Tool

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Standard ML

Milner et al. [1997]

Example

```
datatype 'a option = NONE | SOME of 'a
```

```
fun filter pred l =  
  let  
    fun filterP (x::r, l) =  
      case (pred x)  
        of SOME y => filterP (r, y::l)  
          | NONE   => filterP (r, l)  
    | filterP ([], l) = rev l  
  in  
    filterP (l, [])  
  end
```

The Issue

SML Compilers Behaviour

```
| filterP ([], l) = rev l
```

- ▶ MLton
- ▶ Moscow ML
- ▶ Poly/ML
- ▶ SML/NJ

```
Error: match.sml 9.25.
```

```
Syntax error: replacing EQUALOP with DARROW.
```

The Issue

SML Compilers Behaviour

```
| filterP ([], l) = rev l
```

- ▶ MLton
- ▶ Moscow ML
- ▶ Poly/ML
- ▶ SML/NJ

```
! Toplevel input:
```

```
!     | filterP ([], l) = rev l
```

```
!
```

```
! Syntax error.
```

The Issue

SML Compilers Behaviour

```
| filterP ([], l) = rev l
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Error: => expected but = was found

The Issue

SML Compilers Behaviour

```
| filterP ([], l) = rev l
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- ▶ Moscow ML
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```
stdIn:7.24-7.29 Error: syntax error:  
deleting EQUALOP ID
```

The Issue

Partial SML Grammar

$\langle dec \rangle \rightarrow \mathbf{fun} \langle fvalbind \rangle$
 $\langle fvalbind \rangle \rightarrow \langle sfvb \rangle \mid \langle fvalbind \rangle \mid \langle sfvb \rangle$
 $\langle sfvb \rangle \rightarrow vid \langle atpats \rangle = \langle exp \rangle$

$\langle exp \rangle \rightarrow \mathbf{case} \langle exp \rangle \mathbf{of} \langle match \rangle$
 $\langle match \rangle \rightarrow \langle mrule \rangle \mid \langle match \rangle \mid \langle mrule \rangle$
 $\langle mrule \rangle \rightarrow \langle pat \rangle \Rightarrow \langle exp \rangle$

$\langle atpats \rangle \rightarrow \langle atpat \rangle \mid \langle atpats \rangle \langle atpat \rangle$
 $\langle pat \rangle \rightarrow vid \langle atpat \rangle$
 $\langle atpat \rangle \rightarrow vid$

The Issue

Shift/Reduce Conflict

GNU Bison

```
state 20
```

```
6 exp: "case" exp "of" match .
```

```
8 match: match . '|' mrule
```

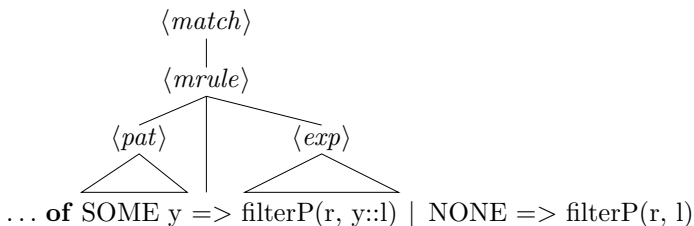
```
'|' shift, and go to state 24
```

```
'|' [reduce using rule 6 (exp)]
```


The Issue

Shift/Reduce Conflict

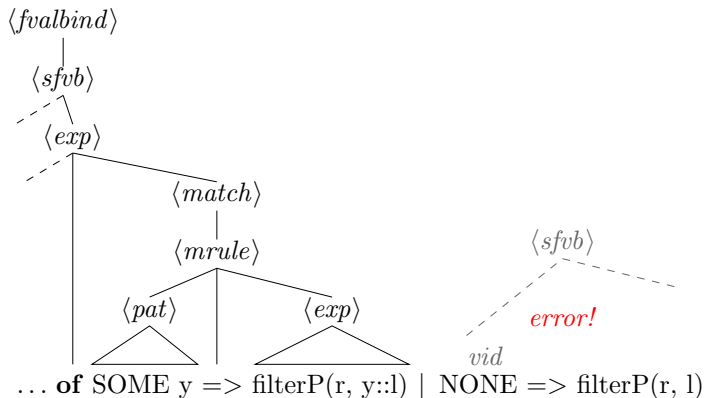
Which action to choose?



The Issue

Shift/Reduce Conflict

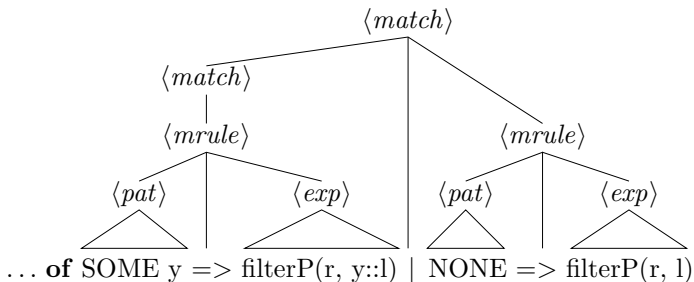
Which action to choose? Reduce?



The Issue

Shift/Reduce Conflict

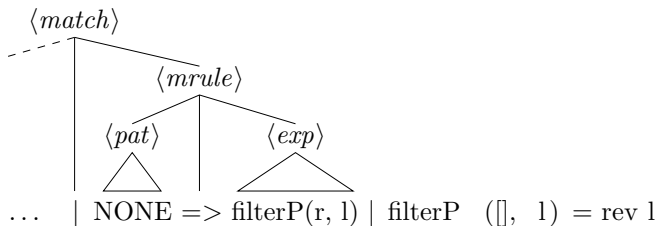
Which action to choose? Shift?



The Issue

Shift/Reduce Conflict

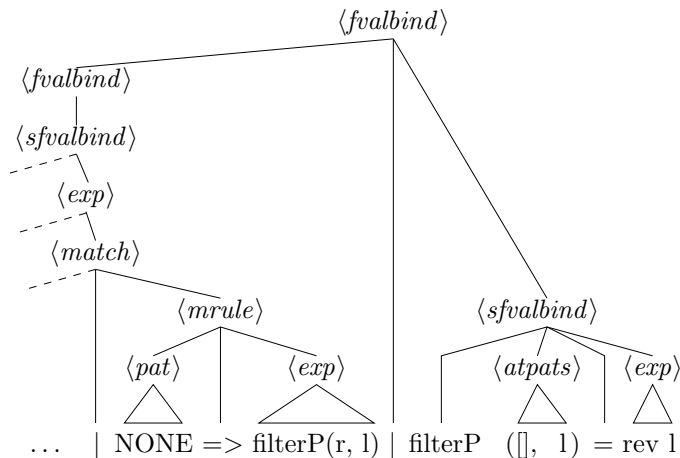
Which action to choose?



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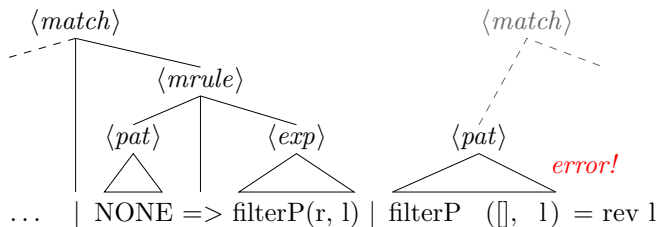
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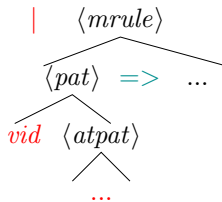
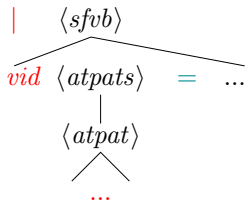
Which action to choose? Shift?



The Issue

Shift/Reduce Conflict

Which action to choose? Lookahead?



An unbounded lookahead is needed.

The Issue

Shift/Reduce Conflict

Which action to choose?

Example

Ambiguous construct:

case a of b \Rightarrow **case b of c** \Rightarrow c | d \Rightarrow d

The Issue

Shift/Reduce Conflict

Which action to choose? Reduce?

Example

Ambiguous construct:

```
case a of
  b => case b of
        c => c
  | d => d
```

The Issue

Shift/Reduce Conflict

Which action to choose? Shift?

Example

Ambiguous construct:

```
case a of
  b => case b of
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```

Here, the standard favors the shift.

Generalized LR

Tomita [1986]

- ▶ checks every possibility
- ▶ considered fast enough
- ▶ painless

Generalized LR

Tomita [1986]

- ▶ checks every possibility
- ▶ considered fast enough
- ▶ painless... or is it?

The Sword of Damocles

Do we catch all the ambiguities?

Example

case a of b \Rightarrow b | c \Rightarrow **case c of d** \Rightarrow d | e \Rightarrow e

The Sword of Damocles

Do we catch all the ambiguities?

Example

case a of

 b \Rightarrow b

| c \Rightarrow **case c of**

 d \Rightarrow d

 | e \Rightarrow e

The Sword of Damocles

Do we catch all the ambiguities?

Example

case a of

 b \Rightarrow b

| c \Rightarrow **case c of**

 d \Rightarrow d

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The Sword of Damocles

Do we catch all the ambiguities?

Example

case a **of** b \Rightarrow b | c \Rightarrow **case** c **of** d \Rightarrow d | e \Rightarrow e

The previous filter is inoperant.

- ▶ warnings in SDF2 and Elkhound, optionally aborts
- ▶ abort in GNU Bison

Detecting Ambiguities

How to find conservatively all the ambiguities in a grammar?

For the whole SML grammar:

- ▶ conflicts in the LALR(1) parser

sml.y: conflicts: 223 shift/reduce, 35 reduce/reduce

- ▶ ?

- ▶ Our tool:

89 potential ambiguities with LR(1) precision detected

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The Algorithm

Results on the SML grammar fragment

2 potential ambiguities with LR(0) precision detected:

```
(match -> mrule . , match -> match . '|' mrule )
```

```
(match -> match . '|' mrule , match -> match '|' mrule . )
```

Principle

- ▶ pairs of items
- ▶ synchronized to read the same prefix
- ▶ find pairs that might reach the end of input simultaneously

The Algorithm

Results on the SML grammar fragment

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- ▶ pairs of items
- ▶ synchronized to read the same prefix
- ▶ find pairs that might reach the end of input simultaneously

Example Run

$(\langle match \rangle \rightarrow \langle match \rangle \cdot ' | ' \langle mrule \rangle, \langle exp \rangle \rightarrow \mathbf{case} \langle exp \rangle \mathbf{of} \langle match \rangle \cdot)$

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$(\langle match \rangle \rightarrow \langle match \rangle \cdot ' | ' \langle mrule \rangle, \langle exp \rangle \rightarrow \mathbf{case} \langle exp \rangle \cdot \mathbf{of} \langle match \rangle)$

$(\langle match \rangle \rightarrow \langle match \rangle \cdot ' | ' \langle mrule \rangle, \langle sfvb \rangle \rightarrow \mathit{vid} \langle atpats \rangle = \langle exp \rangle \cdot)$

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$(\langle match \rangle \rightarrow \langle match \rangle \cdot ' | ' \langle mrule \rangle, \langle match \rangle \rightarrow \langle match \rangle ' | ' \langle mrule \rangle \cdot)$

Our two potential ambiguities

Example Run

$$(\langle match \rangle \rightarrow \langle match \rangle \cdot ' | ' \langle mrule \rangle, \langle match \rangle \rightarrow \langle mrule \rangle \cdot)$$
$$(\langle match \rangle \rightarrow \langle match \rangle \cdot ' | ' \langle mrule \rangle, \langle match \rangle \rightarrow \langle match \rangle ' | ' \langle mrule \rangle \cdot)$$

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Overview of the Algorithm

Schmitz [2006]

Computation of $(\mathbf{mac} \cup \mathbf{mas} \cup \mathbf{mae})^*$ from the initial item pair

$$(S' \rightarrow .S\$, \quad S' \rightarrow .S\$)$$

mac from a pair in conflict to its upwards closure,
mas by shifting a symbol,
mae by downwards closure.

Complexity: $\mathcal{O}(|R|^2)$ with $|R|$ the number of relations for individual items.

LR(0), SLR(1) $\mathcal{O}(|\mathcal{G}|^2)$

LR(1) $\mathcal{O}(|\mathcal{G}|^2 |T|^4)$

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- mas** by shifting a symbol,
- mae** by downwards closure.

Complexity: $\mathcal{O}(|R|^2)$ with $|R|$ the number of relations for individual items.

LR(0), SLR(1) $\mathcal{O}(|\mathcal{G}|^2)$

LR(1) $\mathcal{O}(|\mathcal{G}|^2 |T|^4)$

Some Experiments

Number of conflicts reported by Bison, and number of initial LR(0) conflicting pairs remaining with the LR, LRR and NU tests employing successively LR(0), SLR(1) and LR(1) items.

Grammar	Bison	LR(0) items			SLR(1) items			LR(1) items		
		LR	LRR	NU	LR	LRR	NU	LR	LRR	NU
Pascal	1	243	56	56	5	5	5	1	1	1
Mini C	1	300	16	15	5	5	4	1	1	1
ANSI C	1	377	14	3	14	14	3	1	1	1
ANSI C'	38	387	43	32	43	43	32	29	-	-
Standard ML	258	477	299	271	165	163	135	127	125	104
Small Elsa C++	58	1379	278	226	66	63	58	64	-	-
Elsa C++	115	2094	323	308	91	88	77	-	-	-

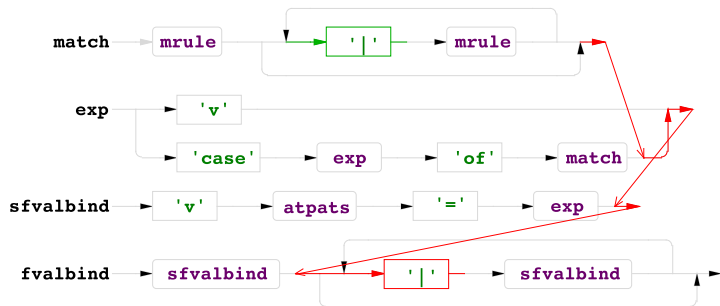
Limitations

Ambiguity Report

► grambiguity [Brabrand et al., 2006]

*** horizontal ambiguity at E[plus]: $\text{Exp} \leftrightarrow '+' \text{Exp}$
 ambiguous string: "x+x+x"

► ANTLRWorks [Parr, 2007]



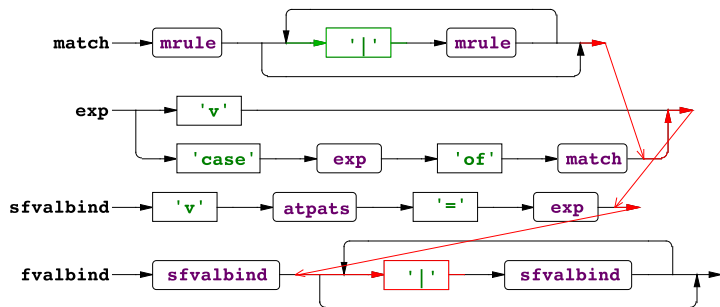
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*** horizontal ambiguity at E[plus]: $\text{Exp} \leftrightarrow '+' \text{Exp}$
 ambiguous string: "x+x+x"

- ▶ ANTLRWorks [Parr, 2007]



Other Limitations

- ▶ memory requirements: a solution could be a NLALR test
- ▶ dynamic disambiguation: inverse problem, some means to deciding equivalence needed

Closing Comments

- ▶ ambiguity hampers the reliability of grammarware
- ▶ conservative detection techniques can be improved
- ▶ prototype tool
- ▶ many, many open questions left

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Dynamic Disambiguation Filters

$$A \rightarrow aBc \mid aaBc, B \rightarrow ab \mid b \quad (\mathcal{G}_1)$$

