# Automates d'arbre

# TD $n^{\circ}3$ : Relations

# October 10th, 2019

#### Exercise 1: Closure properties - back to basics

- 1. Given a recognizable relation, show that all its cylindrifications and projections are recognizable. (provides explicit trees automatas)
- 2. Is the domain and the image of a binary relation recognizable?
- 3. Given R, R' binary relations, show that  $R \circ R'$  is recognizable.
- 4. Give an example of a n-ary relation such that its ith projection followed by its ith cylindrification does not give back the original relation.
- 5. On the contrary, show that ith cylindrification followed by ith projection gives back the original relation.

## Exercise 2: Some relations

- 1. Let  $\mathcal{F} = \{0(2), 1(2), n(0)\}$ . Give an automaton recognizing :  $R_1 = \{(t, t') | t, t' \in T(\mathcal{F}), Pos(t) = Pos(t') \land \forall p \in Pos(t), t(p) = 1 \Rightarrow t'(p) = 1\}$
- 2. Let  $\mathcal{F} = \{f(2), g(1), a(0)\}$ . Is the relation  $R_2 = \{(g(t), t) \mid t \in T(\mathcal{F})\}$  recognizable? And if  $\mathcal{F} = \{g(1), a(0)\}$ ?
- 3. Here assume that  $\mathcal{F} = \{g(1), a(0)\}$ . Is  $R_2^*$  recognizable?
- 4. Is  $R_3 = \{(t, f(t, t')) \mid t, t' \in T(\mathcal{F})\}$  recognizable?
- 5. Design two relations, one recognizable and one which is not. Challenge your classmates with them.

## Exercise 3: Rewriting systems

Let  $\mathcal{F} = \{a_i(1) \mid 1 \le i \le n\} \cup \{0(0)\}.$ 

- 1. Prove that any rewrite system  $\rightarrow$  (i.e. the one step rewriting relation) on  $\mathcal{F}$  is recognizable.
- 2. Prove that  $S = \{(a_1^k(a_1(a_2(a_2^l(0)))), a_1^k(a_2^l(0))) \mid k, p \in \mathbb{N}\}$  is recognizable.
- 3. Prove that  $S^*$  is not recognizable.