

Complexité avancée - TD 11

Simon Halfon

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Exercise 1: how big is PP?

Show that:

1. PP has complete problems.
2. $BPP \subseteq PP$
3. $NP \subseteq PP$
4. $MA \subseteq PP$

Exercise 2: Polynomial Identity Testing

From previous exercise sheet.

Exercise 3: Counting versus Deciding

Define FP to be the class of functions $\Sigma^* \rightarrow \Sigma^*$ that are computable by a Turing machine working in polynomial time. Remember the definition of #P and PP from this year's homework assignment.

Prove that $PP = P$ iff $\#P = FP$.

Exercise 4: how big is P/poly?

1. Show that $P/poly = P^{SPARSE}$, where SPARSE is the class of all sparse languages. A language L is said to be sparse if there is a polynomial p such that for all $n \in \mathbb{N}$, $|L \cap \Sigma^n| \leq p(n)$.
Hint: show $P/poly \subseteq P^{SPARSE} \subseteq P^P/poly \subseteq P/poly$.
2. Show that $P/poly = P^{UNARY}$, where UNARY is the class of all unary languages.
3. Deduce that $EXPTIME = EXPSPACE$ iff $PSPACE \cap P/poly = P$.
4. Deduce that if $P \neq RP$, then $EXPTIME \neq EXPSPACE$.

Exercise 5: P versus RP

Define a random language A as a language such that for every word $w \in \Sigma^*$, $w \in A$ with probability $1/2$.

Show that with probability 1 over the choice of a random A , $P^A = RP^A$.

Hint: use the random language as a random bit generator.