A few things on Noetherian spaces

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Noetherian spaces

- * **Defn.** A space is **Noetherian** iff every open is compact.
- * Here compact does not entail any kind of separation.
- Fact. The following are equivalent:
 (1) X is Noetherian
 - (2) Every subspace of *X* is compact
 - (3) Ascending sequences $U_1 \subseteq \overline{U}_2 \subseteq ... \subseteq U_n \subseteq ...$ of opens stabilize
 - (4) Descending sequences $C_1 \supseteq C_2 \supseteq ... \supseteq C_n \supseteq ...$ of closed sets stabilize

* We shall see other characterizations later.

Outline

- Characterizations of Noetherian spaces (half of them well-known)
- Transfering results from wqo theory to topology
- Applications in software verification
- Representations
- Conclusion

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Noetherian spaces, classically

- Defn. A space is Noetherian iff every open is compact.
- Prop. The spectrum of a Noetherian ring is a Noetherian space.
- * E.g., the spectrum of a polynomial ring over Q, R, or C.
 Not my first source of inspiration here.
 We shall see many (simpler) examples.
- Note. Noetherian + Hausdorff ⇔ finite, so we shall definitely drop Hausdorffness.

Noetherian spaces Hausdorff spaces

spaces only

Finite

Well-quasi-orders

- * Fact. The following are equivalent for a quasi-ordering ≤:
 (1) Every sequence (x_n)_{n∈ℕ} is good: x_m ≤ x_n for some m<n
 (2) Every sequence (x_n)_{n∈ℕ} is perfect: has a monotone subsequence
 (3) ≤ is well-founded and has no infinite antichain.
- Defn. Such a quasi-ordering ≤ is called a well-quasi-order (wqo).



 Applications: classification of graphs (Kuratowski, Robertson-Seymour) verification (computer science) model theory (logic: Fraïssé, Jullien, Pouzet)

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- * Given a qo (X, ≤), its **Alexandroff topology** has as opens U all upwards-closed subsets of X.
- * **Prop.** Let (X, ≤) be wqo. With its Alexandroff topology, X is Noetherian.
- * *Proof.* Consider an infinite ascending sequence $U_1 \subsetneq U_2 \subsetneq ... \subsetneq U_n \subsetneq ...$ of opens. Pick x_n in U_n , not in any previous U_m . By wqo, $x_m \le x_n$ for some m < n. Since $x_m \in U_m$ upwards-closed, $x_n \in U_m$: contradiction.

* Plenty of wqos ⇒ plenty of Noetherian spaces.

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 U_1

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Noetherian + Alexandroff

- * **Prop.** Let (X, \leq) be wqo. With its Alexandroff topology, X is Noetherian.
- There are also Noetherian spaces that are not Alexandroff:
 spectra of rings, with the Zariski topology
 powersets (see later)
- Conversely, the qo sets (X, ≤) that are Noetherian in their Alexandroff topology are exactly the wqo sets.
- * *Proof.* From $(x_n)_{n \in \mathbb{N}}$ define $U_n = \uparrow \{x_1, ..., x_n\}$. This stabilizes at $n: U_{n-1} = U_n$, so $x_n \in \uparrow \{x_1, ..., x_{n-1}\}$.

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Basic constructions

- Prop. (1) Every wqo is Noetherian in its Alexandroff topology
 (2) The spectrum of a Noetherian ring is Noetherian
 (3) Finite products of Noetherian spaces are Noetherian
 (4) Finite coproducts of Noetherian spaces are Noetherian
 (5) Subspaces of Noetherian spaces are Noetherian
 (6) Topologies coarser than a Noetherian topology are Noetherian
 (7) Continuous images of Noetherian spaces are Noetherian
 (6) are noetherian spaces are Noetherian
- * We shall see other constructions that preserve Noetherianness.
- * We need additional characterizations of Noetherianness.

Cluster points

• **Prop.** *X* is Noetherian iff every net $(x_i)_{i \in I}$ contains a cluster point x_i . (The important point is: the cluster point x_i belongs to the net.)

* Proof.
(⇒) If X Noetherian, then subspace K={x_i | i∈I} is compact, hence (x_i)_{i∈I} has a cluster point in K.
(⇐) Let U be open in X. Every net (x_i)_{i∈I} inside U has a cluster point in U, viz. some x_i. So U is compact. □

Note: in Alexandroff spaces, x_i cluster point means that for some *i*, infinitely many entries x_j are above x_i. (Take the open \\$x_i.)
 ... hence all sequences are good.

Self-convergent nets

* A net $(x_i)_{i \in I}$ is **self-convergent** iff it converges to **every** x_i . (A very much non-Hausdorff notion!)

◆ **Thm.** *X* is Noetherian iff every net $(x_i)_{i \in I}$ has a self-convergent subnet.

* *Proof.* (⇒) Let *J* be {*i*∈*I* | *x_i* is a cluster point of the net}. By previous Prop., *J* is non-empty. Check: *J* is cofinal and directed in *I*; so (*x_j*)_{*j*∈*J*} is a subnet. By Kelley's Theorem, (*x_j*)_{*j*∈*J*} has a further subnet that is an ultranet. Check that this ultranet is self-convergent.
(⇐) Obvious, using previous Prop. □

In Alexandroff spaces, (x_i)_{i∈I} self-convergent iff eventually monotone
 hence all sequences are perfect.

Ultrafilters

A similar characterization (lim 𝒴 = set of limits of 𝒴):

* **Thm.** X is Noetherian iff every ultrafilter \mathcal{U} is compact: $\lim \mathcal{U} \in \mathcal{U}$.

* Proof. (⇒) Let U be (open) complement of lim U.
If lim U not in U, then U is in U (ultrafilter).
Since U is compact, U has a limit in U.
So lim U intersects U: contradiction.
(⇐) Fix an open U. Let U be an arbitrary ultrafilter containing U.
Since lim U ∈ U, lim U ∩ U ∈ U, so lim U ∩ U≠Ø.
Hence U has a limit in U: U is compact. □

Application: finite products

- Well-known: finite products of Noetherian spaces are Noetherian. Here is a simple proof. (Warning: I'm lying a bit about what a subnet is.)
- Let X, Y be Noetherian.
 Let (x_i, y_i)_{i∈I} be a net in X×Y.
 Extract a self-convergent subnet (x_j)_{j∈J}.
 From (y_j)_{j∈J} extract a further self-convergent subnet (y_k)_{k∈K}.
 Then (x_k, y_k)_{k∈K} is a self-convergent subnet of the original net.
- This is a topological version of the Ramsey argument behind the classical wqo proofs.

Stone duality in a nutshell

- * There is a functor $\mathbf{O} : \mathbf{Top} \to \mathbf{Frame}^{op}$ that: — maps each space X to its frame $\mathbf{O}X$ of opens — maps $f:X \to Y$ to $\mathbf{O}f : \mathbf{O}Y \to \mathbf{O}X : V \mapsto f^1(V)$.
- * **O** is left-adjoint to a functor **pt** : **Frame**^{op} \rightarrow **Top**.
- * **S**=**pt O** is the **sobrification** functor.
- Defn. A space X is sober iff it is of the form pt L for some frame L iff it is of the form SY for some space Y iff X=SX (all that, up to iso.)

The specialization quasi-ordering

- * **Defn** (specialization, \leq). In a space *X*, $x \leq y$ iff every open containing *x* also contains *y* iff $x \in cl(\{y\})$.
- * *X* is T_0 iff \leq is antisymmetric (an ordering).
- Every open *U* is upwards-closed.
 Every closed set *C* is downwards-closed.
- * The closure $cl({x})$ is $\downarrow x = \{z \mid z \le x\}$.



Sober spaces

- Call *C* irreducible closed iff closed and:
 if C ⊆ ∪_{i=1}ⁿ C_i then C ⊆ C_i for some *i*.
 E.g., ↓x = cl({x}) is irreducible closed, for every point *x*.
- Fact. X is sober iff T₀ and all irreducible closed sets are of this form.
- * All Hausdorff spaces are sober, but there are more (e.g., continuous and quasi-continuous dcpos in domain theory).

Sobrification

- * The sobrification functor can be described more concisely as:

 - For $f: X \to Y$ to, $\mathbf{S}f : \mathbf{S}X \to \mathbf{S}Y : C \mapsto \operatorname{cl}(f(C))$.
 - $-X \text{ embeds into } \mathbf{S}X \text{ through } \eta : X \rightarrow \mathbf{S}X : x \mapsto \downarrow x.$
- Fact. X is Noetherian iff SX is Noetherian.
- * *Proof.* \diamond : **O***X* \rightarrow **OS***X* iso, and Noetherianness is a property of opens (ascending sequences of opens stabilize). \Box
- Fact. The Noetherian sober spaces X are the Stone duals pt L of distributive lattices L with the ascending chain condition.

Sober Noetherian spaces

- * An **order-theoretic** characterization. Call sets of the form $\downarrow \{x_1, ..., x_n\}$ **finitary**.
- **Thm.** A sober space X is Noetherian iff: (1) \leq is **well-founded**, and (2) the set of lower bounds of any finite set is **finitary**.
 - Then:
 - (3) The topology of *X* is the upper topology of \leq
 - (4) The closed sets are the finitary sets.
- *Proof.* Folklore, or see (JGL 2013) or (Dickmann,Schwartz&Tressl).



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Higman's Lemma



- * Lemma (Higman 1954). If *X*, ≤ is wqo, then so is the qoset *X*^{*} of finite words under **subword** relation \leq *.
- Thm (Topological Higman Lemma, JGL 2013). If X is Noetherian, then so is X* with the subword topology.
- My aim here is to give you a proof of that, imitating Nash-Williams' classic proof (1963) of Higman's Lemma.

The subword topology

*** Defn.** For opens *U*₁, *U*₂, ..., *U_n* in *X*, let [*U*₁ *U*₂... *U_n*] be set of words

 ... *a*₁ ... *a*₂ ... (etc.) ... *a_n* ...
 with *a*₁ ∈ *U*₁, *a*₂ ∈ *U*₂, ..., *a_n* ∈ *U_n*.

 The **subword topology** on *X*^{*} is generated by those sets [*U*₁ *U*₂... *U_n*].

* Specialization quasi-ordering is ≤*

 If X is Alexandroff, then subword topology=Alexandroff on X*, so Higman's Lemma is a special case of the topological Higman Lemma.

Bad sequences

- * Let *X* be a topological space, with subbase \mathscr{B} . A **bad sequence** is a sequence $(U_n)_{n \in \mathbb{N}}$ of elements of \mathscr{B} such that no U_n is included in a union $\bigcup_{m < n} U_n$ of previous elements.
- Lemma. If X is not Noetherian, then (whatever the subbase) it has a bad sequence.
- * *Proof.* Let *U* be non-compact open. By Alexander's Subbase Lemma, *U* has a cover (*U_i*)_{*i*∈*I*} by elements of *C* that has no finite subcover. Pick some *U_{i1}*. Some point *x*₁ in *U* is not in *U_{i1}*. Pick *U_{i2}* containing *x*₁. Some point *x*₂ in *U* is not in *U_{i1}* ∪ *U_{i2}*. Pick *U_{i3}* containing *x*₂. Some point ... etc. □

Minimal bad sequences

- * Assume additionally a well-founded \sqsubseteq ordering on \mathscr{B} . A **bad sequence** $(U_n)_{n \in \mathbb{N}}$ is **minimal** iff every \sqsubseteq -lexicographically smaller sequence $(V_n)_{n \in \mathbb{N}}$ is **good** (i.e., not bad). [i.e., $V_0 = U_0$, $V_1 = U_1$, ..., $V_{n-1} = U_{n-1}$, and $V_n \sqsubset U_n$ (strictly) for some n]
- Lemma. If X is not Noetherian, then (whatever *B* and ⊆) it has a minimal bad sequence.
- * *Proof.* Find $U_0 \sqsubseteq$ -minimal so that it starts a bad sequence. Given U_0 , find $U_1 \sqsubseteq$ -minimal so that U_0 , U_1 start a bad sequence. Given U_0 , U_1 , find $U_2 \sqsubseteq$ -minimal so that ... etc. \Box
- * Note: Similar to wqos, where bad sequences are sequences of **points**.

Proof plan

- * On X*, let *C* consist of the subbasic opens [U₁ U₂... U_m]. Let [U₁ U₂... U_m] ⊆ [V₁ V₂... V_n] iff there is a (strictly) increasing map f:{1,2,...,m} → {1,2,...,n} such that U_k = V_{f(k)}. (I.e., the V_ps are obtained by inserting new opens in the list of U_ks.)
- * If *X** is not Noetherian, then extract some minimal bad sequence.
- Using the zoom-in Lemma (next slide), find a smaller sequence: that one must be good, leading to a contradiction.

The zoom-in Lemma

- * **Lemma.** Let *X* be Noetherian, and $a_n \in U_n$ open for each $n \in \mathbb{N}$. There is a subsequence $(a_{n(k)})_k$ s.t. $a_{n(k)} \in U_{n(0)} \cap ... \cap U_{n(k)}$ for every *k*.
- * *Proof.* Pick a cluster point $a_{n(0)}$ (inside the sequence itself).

Infinitely many $a_n s$ with n > n(0) are in $U_{n(0)}$, forming a subsequence. Pick a cluster point $a_{n(1)}$ from that subsequence.

Infinitely many $a_n s$ with n > n(1) from that subsequence are in $U_{n(0)} \cap U_{n(1)}$, forming a sub-subsequence. Pick a cluster point $a_{n(2)}$ from that sub-subsequence... etc.

 Thm (Topological Higman Lemma, JGL 2013). If X is Noetherian, then so is X* with the subword topology.

* Proof (1/3).
Imagine X* is not Noetherian, and let
𝒰_n = [U_{n1} U_{n2} ... U_{nm} ...] form a minimal bad sequence.
Pick a word 𝔹_n in 𝔅_n that is in no previous 𝔅_m.
Let 𝔅_n = [U_{n2} ... U_{nm} ...] be « 𝔅_n without its first open U_{n1}».

By definition, $w_n = l_n a_n r_n$ where $a_n \in U_{n1}, r_n \in \mathcal{R}_n$. By zoom-in, extract $(a_{n(k)})_k$ s.t. $a_{n(k)} \in U_{n(0)1} \cap ... \cap U_{n(k)1}$ for every k. By minimality, $\mathcal{U}_0, \mathcal{U}_1, ..., \mathcal{U}_{n(0)-1}, \mathcal{R}_{n(0)}, \mathcal{R}_{n(1)}, ...$ is **good**. So, for some k: $\mathcal{R}_{n(k)} \subseteq \mathcal{U}_0 \cup ... \cup \mathcal{U}_{n(0)-1} \cup \mathcal{R}_{n(0)} \cup ... \cup \mathcal{R}_{n(k-1)}$.

- Thm (Topological Higman Lemma, JGL 2013). If X is Noetherian, then so is X* with the subword topology.
- * *Proof* (2/3). Recall: $\mathcal{U}_n = [U_{n1} \ U_{n2} \dots \ U_{nm} \dots], w_n \text{ in } \mathcal{U}_n, \text{ in no previous } \mathcal{U}_m.$ $\mathcal{R}_n = [U_{n2} \dots \ U_{nm} \dots]$ $w_n = l_n \ a_n \ r_n \text{ where } a_n \in U_{n1}, \ r_n \in \mathcal{R}_n.$ $a_{n(k)} \in U_{n(0)1} \cap \dots \cap U_{n(k)1}$ $\mathcal{R}_{n(k)} \subseteq \mathcal{U}_0 \cup \dots \cup \mathcal{U}_{n(0)-1} \cup \mathcal{R}_{n(0)} \cup \dots \cup \mathcal{R}_{n(k-1)}.$

* Note $r_{n(k)} \in \mathcal{R}_{n(k)}$.

- Thm (Topological Higman Lemma, JGL 2013). If X is Noetherian, then so is X* with the subword topology.
- * Proof (2/3). Recall: $\mathcal{U}_n = [\mathcal{U}_{n1} \ \mathcal{U}_{n2} \dots \mathcal{U}_{nm} \dots], w_n \text{ in } \mathcal{U}_n, \text{ in no previous } \mathcal{U}_m.$ $\mathcal{R}_n = [\mathcal{U}_{n2} \dots \mathcal{U}_{nm} \dots]$ $w_n = l_n \ a_n \ r_n \text{ where } a_n \in \mathcal{U}_{n1}, \ r_n \in \mathcal{R}_n.$ $a_{n(k)} \in \mathcal{U}_{n(0)1} \cap \dots \cap \mathcal{U}_{n(k)1}$ $\mathcal{R}_{n(k)} \subseteq \mathcal{U}_0 \cup \dots \cup \mathcal{U}_{n(0)-1} \cup \mathcal{R}_{n(0)} \cup \dots \cup \mathcal{R}_{n(k-1)}.$ * Note $r_{n(k)} \in \mathcal{R}_{n(k)}.$ Case 1: $r_{n(k)} \in \mathcal{U}_m$ for some m in 0, ..., n(0)-1. Then the larger $w_{n(k)}$ is in \mathcal{U}_m , too. (Opens are upwards-closed.) Impossible since \mathcal{U}_m is previous (m < n(k)).

- Thm (Topological Higman Lemma, JGL 2013). If X is Noetherian, then so is X* with the subword topology.
- * Proof (3/3). Recall: $\mathcal{U}_n = [\mathcal{U}_{n1} \ \mathcal{U}_{n2} \dots \mathcal{U}_{nm} \dots], w_n \text{ in } \mathcal{U}_n, \text{ in no previous } \mathcal{U}_m.$ $\mathcal{R}_n = [\mathcal{U}_{n2} \dots \mathcal{U}_{nm} \dots]$ $w_n = l_n \ a_n \ r_n \text{ where } a_n \in \mathcal{U}_{n1}, \ r_n \in \mathcal{R}_n.$ $a_{n(k)} \in \mathcal{U}_{n(0)1} \cap \dots \cap \mathcal{U}_{n(k)1}$ $\mathcal{R}_{n(k)} \subseteq \mathcal{U}_0 \cup \dots \cup \mathcal{U}_{n(0)-1} \cup \mathcal{R}_{n(0)} \cup \dots \cup \mathcal{R}_{n(k-1)}.$ * Note $r_{n(k)} \in \mathcal{R}_{n(k)}.$ Case 2: $r_{n(k)} \in \mathcal{R}_{n(j)}$ for some j in 0, ..., k-1. Note that $a_{n(k)} \in \mathcal{U}_{n(j)1}.$ Hence $w_{n(k)} = l_{n(k)} \ a_{n(k)} \ r_{n(k)}$ is in $\mathcal{U}_{n(j)}$, too. Impossible since $\mathcal{U}_{n(j)}$ is previous (j < k).

Kruskal's Theorem



- **Thm** (Kruskal 1960). If *X*, ≤ is wqo, then so is the qoset $\mathcal{T}(X)$ of finite trees labeled by *X* under **homeomorphic embedding** relation \leq_{\leq} .
- * **Thm** (Topological Kruskal Theorem, JGL 2013). If X is Noetherian, then so is $\mathcal{T}(X)$ with the **tree** topology.
- * Admitted. Slightly more complex.

Powersets

- * Let $\mathbb{P}(X)$ come with the lower Vietoris topology, with subbase $U = \{A \mid A \cap U \neq \emptyset\}, U \in \mathbf{O}X$.
- * **Thm** (JGL, 2007). If X is Noetherian, then so is $\mathbb{P}(X)$.
- * *Proof.* If $\mathbb{P}(X)$ not Noetherian, let $(\diamond U_n)_{n \in \mathbb{N}}$ be a bad sequence: no $\diamond U_n$ is included in $\bigcup_{m < n} \diamond U_m = \diamond \bigcup_{m < n} U_m$.

Since \diamond is monotonic, no U_n is included in $\bigcup_{m < n} U_m$. Therefore $(U_n)_{n \in \mathbb{N}}$ is bad: contradiction. \Box

* Specialization qo: $A \leq {}^{\flat}B$ iff every $a \in A$ is below some $b \in B$.

Powersets, or: beyond wqos

- * Let $\mathbb{P}(X)$ come with the lower Vietoris topology, with subbase $U = \{A \mid A \cap U \neq \emptyset\}, U \in \mathbf{O}X$.
- * **Thm** (JGL, 2007). If X is Noetherian, then so is $\mathbb{P}(X)$.
- * Specialization qo: $A \leq {}^{\flat}B$ iff every $a \in A$ is below some $b \in B$. Pretty remarkable, since:
- **Prop** (Rado, 1957). There are wqos X, ≤ such that $\mathbb{P}(X)$, ≤ ^b is **not** wqo.
 - $X_{\text{Rado}} = \{(i,j) \mid i \leq j\}, (i,j) \sqsubseteq (k,l) \text{ iff } i = k \text{ and } j \leq l \\ \text{or } j \leq k.$



A catalogue of Noetherian spaces

D ::= A \mathbb{N} $D_1 \times D_2 \times \ldots \times D_n$ $D_1 + D_2 + \ldots + D_n$ $\mathcal{S}(D)$ $\mathbb{P}(D)$ $\mathbb{P}^*(D)$ $\mathcal{H}(D)$ $\mathcal{H}_{\emptyset}(D)$ $\operatorname{Spec}(R)$ \mathbb{C}^k D^* D^* $\triangleright_{n=1}^{+\infty} D_n$ $\mathcal{T}(D)$

finite poset natural numbers products sums sobrification powerset non-empty powerset extended Hoare powerspace Hoare powerspace spectrum of a ring complex vector space (Zariski) words / embedding multisets words / prefix trees / embedding

A catalogue of Noetherian spaces

| D ::= | $A \\ \mathbb{N} \\ D_1 \times D_2 \times \ldots \times D_n \\ D_1 + D_2 + \ldots + D_n$ | finite poset natural numbers products sums | Noetherian, |
|-------|--|---|-------------|
| | $\mathcal{S}(D)$ $\mathbb{P}(D)$ | sobrification | not wqo |
| | $\mathbb{P}^*(D)$ | | - |
| | $\mathcal{H}(D)$ | | |
| | $\mathcal{H}_{\emptyset}(D)$ | Hoare powerspace | |
| | $\operatorname{Spec}(R)$ | spectrum of a ring | |
| | \mathbb{C}^k | complex vector space (Zariski) | |
| | D^* | words / embedding | |
| | D^{\circledast} | multisets | |
| | $\triangleright_{n=1}^{+\infty} D_n$ | words / prefix | |
| | $\mathcal{T}(D)$ | trees / embedding | |

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Verification

- * How do you ensure a software/hardware system Sys is correct?
- * **Testing**: fine and useful, but not exhaustive
- * Verification: given a desirable property P, check that
 Sys ⊧ P
- * That check should be done by an **algorithm**.



Verification

- A paradigmatic case is given by:
 Sys is a transition system

 (a directed graph, vertices=states)

 P is a (non-)reachability property

 «can Sys evolve from an initial state s to a state in the set Bad?»
- Verification is undecidable in general.
 Decidable for Sys finite.
 But most systems are infinite.
- Classes of infinite-state systems/properties for which verification would be decidable?





Well-Structured Transition Systems (WSTS)

- * A transition system: state space *X*, transition relation $\xrightarrow{\delta}$
- ✤ with a wqo ≤
- satisfying monotonicity



(Finkel 1990, Abdulla,Čerāns, Jonsson&Tsay 2000, Finkel&Schnoebelen 2001)



Lossy channel systems





... and many other examples

Topological WSTS

- * A transition system: state space *X*, transition relation $\xrightarrow{\delta}$
- * with a **Noetherian** topology
- * satisfying **lower semi-continuity**: for every open U, $\delta^{-1}(U)$ open.

| Potrinote NADPH+H+ ATP | COo | | |
|---|---|--|--|
| while (*) { x = x * y - 6; y = $if (x^2 - 3 * x * y)$ while (*) x = x * y - 6; y = $if (x^2 - 3 * x * y)$ while (*) else send (S) $x = x^2 + x * y;$ $ recv (SIG_QUIT) \Rightarrow rockson y$ | $\begin{array}{c} \text{polynomial } p \\ \text{f (*) } \{ x = 2; y = 3; \} \\ \text{lse } \{ x = 3; y = 2; \} \\ = 0; \\ y == 0 \\) \{ x = x + 1; y = y - 1; \}; \\ \text{SIG_FUZZ}; \\ \leftarrow \qquad \qquad$ | $\begin{array}{c} a = *; b = \\ while (*) \\ \hline \\ a = c \\ \hline \\ c \\ c$ | |

The standard backward algorithm

- * **Defn** (coverability). **INPUT**: state *s*, and open state set **Bad QUESTION**: $s \rightarrow_{\delta}$ * **Bad**?
- Prop. Given an effective topological WSTS, coverability is decidable.
- * *Proof.* The function pre* computes $U_0 = \text{Bad}, U_{n+1} = U_n \cup \delta^{-1}(U_n).$ This **terminates** because $U_0 \subseteq U_1 \subseteq ... \subseteq U_n \subseteq ... \text{ stabilizes}$ (Noetherianness). At the end, $U_n = \{s \mid s \rightarrow_{\delta} * \text{Bad}\}.$

```
fun pre* U =
   let V = pre U
   in
        if V⊆U
        then U
        else pre* (U ∪ V)
   end;
fun coverability (s, bad) =
        s in pre* (bad);
```

(Don't be fooled by the simplicity of the algorithm: complexity is not even primitive recursive in general —and I mean the complexity of the problem, independently of the algorithm.)

Effective?

- * By an **effective** topological WSTS, we mean one where: — opens *U* are **representable** by some data structure — the inclusion test $U \subseteq V$ is **decidable** — one-step predecessors $\delta^{-1}(U)$ of open sets are **computable**.
- When the Noetherian state space is Alexandroff (a wqo), there is a standard representation of open (upwards-closed) sets:
- **Prop.** In a wqo, every upwards-closed subset is the upward closure \$\$\{x_1, ..., x_n\$}\$ of finitely many points.
- * No longer true in more general Noetherian spaces.

Sobrification and closed sets

* Recall that, in a sober Noetherian space, every closed set *C* is **finitary**: $C = \bigcup \{x_1, ..., x_n\}.$



- Provides alternate representation of open sets U in Noetherian X:
 - *U* is also open in larger space SX
 - represent U by its closed complement in SX,
 - i.e., by finite sets of **points** in **S**X.
- ✤ ⇒ Find computable representations of points of SX.

Outline

- Characterizations of Noetherian spaces (half of them well-known)
- Transfering results from wqo theory to topology
- Applications in software verification
- Representations
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A simple case

- * Consider \mathbb{N} with the (Alexandroff topology of) its ordering \leq .
- * Its closed subsets are \emptyset , $[n] = \{0, 1, ..., n\}$ and the whole of \mathbb{N} .
- * All except \emptyset are irreducible. Hence:
- * **Prop.** A representation for **S** \mathbb{N} is \mathbb{N}_{ω} , i.e., \mathbb{N} plus a top element ω .
- * This is an effective representation.
- I'm not giving the topology on N_ω:
 this must be the upper topology of its ordering.

Products

- * Another simple case. We know that $\mathbf{S}(\prod_i X_i) = \prod_i \mathbf{S} X_i$, up to iso (R.-E. Hoffmann 1979). Hence:
- **Prop.** A representation of S(X₁ × ... × X_n) is the Cartesian product of representations for SX_i.
- This is effective,
 provided the representations for SX_i are.

Words and regular expressions

- * Thm (Finkel&JGL 2009). A representation for $S(X^*)$ is the space of word products, i.e., regular expressions of the form: $R_1 R_2 \dots R_n$ where each R_i is of the form: $-(\downarrow a)^?$ with $a \in S(X)$, or $-(\downarrow \{a_1, ..., a_k\})^*$ with $a_1, ..., a_k \in S(X)$
- Proof omitted. Again, effectivity is preserved.
 Was already known for wqos (Kabil&Pouzet 1992).
- * Embedding $X^* \rightarrow \mathbf{S}(X^*)$ maps $a_1a_2...a_n$ to $(\downarrow a_1)^? (\downarrow a_2)^? ... (\downarrow a_n)^?$ Limit elements include $(\downarrow a_1)^? (\downarrow \{a_2, a_3\})^* (\downarrow a_1)^?$, for example.

 Thm (Finkel&JGL, unpublished).
 For all the spaces X in our catalogue of Noetherian spaces,
 SX has effective representations.

A catalogue of Noetherian spaces

| D ::= | $A \\ \mathbb{N} \\ D_1 \times D_2 \times \ldots \times D_n \\ D_1 + D_2 + \ldots + D_n$ | finite poset natural numbers products sums | Noetherian. |
|-------|---|--|-------------|
| | $\mathcal{S}(D)$ $\mathbb{P}(D)$ $\mathbb{P}^{*}(D)$ $\mathcal{H}(D)$ | sobrification powerset non-empty powerset extended Hoare powerspace | not wqo |
| | $R_{0}(D)$ Spec(R) \mathbb{C}^{k} | spectrum of a ring complex vector space | |
| | $ \begin{array}{c} D^{*} \\ D^{*} \\ \bigtriangledown_{n=1}^{+\infty} D_{n} \end{array} $ | words / embedding multisets words / prefix | |
| T | $\mathcal{T}(D)$ | trees / embedding | |

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...

- Thm (Finkel&JGL, unpublished).
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 SX has effective representations.
- ... up to a small change: replace
 Spec(R) and C^k (Zariski) with some concrete spectrum, say,
 Spec(Q[X₁,X₂,...,X_n]).

A catalogue of Noetherian spaces



- Thm (Finkel&JGL, unpublished).
 For all the spaces X in our catalogue of Noetherian spaces,
 SX has effective representations.
- ... up to a small change: replace Spec(R) and C^k (Zariski) with some concrete spectrum, say, Spec(Q[X₁,X₂,...,X_n]).
- Including (infinite) powersets!

A catalogue of Noetherian spaces



Representing powersets?

- * Recall $\mathbb{P}(X)$ has subbase $\diamond U = \{A \mid A \cap U \neq \emptyset\}, U \in \mathbf{O}X$. Let $\mathbb{F}(Y)$ be the set of **finite** subsets of *Y*, with subspace topology.
- **Prop.** S(P(X)) = F(SX), up to iso.
 Hence a representation for S(P(X)) is given by finite sets of elements from a representation of SX.
- * Proof. Let H(X) be subspace of closed subsets of X.
 cl : P(X) → H(X) is a quasi-iso, i.e. O(cl) : OH(X) → OP(X) is iso. (Exercise: use A ∩ U≠Ø iff cl(A) ∩ U≠Ø.)
 Hence S(P(X)) = S(H(X)).
 Since OX=OSX, H(X)=H(S(X)).
 Now recall that all elements of H(S(X)) are finitary.

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Conclusion

 There is more to Noetherian spaces than algebraic geometry. A deep connection with wqos. Basic theory in (JGL 2013), Section 9.7.

- * Any topological analogue of **better** quasi-orderings?
- Any topological analogue of the Robertson-Seymour theorem (for undirected finite graphs with labels in a Noetherian space)?
- Any topological analogue of the theory of maximal order types of wqos? (Hint: ordinal height of H(X), of SX.) Application to complexity of WSTS algorithms (à la Schnoebelen, Schmitz, Halfon).

Non-Hausdorff Topology and Domain Theory

Jean Goubault-Larrecq