Advanced Complexity

TD $n^{\circ}7$

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1 The homework is back!

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Exercise 1: A many to one reduction
Assuming that BestClique is NP-hard, show that BestClique is \nablaNP-complete.
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- **Exercise 2: Another many to one** Assuming $NP \neq coNP$, show that $BestClique \notin NP$.
- Exercise 3: Not a many to one
 - Assuming $NP \neq coNP$, show that BestClique $\notin coNP$.

2 The truth about the homework

Definition 1 A language L is in DP if and only if there are two language $L_1 \in NP$ and $L_2 \in coNP$ such that $L = L_1 \cap L_2$

The class DP (the D stands for differential) was introduced by Papadimitriou and Yannakis in [2]. Some problems were then proved to be complete for this class in [3], and the strictness properties where studied in [1].

Definition 2 The TSP problem is defined as :

- INPUT : A distance matrix and an integer B.
- QUESTION : Is there a tour of length B or less?
- The ExactTSP problem is defined as :
- INPUT : A distance matrix and an integer B.
- QUESTION : Is the length of the shortest tour B ?

Exercise 4: Money Money Money

- 1. Show that TSP is NP-complete. *Hint* : HAMILTONPATH *is* NP-*complete*
- 2. Why it is unlikely that $\mathsf{ExactTSP} \in \mathsf{NP}$ or $\mathsf{ExactTSP} \in \mathsf{coNP}$?
- 3. Prove that $\mathsf{ExactTSP} \in \mathsf{DP}$.
- 4. Modify the reduction from 3SAT to HAMILTONPATH, so that , wheter or not the expressions are satisfiable, the produced graph will always contain a broken hamiltonian path, that is, two node-disjoint paths that cover all nodes.
 You do not need to know the reduction of 3SAT to HAMILTONPATH, you just need to

You do not need to know the reduction of 3SAT to HAMILTONPATH, you just need to know that it produces a graph with a "starting" node and and an "ending" node, and if there is only one not satisfied clause, then there is a broken hamiltonian path in the graph.

5. Prove that ExactTSP is DP-complete.

Références

- E.W. Leggett and Daniel J. Moore. Optimization problems and the polynomial hierarchy. Theoretical Computer Science, 15(3):279 - 289, 1981.
- [2] C.H. Papadimitriou and M. Yannakakis. The complexity of facets (and some facets of complexity). Journal of Computer and System Sciences, 28(2):244-259, 1984.
- [3] Christos H. Papadimitriou and David Wolfe. The complexity of facets resolved. Journal of Computer and System Sciences, 37(1):2 13, 1988.