Algorithmic Complexity of Well-Quasi-Orders

Sylvain Schmitz

Séminaire Automates, IRIF

February 9, 2018

Outline

well-quasi-orders (wqo):

- proving algorithm termination

thesis: a toolbox for wqo complexity

- upper bounds
- lower bounds
- complexity classes

this talk: focus on one problem

- reachability in vector addition systems

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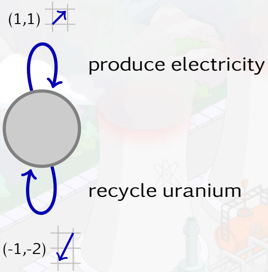
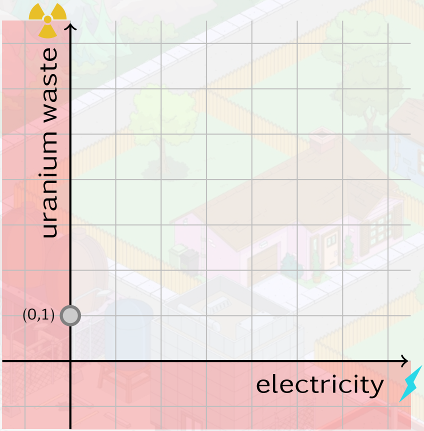
Vector Addition Systems

Vector Addition Systems

Vector Addition Systems

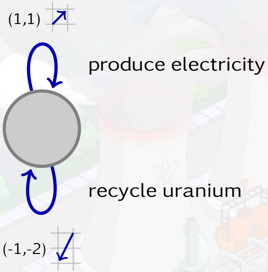
Springfield Power Plant

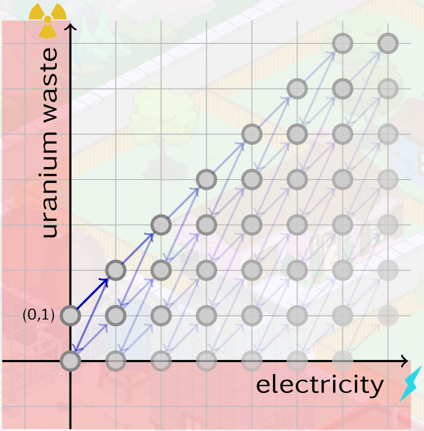




Vector Addition Systems

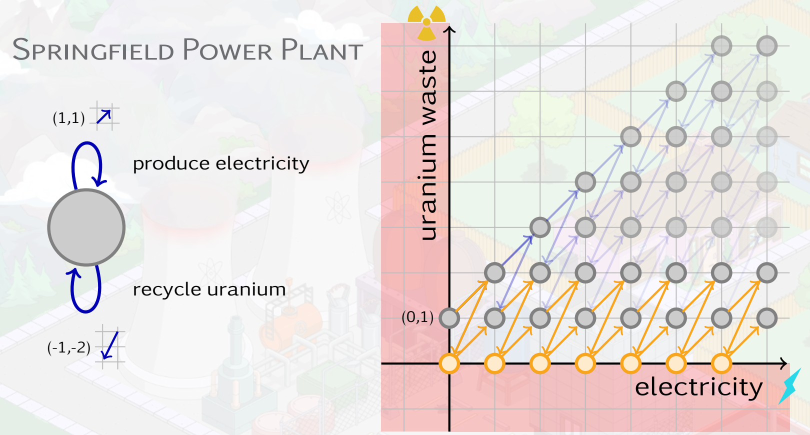
Springfield Power Plant





Can we produce unbounded electricity with no left-over uranium waste?

Vector Addition Systems



Can we produce unbounded electricity with no left-over uranium waste? Yes, $(\infty,0)$ is reachable

Importance of the Problem

Reachability Problem

input: *a vector addition system and two configurations* **source** *and* **target**

question: **source** $\rightarrow^*$ **target***?*

Discrete Resources

- modelling: items, money, energy, molecules, ...
- distributed computing: active threads in thread pool
- data: isomorphism types in data logics and data-centric systems

Importance of the Problem

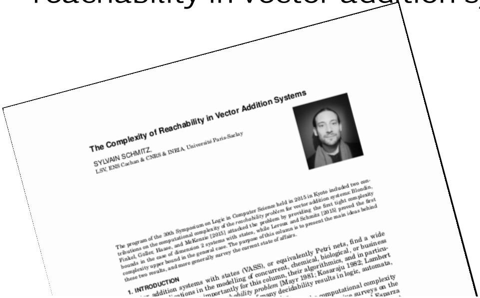
Reachability Problem

input: *a vector addition system and two configurations* **source** *and* **target**

question: **source** $\rightarrow^*$ **target**?

Central Decision Problem [invited survey S., SIGLOG'16]

Large number of problems interreducible with reachability in vector addition systems



The Complexity of Reachability in Vector Addition Systems

SYLVAIN SCHMITZ,
LSV, ENS Cachan & CNRS & INRIA, Université Paris-Saclay

The program of the 30th Symposium on Logic in Computer Science held in 2015 in Kyoto included two contributions on the computational complexity of the *reachability problem* for vector addition systems: Blondin, Finkel, Goller, Haase, and McKenzie [2015] attacked the problem by providing the first tight complexity bounds in the case of dimension 2 systems with states, while Leroux and Schmitz [2015] proved the first complexity upper bound in the general case. The purpose of this column is to present the main ideas behind these two results, and more generally survey the current state of affairs.

1. INTRODUCTION

... addition systems with states (VASS), or equivalently Petri nets, find a wide
... in the modelling of concurrent, chemical, biological, or business
... importantly for this column, their algorithmics, and in particu-
... *ability problem* [Mayr 1981; Kosaraju 1982; Lambert
... many decidability results in logic, automata,

Importance of the Problem

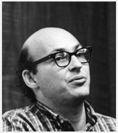
Reachability Problem

input: *a vector addition system and two configurations* **source** *and* **target**

question: **source** $\rightarrow^*$ **target**?

Theorem (Minsky'67)

Reachability is undecidable in 2-dimensional Minsky machines (vector addition systems with zero tests).



Importance of the Problem

1962  C. A. Petri: Petri nets

1969   R. M. Karp & R. E. Miller: coverability trees

1976  R. J. Lipton: EXPSPACE lower bound

1979   J. E. Hopcroft & J.-J. Pansiot: dim. $\geqslant 3$ not definable in Presburger arithmetic

1981  E. W. Mayr: decidability by decomposition

1982  S. R. Kosaraju: decidability by decomposition

1992  J.-L. Lambert: decidability by decomposition

2011  J. Leroux: decidability by Presburger inductive invariants

2015 this talk: Leroux & S., LICS'15

Demystifying Reachability in Vector Addition Systems

[Leroux & S., LICS'15]

Upper Bound Theorem

Reachability in vector addition systems is in cubic Ackermann.

Ideal Decomposition Theorem

The Decomposition Algorithm computes the ideal decomposition of the set of runs from source to target.

DEMYSTIFYING REACHABILITY IN VECTOR ADDITION SYSTEMS

[Leroux & S., LICS'15; S., 2017]

UPPER BOUND THEOREM

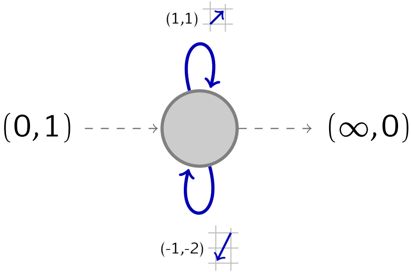
Reachability in vector addition systems is in quadratic Ackermann.

IDEAL DECOMPOSITION THEOREM

The Decomposition Algorithm computes the ideal decomposition of the set of runs from source to target.

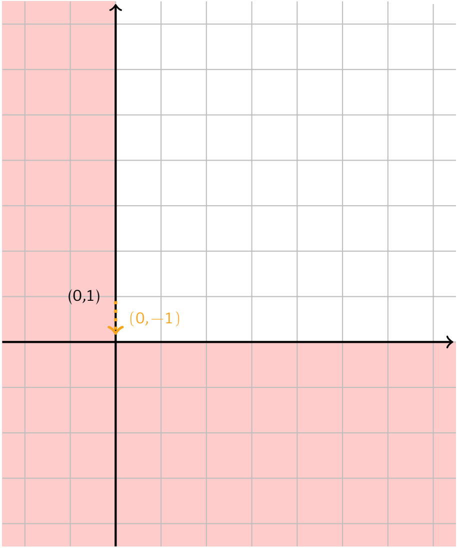
Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]



Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]



Decomposition Algorithm

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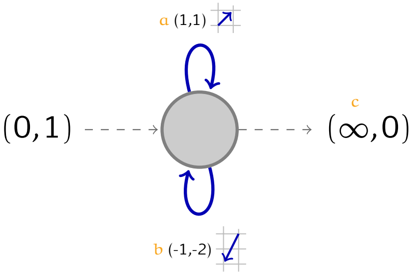
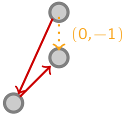


Equations

$$0 + 1 \cdot a - 1 \cdot b = c$$
$$1 + 1 \cdot a - 2 \cdot b = 0$$

Solution Path

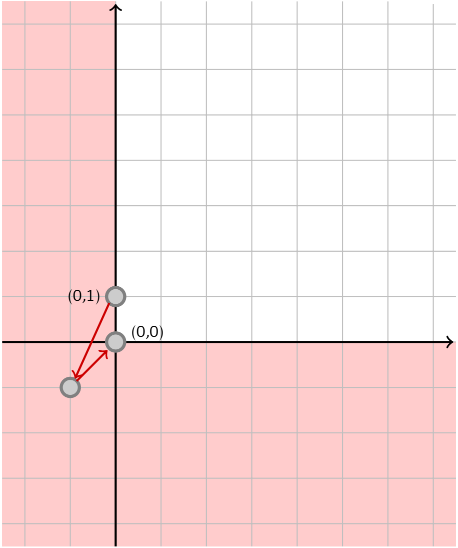


Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

solution path

$\times 1$



Decomposition Algorithm

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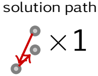

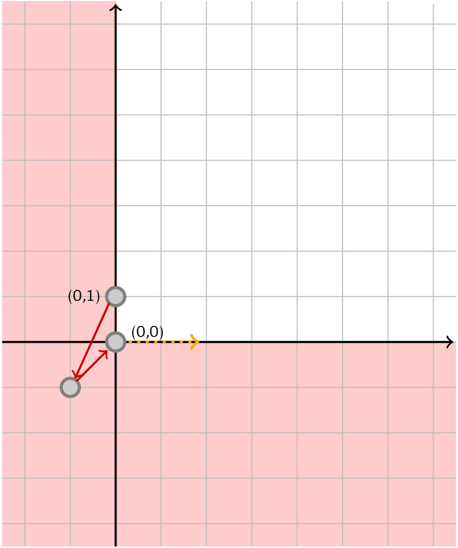

Decomposition Algorithm

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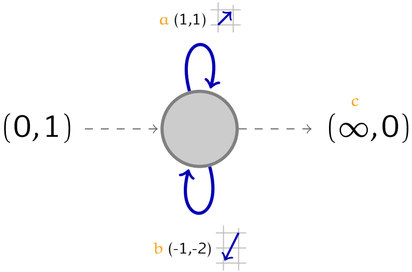
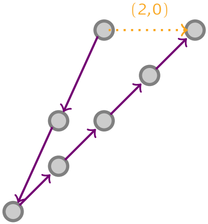


Equations

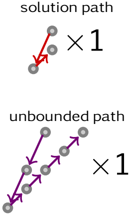
$$1 \cdot a - 1 \cdot b = c$$
$$1 \cdot a - 2 \cdot b = 0$$
$$a, b, c > 0$$

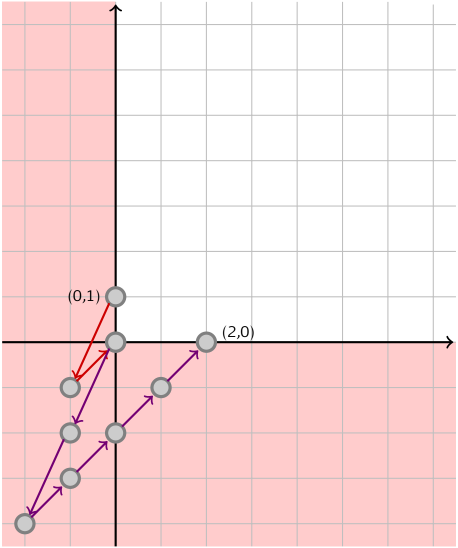
Unbounded Path



Decomposition Algorithm

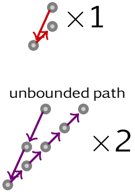
[Mayr'81, Kosaraju'82, Lambert'92]

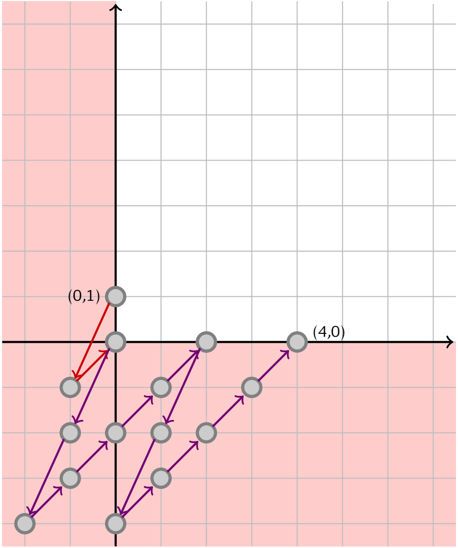




Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]





Decomposition Algorithm

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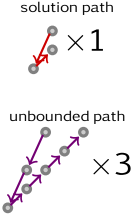
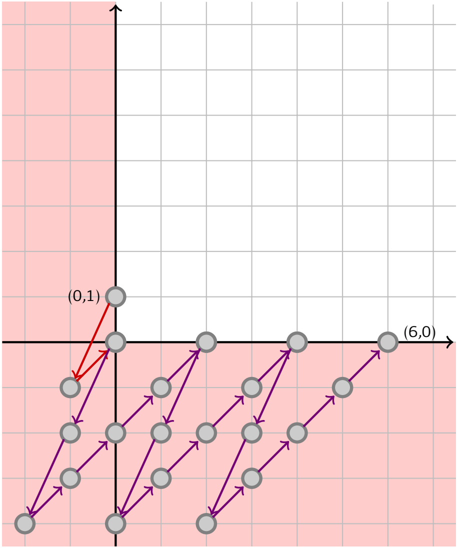
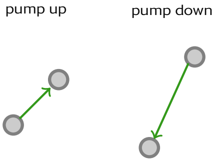




Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

Pumpable Paths



Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

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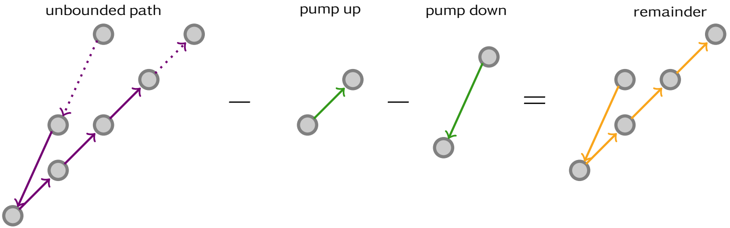
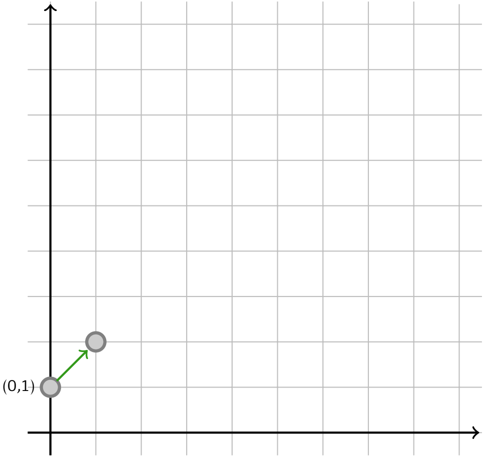


Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

pump up

 $\times 1$

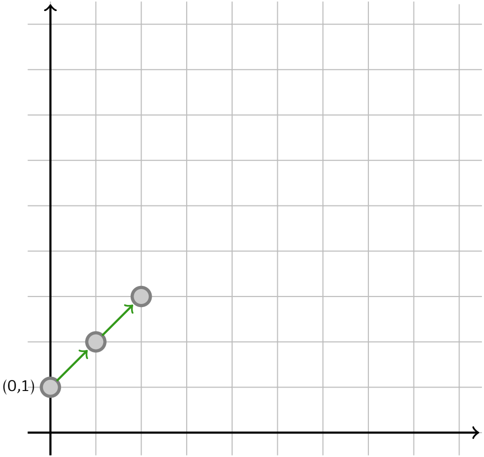


Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

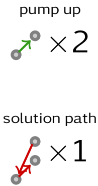
pump up

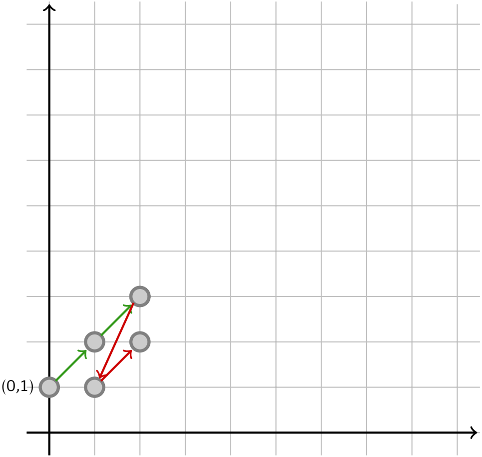
$\times 2$



Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]





Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

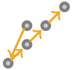
pump up

×2

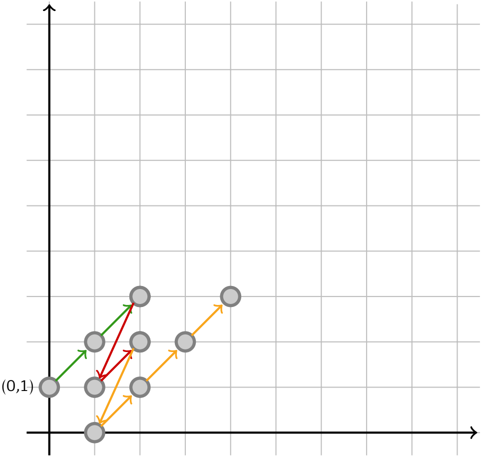
solution path

×1

remainder

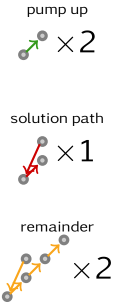


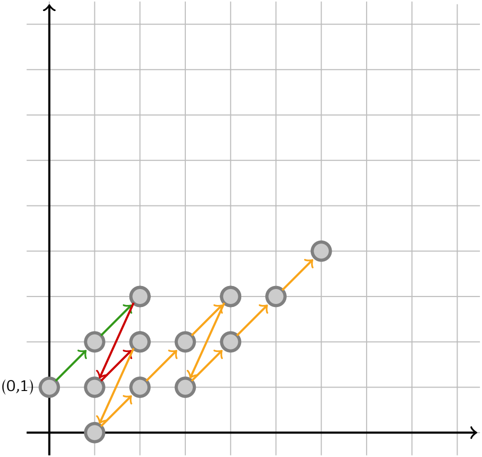
×1



Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]





Decomposition Algorithm

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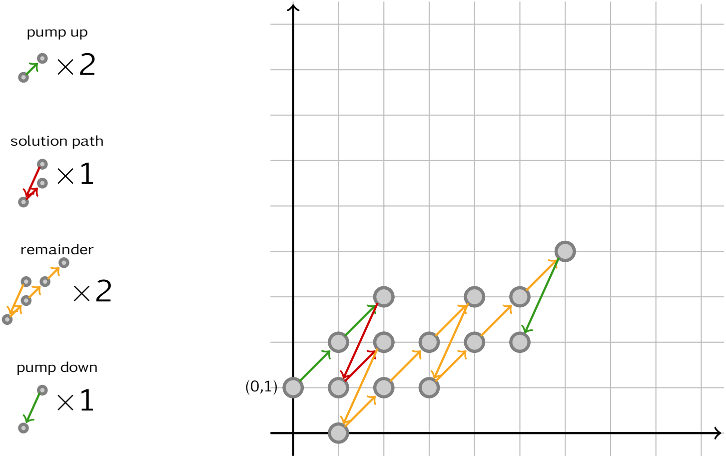

Decomposition Algorithm

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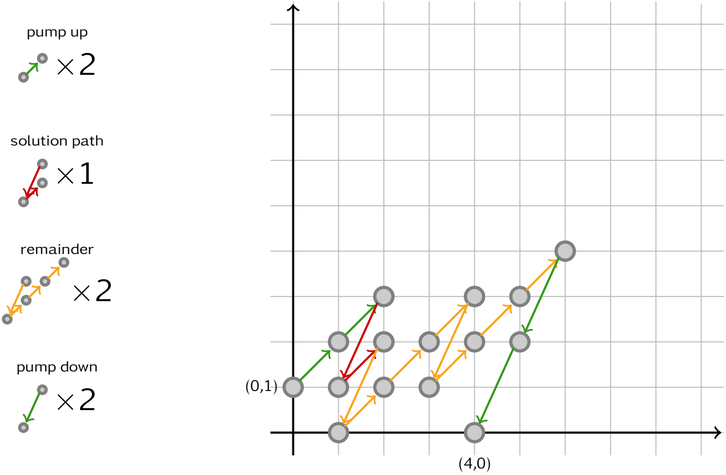


Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

pump up

$\times 3$

solution path

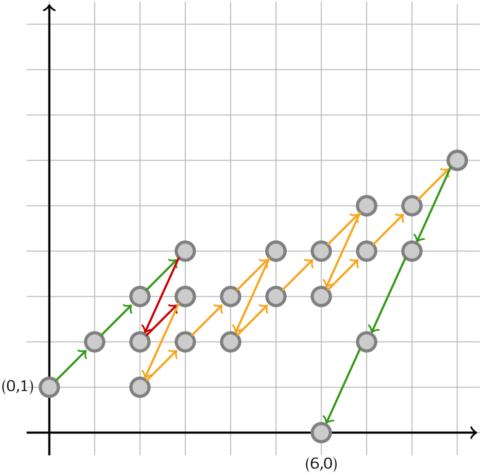
$\times 1$

remainder

$\times 3$

pump down

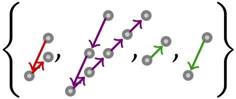
$\times 3$



Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

can we build a simple run?



Decomposition Algorithm

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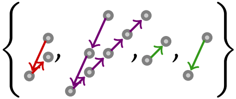


"Θ Condition"

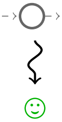
in ExpSpace

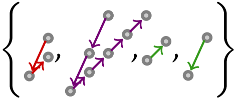
[e.g. Rackoff'78, Demri'13, Blockelet & S., MFCS'11]

Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

can we build a simple run? yes





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Decomposition Algorithm

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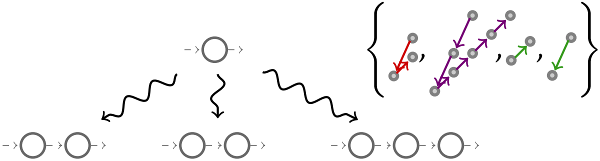
can we build a simple run? no

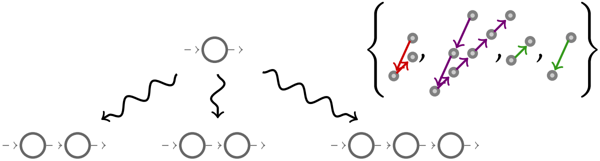


decompose

Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

can we build a simple run? no



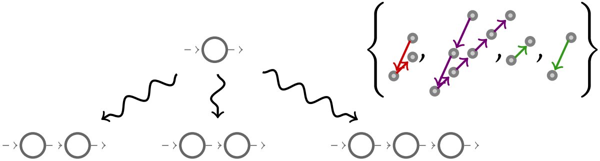
decompose

uses *coverability trees* [Karp & Miller'69]

Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

can we build a simple run? no



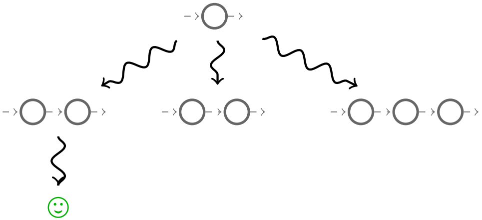
decompose

uses *coverability trees* [Karp & Miller'69]

which use *Dickson's Lemma* [Dickson, 1913]

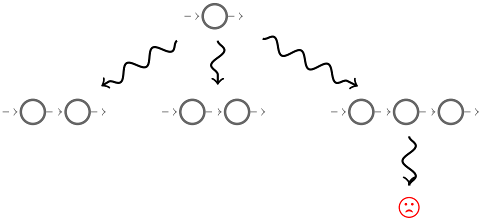
Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]



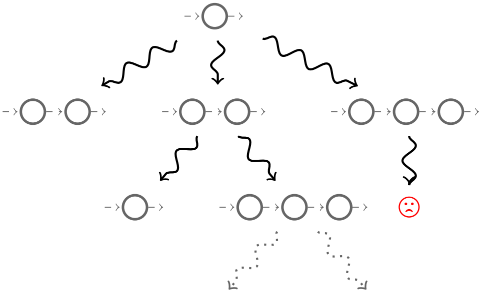
Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]



Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]



Termination

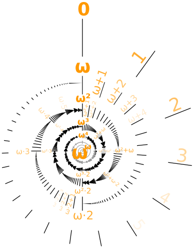
"Finally the checker has to verify that the process comes to an end. Here again he should be assisted by the programmer giving a further definite assertion to be verified. This may take the form of a quantity which is asserted to decrease continually and vanish when the machine stops."

[Turing'49]

Termination

"Finally the checker has to verify that the process comes to an end. Here again he should be assisted by the programmer giving a further definite assertion to be verified. This may take the form of a quantity which is asserted to decrease continually and vanish when the machine stops. To the pure mathematician it is natural to give an ordinal number."

[Turing'49]



Termination of the Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

Ranking Function

$$\omega^{\omega^2}$$

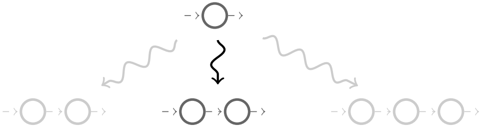
$$\vee$$

$$\alpha_0$$

Termination of the Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

Ranking Function



$$\omega^{\omega^2}$$

$$\vee$$

$$\alpha_0$$

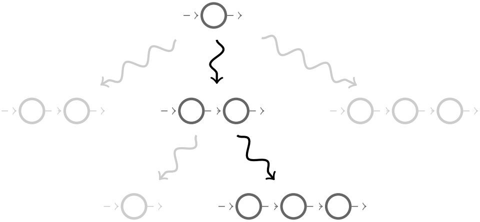
$$\vee$$

$$\alpha_1$$

Termination of the Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

Ranking Function

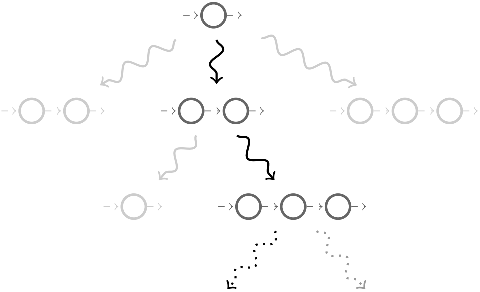


$$\omega^{\omega^2} > \alpha_0 > \alpha_1 > \alpha_2$$

Termination of the Decomposition Algorithm

[Mayr'81, Kosaraju'82, Lambert'92]

Ranking Function



$$\omega^{\omega^2} > \alpha_0 > \alpha_1 > \alpha_2 > \vdots$$

Demystifying Reachability in Vector Addition Systems

[Leroux & S., LICS'15; S., 2017]

Upper Bound Theorem
Reachability in vector addition systems is in quadratic Ackermann.

Ideal Decomposition Theorem
The Decomposition Algorithm computes the ideal decomposition of the set of runs from source to target.

Upper Bounds

How to bound the running time of algorithms with ordinal-based termination proofs?

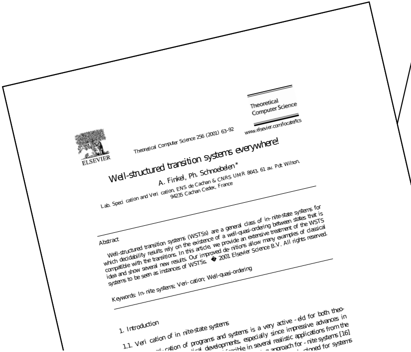
Upper Bounds

How to bound the running time of algorithms with wqo-based termination proofs?

Upper Bounds

How to bound the running time of algorithms with wqo-based termination proofs?

wqos ubiquitous in infinite-state verification


Theoretical Computer Science 256 (2001) 63–92

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Well-structured transition systems everywhere!

A. Finkel, Ph. Schnoebelen*

Lab. Specification and Verification, ENS de Cachan & CNRS UMR 8643, 61 av. Pdt Wilson, 94235 Cachan Cedex, France

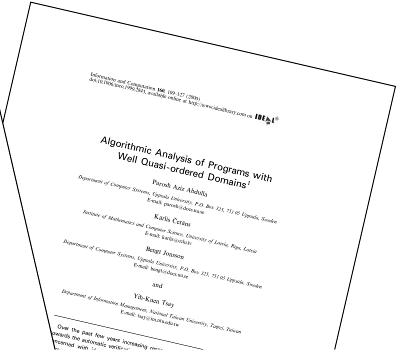
Abstract

Well-structured transition systems (WSTSs) are a general class of infinite-state systems for which decidability results rely on the existence of a well-quasi-ordering between states that is compatible with the transitions. In this article, we provide an extensive treatment of the WSTS idea and show several new results. Our improved definitions allow many examples of classical systems to be seen as instances of WSTSs. © 2001 Elsevier Science B.V. All rights reserved.

Keywords: Infinite systems; Verification; Well-quasi-ordering

1. Introduction

1.1. Verification of infinite-state systems


Information and Computation 160, 109–127 (2000)
doi:10.1006/inco.1999.2843, available online at http://www.idealibrary.com on IDEAL

Algorithmic Analysis of Programs with Well Quasi-ordered Domains[1]

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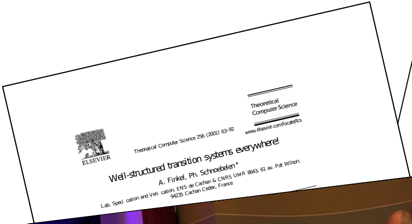
Upper Bounds

How to bound the running time of algorithms with wqo-based termination proofs?

wqos ubiquitous in infinite-state verification


Theoretical Computer Science 256 (2001) 63–92

Theoretical Computer Science

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Well-structured transition systems everywhere!

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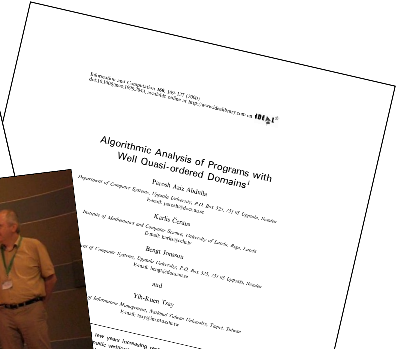
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Bad Sequences

Over a qo $(X, \leqslant)$

- $x_0, x_1, \ldots$ is bad if $\forall i < j \,.\, x_i \not\leqslant x_j$
- $(X, \leqslant)$ wqo iff all bad sequences are finite
- but can be of arbitrary length

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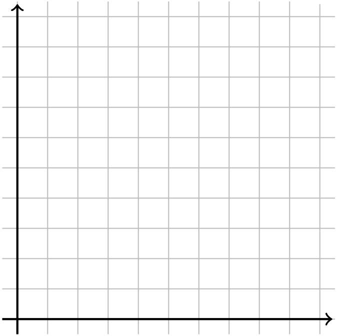

Example (over $\mathbb{N}^2$)

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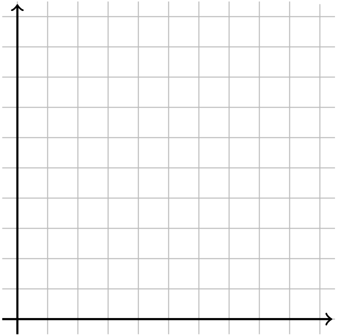


Example (over $\mathbb{N}^2$)

Controlled Bad Sequences

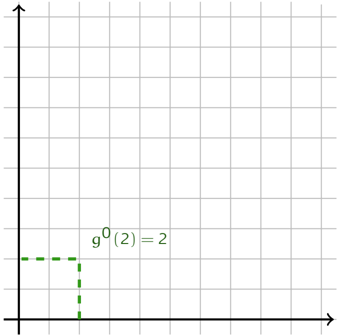
Over a qo $(X, \leqslant)$ with norm $\|\cdot\|$

$x_0, x_1, \ldots$ is bad if $\forall i < j \,.\, x_i \not\leqslant x_j$

$(X, \leqslant)$ wqo iff all bad sequences are finite

controlled by $g\colon \mathbb{N} \to \mathbb{N}$ and $n \in \mathbb{N}$ if $\forall i \,.\, \|x_i\| \leqslant g^i(n)$

[Cichoń & Tahhan Bittar'98]



Example (over $\mathbb{N}^2$ with $n = 2$ and $g(n) = n + 1$)

Controlled Bad Sequences

Over a qo $(X, \leqslant)$ with norm $\|\cdot\|$

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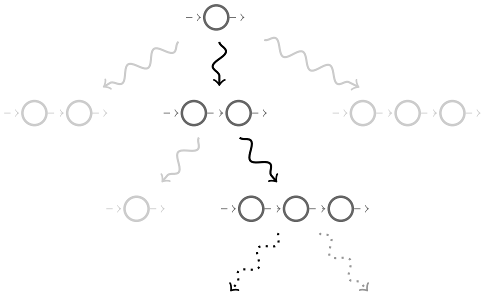
[Cichoń & Tahhan Bittar'98]

Proposition
Assuming $\{x \in X \mid \|x\| \leqslant n\}$ finite $\forall n$, controlled bad sequences have bounded length.

The Length of Descending Sequences



$$\alpha_0 > \alpha_1 > \alpha_2 > \cdots$$

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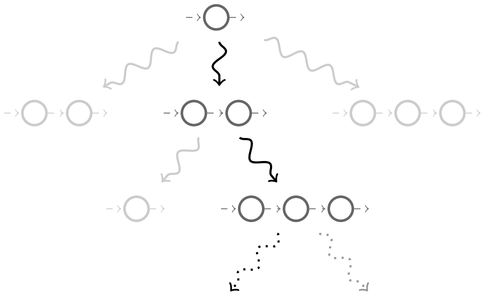

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Length Function Theorem (for Ordinals [invited talk S., RP'14])

Descending sequences over ω^{ω^2} controlled by Ackermannian functions are of at most quadratic Ackermannian length.

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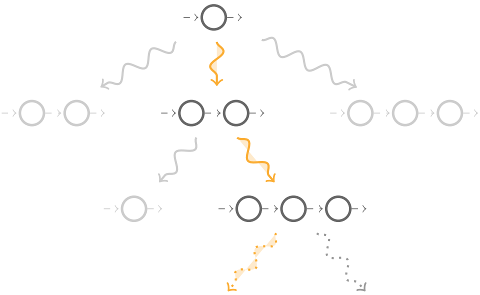

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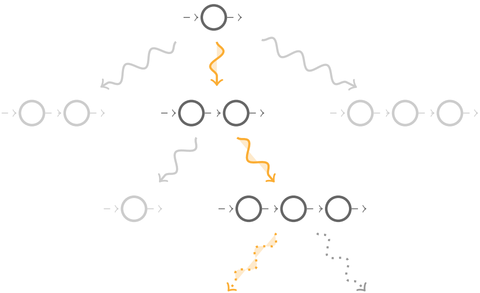

$$\alpha_0 > \alpha_1 > \alpha_2 > \vdots$$

Length Function Theorem (for Dickson's Lemma [Figueira, Figueira, S. & Schnoebelen, LICS'11])

Bad sequences over $\mathbb{N}^d$ controlled by primitive recursive functions are of at most Ackermannian length.

Fast-Growing Functions

Ackermann Function

$$A(1,n) = 2n$$
$$A(2,n) = 2^n$$
$$A(3,n) = \text{tower}(n) \stackrel{\text{def}}{=} 2^{\cdot^{\cdot^{\cdot^{2}}}}\Big\}n \text{ times}$$
$$\vdots$$

- $\text{ackermann}(n) \stackrel{\text{def}}{=} A(n,n)$ not primitive recursive
- quadratic Ackermann function F_{ω^2}: 3-arguments variant

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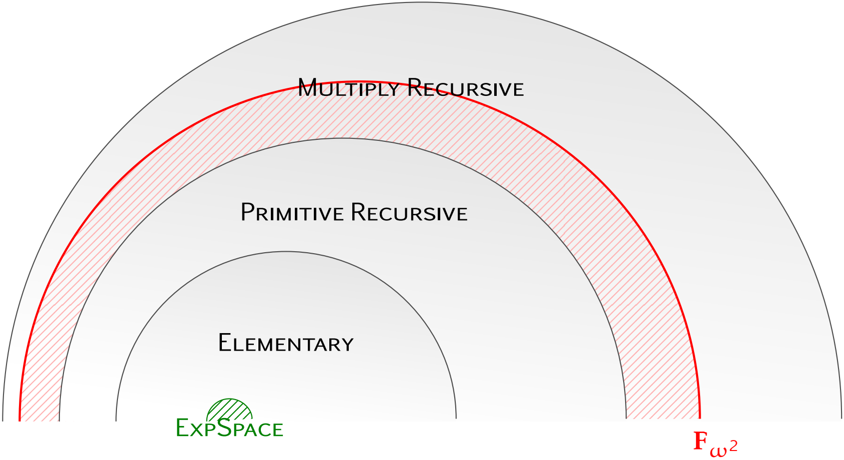
Complexity Classes Beyond Elementary

[S., ToCT'16]



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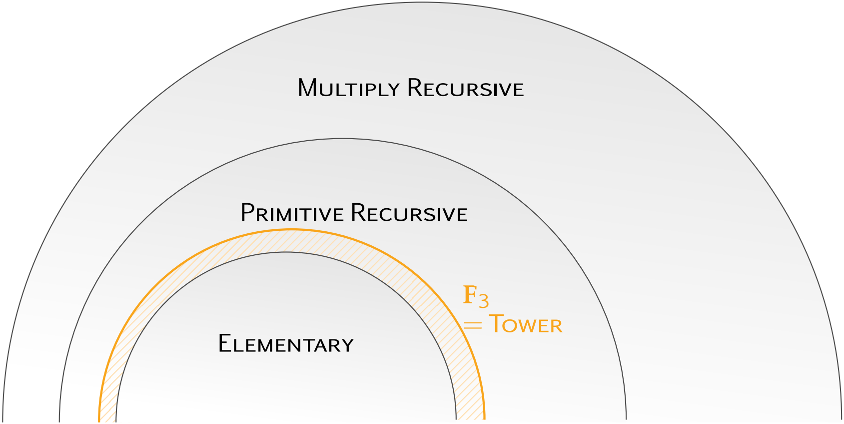

$$\mathbf{F}_3 \stackrel{\text{def}}{=} \bigcup_{e \text{ elementary}} \mathrm{DTime}\big(\mathrm{tower}(e(n))\big)$$

Complexity Classes Beyond Elementary

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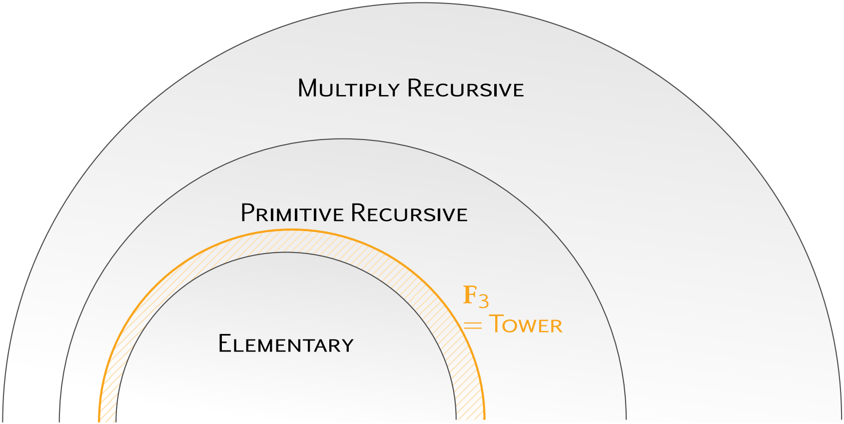
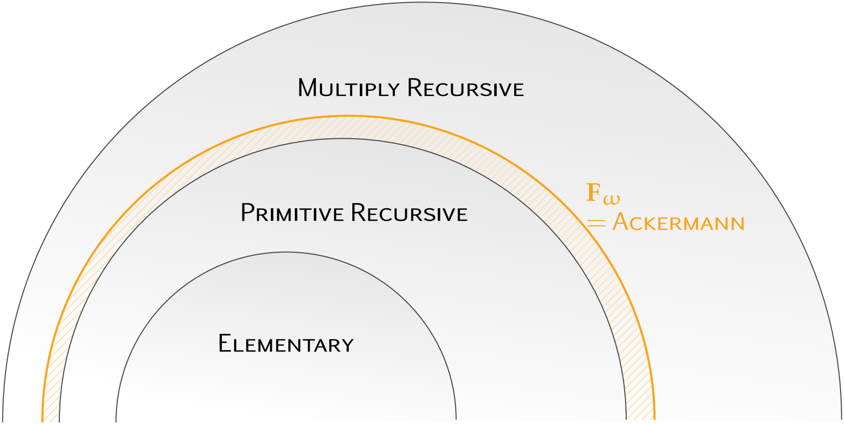


Examples of Tower-Complete Problems:

- satisfiability of first-order logic on words [Meyer'75]
- β-equivalence of simply typed λ terms [Statman'79]
- model-checking higher-order recursion schemes [Ong'06]

Complexity Classes Beyond Elementary

[S., ToCT'16]



$$\mathbf{F}_\omega \stackrel{\text{def}}{=} \bigcup_{p \text{ primitive recursive}} \text{DTime}\big(\text{ackermann}(p(n))\big)$$

Complexity Classes Beyond Elementary

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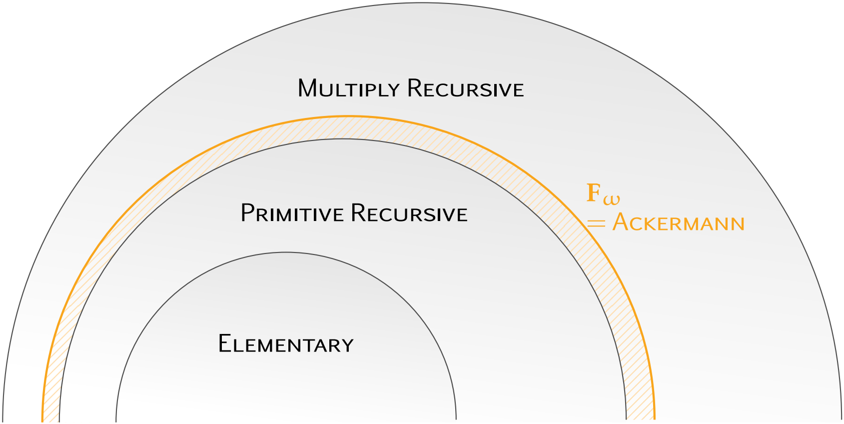

Examples of Ackermann-Complete Problems:

- reachability in lossy Minsky machines [Urquhart'98, Schnoebelen'02]
- satisfiability of safety Metric Temporal Logic [Lazić et al.'16]
- satisfiability of Vertical XPath [Figueira and Segoufin'17]

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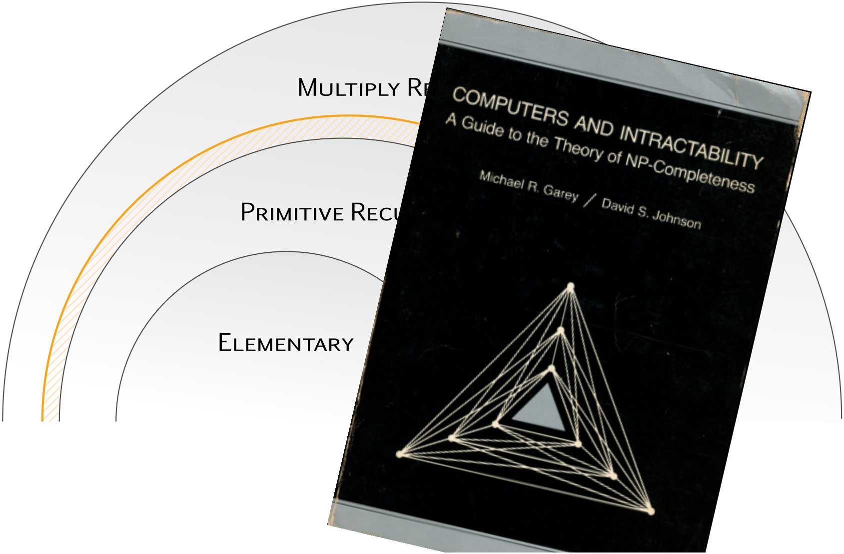

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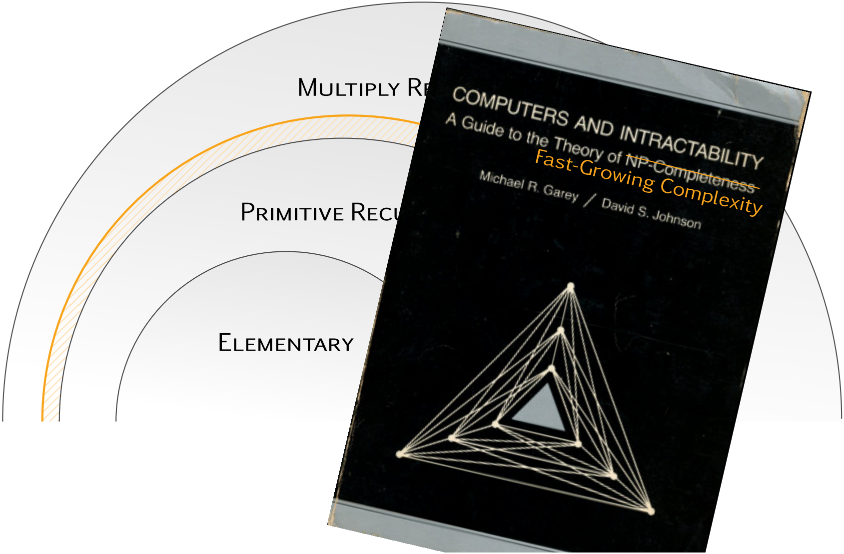


SUMMARY

well-quasi-orders (wqo):

- proving algorithm termination

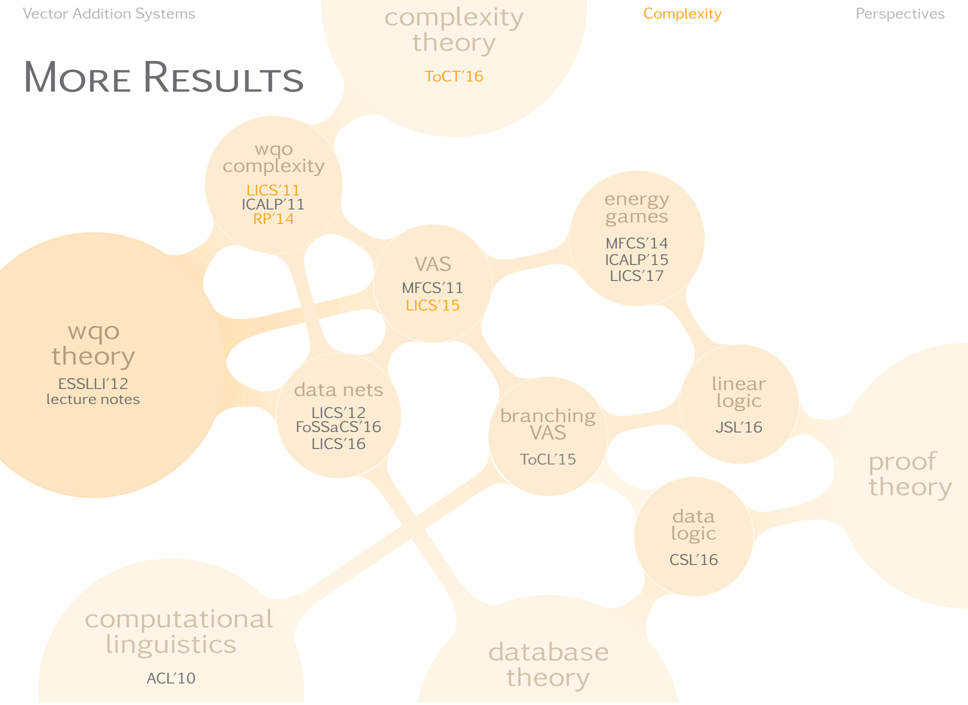
thesis: a toolbox for wqo complexity

- upper bounds: length function theorems (for ordinals, Dickson's Lemma, Higman's Lemma, and combinations)
- lower bounds
- complexity classes: $(\mathbf{F}_\alpha)_\alpha$

this talk: focus on one problem

- reachability in vector addition systems in $\mathbf{F}_{\omega^2}$

More Results



Perspectives

1. complexity gap for VAS reachability
   - ExpSpace-hard [Lipton'76]
   - decomposition algorithm: at least $\mathbf{F}_\omega$ (Ackermannian) time

2. parameterisations for counter systems
   - the dimension is the main source of complexity
   - find better parameters with tight bounds? [Kristiansen & Niggl'04]

3. beyond wqos: FAC qos, Noetherian spaces [Goubault-Larrecq'06]
   - complexity?

4. reachability in VAS extensions
   - decidable in VAS with hierarchical zero tests [Reinhardt'08]
   - what about
     - branching VAS
     - unordered data Petri nets
     - pushdown VAS

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Demystifying Reachability in Vector Addition Systems

[Leroux & S., LICS'15; S., 2017]

Upper Bound Theorem
Reachability in vector addition systems is in quadratic Ackermann.

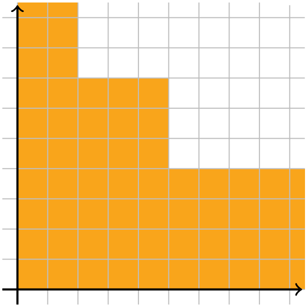
Ideal Decomposition Theorem
The Decomposition Algorithm computes the ideal decomposition of the set of runs from source to target.

Ideals of Well-Quasi-Orders $(X, \leqslant)$

- Canonical decompositions
  [Bonnet'75]
  if $D \subseteq X$ is $\downarrow$-closed, then

  $$D = I_1 \cup \cdots \cup I_n$$

  for (maximal) ideals $I_1, \ldots, I_n$



Example (over $\mathbb{N}^2$)

$$D = (\{0,\ldots,2\} \times \mathbb{N}) \cup (\{0,\ldots,5\} \times \{0,\ldots,7\}) \cup (\mathbb{N} \times \{0,\ldots,4\})$$

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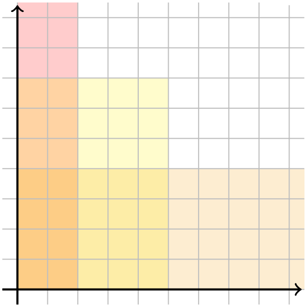

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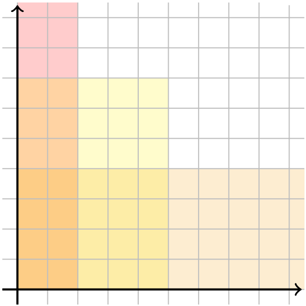
- Effective representations
  [Goubault-Larrecq et al.'17]



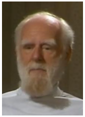
Example (over $\mathbb{N}^2$)

$$D = [\![(2, \infty)]\!] \cup [\![(5, 7)]\!] \cup [\![(\infty, 4)]\!]$$

Decomposition Theorem

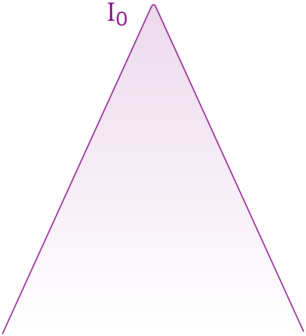
Well-Quasi-Order on Runs

combination of Dickson's and Higman's lemmata



Syntax

Semantics



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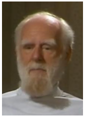

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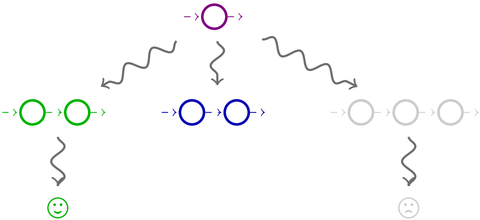

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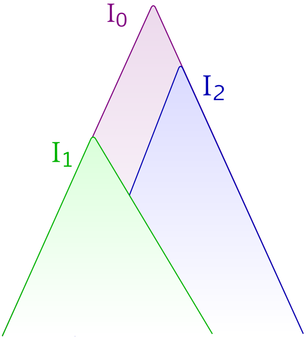

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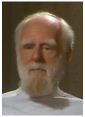

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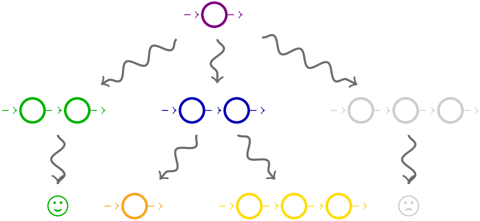

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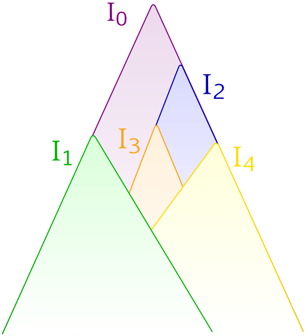

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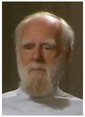

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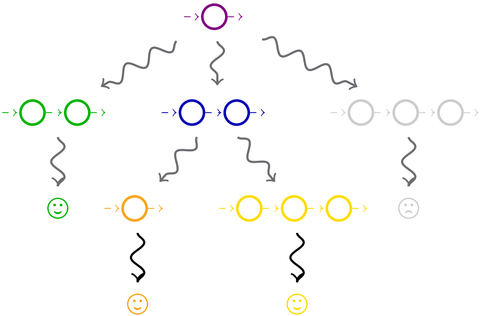

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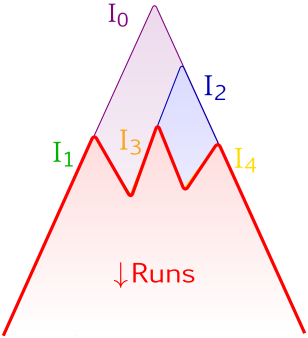
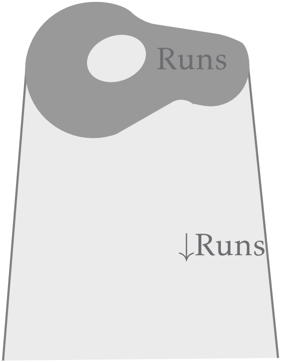


Adherence Membership

- I is adherent to Runs if $I \subseteq \downarrow(I \cap \text{Runs})$
- semantic equivalent to Θ condition
- undecidable for arbitrary ideals
- decidable for the ideals arising in the decomposition algorithm



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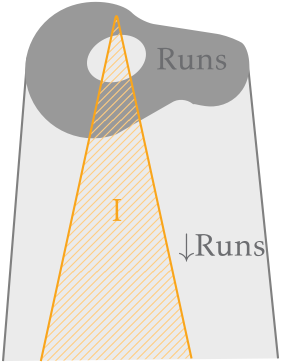

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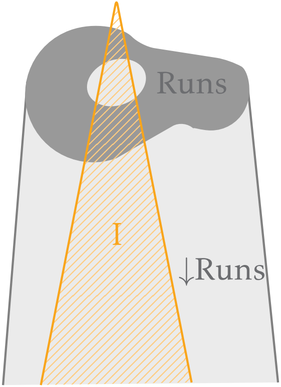

I not adherent

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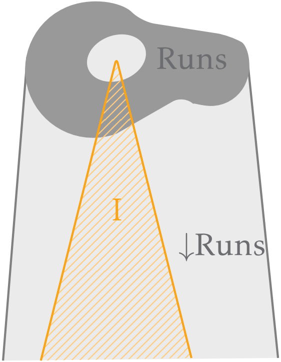

I not adherent

Adherence Membership

- I is adherent to Runs if $\mathrm{I} \subseteq \downarrow(\mathrm{I} \cap \mathrm{Runs})$
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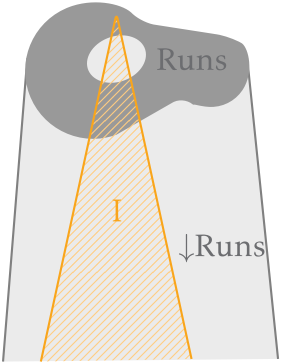

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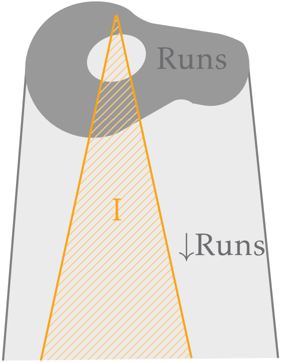

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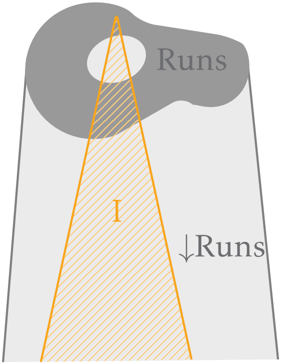

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Branching VAS Reachability

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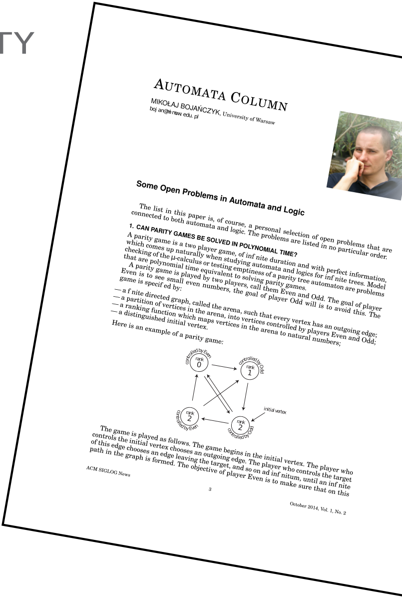
- application domains:



AUTOMATA COLUMN

MIKOŁAJ BOJAŃCZYK, University of Warsaw
boj an@mi muw. edu. pl

Some Open Problems in Automata and Logic

The list in this paper is, of course, a personal selection of open problems that are connected to both automata and logic. The problems are listed in no particular order.

1. CAN PARITY GAMES BE SOLVED IN POLYNOMIAL TIME?

A parity game is a two player game, of inf nite duration and with perfect information, which comes up naturally when studying automata and logics for inf nite trees. Model checking of the μ-calculus or testing emptiness of a parity tree automaton are problems that are polynomial time equivalent to solving parity games.

A parity game is played by two players, call them Even and Odd. The goal of player Even is to see small even numbers, the goal of player Odd will is to avoid this. The game is specif ed by:

— a f nite directed graph, called the arena, such that every vertex has an outgoing edge;
— a partition of vertices in the arena, into vertices controlled by players Even and Odd;
— a ranking function which maps vertices in the arena to natural numbers;
— a distinguished initial vertex.

Here is an example of a parity game:

The game is played as follows. The game begins in the initial vertex. The player who controls the initial vertex chooses an outgoing edge. The player who controls the target of this edge chooses an edge leaving the target, and so on ad inf nitum, until an inf nite path in the graph is formed. The objective of player Even is to make sure that on this

ACM SIGLOG News 3 October 2014, Vol. 1, No. 2

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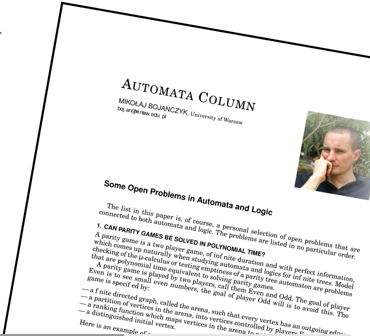
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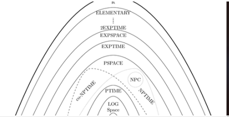
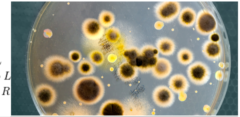
Complexity Theory	Distributed Computing	Proof Theory	Computational Biology



$$\dfrac{\dfrac{\dfrac{\dfrac{}{X \vdash X}\,\mathrm{ax} \qquad \dfrac{\dfrac{}{0 \vdash Y, Z}\,0L}{\vdash Y, 0 \multimap Z}\multimap R}{X \vdash Y, X \otimes (0 \multimap Z)}\otimes R}{\vdash X \multimap Y, X \otimes (0 \multimap Z)}\multimap R \qquad \dfrac{}{0 \vdash}\,0L}{\dfrac{(X \multimap Y) \multimap 0 \vdash X \otimes (0 \multimap Z)}{\vdash ((X \multimap Y) \multimap 0) \multimap (X \otimes (0 \multimap Z))}\multimap R}\multimap L$$



Programming Languages	Database Theory	Security	Computational Linguistics

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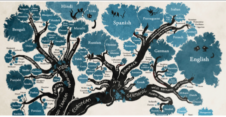

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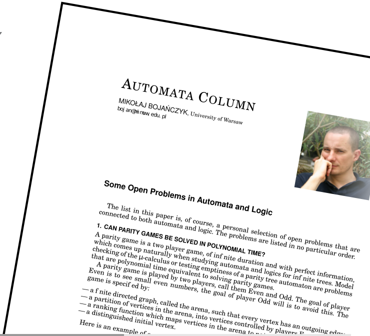
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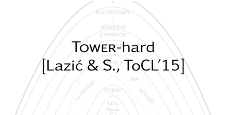
Complexity Theory

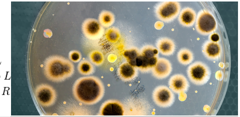


Distributed Computing

Proof Theory

$$\dfrac{\dfrac{\dfrac{X \vdash X \;\text{ax} \qquad \dfrac{\overline{0 \vdash Y, Z}\;0L}{\vdash Y, 0 \multimap Z}\multimap R}{X \vdash Y, X \otimes (0 \multimap Z)}\otimes R}{\vdash X \multimap Y, X \otimes (0 \multimap Z)}\multimap R \qquad \overline{0 \vdash}\;0L}{\dfrac{(X \multimap Y) \multimap 0 \vdash X \otimes (0 \multimap Z)}{\vdash ((X \multimap Y) \multimap 0) \multimap (X \otimes (0 \multimap Z))}\multimap R}\multimap L$$

Computational Biology



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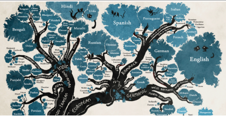
Database Theory

Security

Computational Linguistics



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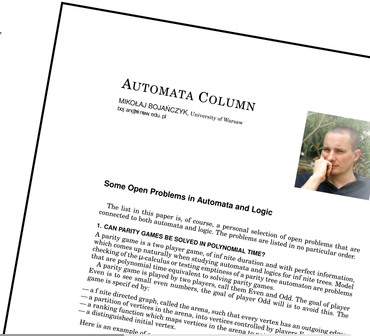
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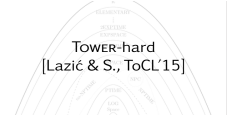
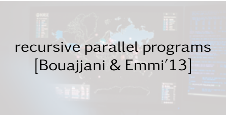
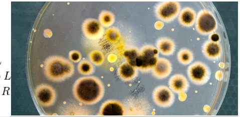
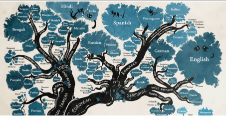
Complexity Theory	Distributed Computing	Proof Theory	Computational Biology
 Tower-hard [Lazić & S., ToCL'15]	 recursive parallel programs [Bouajjani & Emmi'13]	$\dfrac{\dfrac{X \vdash X}{}\text{ax} \quad \dfrac{\dfrac{}{0 \vdash Y, Z}0L}{\vdash Y, 0 \multimap Z}\multimap R}{\dfrac{X \vdash Y, X \otimes (0 \multimap Z)}{\vdash X \multimap Y, X \otimes (0 \multimap Z)}\multimap R}\otimes R \qquad \dfrac{}{0 \vdash}0L$; $\dfrac{(X \multimap Y) \multimap 0 \vdash X \otimes (0 \multimap Z)}{\vdash ((X \multimap Y) \multimap 0) \multimap (X \otimes (0 \multimap Z))}\multimap R$ (from $\multimap L$)	

Programming Languages	Database Theory	Security	Computational Linguistics
see code below			

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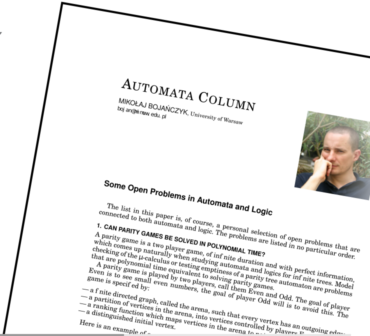
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Complexity Theory

Tower-hard
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Distributed Computing

recursive parallel programs
[Bouajjani & Emmi'13]

Proof Theory

linear and relevance logics
[de Groote et al.'04
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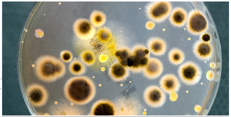
Computational Biology



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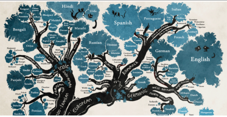
Database Theory

Security

Computational Linguistics



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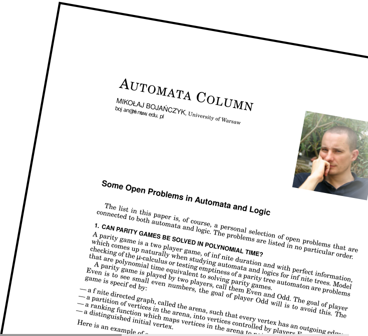
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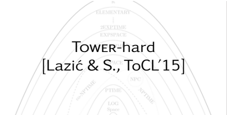
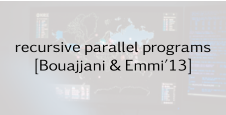


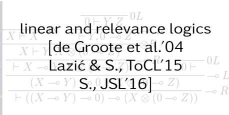
Tower-hard [Lazić & S., ToCL'15]

Distributed Computing



recursive parallel programs [Bouajjani & Emmi'13]

Proof Theory



linear and relevance logics [de Groote et al.'04 Lazić & S., ToCL'15 S., JSL'16]

Computational Biology

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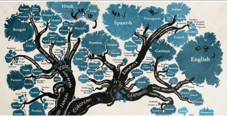

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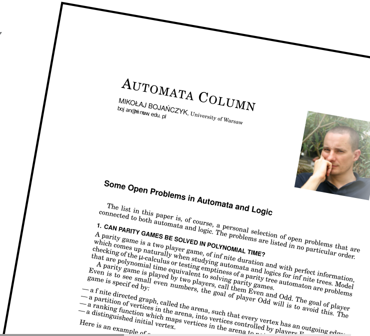
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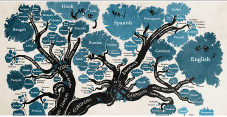
Programming Languages	Database Theory	Security	Computational Linguistics
observational equivalence [Cotton-Barratt et al.'17]			

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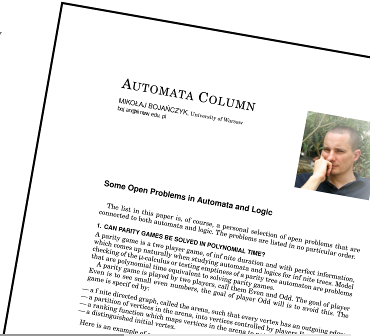
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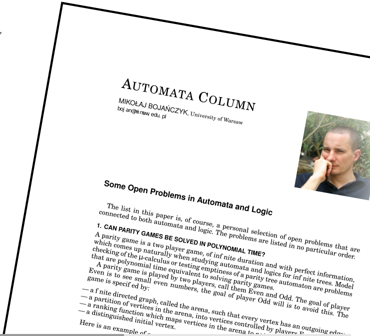
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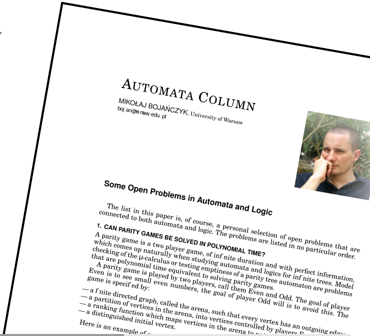
observational equivalence [Cotton-Barratt et al.'17]	data logics [Bojańczyk et al.'09, Abriola et al.'17]	security protocols [Verma & Goubault-Larrecq'05]	

Branching VAS Reachability

- important open problem [Bojańczyk'14]
- incorrect decidability proof in [Bimbó'15]
- application domains: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Some Open Problems in Automata and Logic

The list in this paper is, of course, a personal selection of open problems that are connected to both automata and logic. The problems are listed in no particular order.

1. CAN PARITY GAMES BE SOLVED IN POLYNOMIAL TIME?

A parity game is a two player game, of infinite duration and with perfect information, which comes up naturally when studying automata and logics for infinite trees. Model checking of the μ-calculus or testing emptiness of a parity tree automaton are problems that are polynomial time equivalent to solving parity games.

A parity game is played by two players, call them Even and Odd. The goal of player Even is to see small even numbers, the goal of player Odd will is to avoid this. The game is specified by:

— a finite directed graph, called the arena, such that every vertex has an outgoing edge

— a partition of vertices in the arena, into vertices controlled by players

— a ranking function which maps vertices in the arena to

— a distinguished initial vertex.

Here is an example of a

Complexity Theory	Distributed Computing	Proof Theory	Computational Biology
Tower-hard [Lazić & S., ToCL'15]	recursive parallel programs [Bouajjani & Emmi'13]	linear and relevance logics [de Groote et al.'04 Lazić & S., ToCL'15 S., JSL'16]	population protocols [Bertrand et al.'17]
Programming Languages	**Database Theory**	**Security**	**Computational Linguistics**
observational equivalence [Cotton-Barratt et al.'17]	data logics [Bojańczyk et al.'09, Abriola et al.'17]	security protocols [Verma & Goubault-Larrecq'05]	dominance grammars [Rambow'94; S., ACL'10] minimalist syntax [Salvati'10]

Summary

- well-quasi-orders
  ubiquitous in termination proofs

- complexity toolbox
  upper & lower bounds, fast-growing complexity classes

- application
  VAS reachability

Perspectives

1. complexity gap
   for VAS reachability

2. parameterisations
   for counter systems

3. beyond wqos
   FAC orders, Noetherian spaces

4. reachability in VAS extensions

Thank you!