Testing with asynchronous communication

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Joint work with Puneet Bhateja and Madhavan Mukund

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Outline

1. Introduction

Input/Output semantics

IO-Blocks semantics

Queue semantics (Tretman)

Conclusion

Introduction

Verification of software or hardware

- Proof
- Model checking
- Test

Synchronous testing

- The tester interacts synchronously with the system.
- The tester proposes an action which is either refused or accepted and executed by the system.
- The tester has an immediate feedback.

Asynchronous testing

- The tester communicate asynchronously with the system
- The tester provides inputs and observes outputs.
- The tester does not necessarily know whether its inputs have been used by the system or not.

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- Tests are computed in advance and are sent as a whole stream to the system
- The tester then observes the output streams generated by the system

on the fly test generation – IO-Blocks semantics

- Inputs are supplied incrementally.
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Test equivalence

- Equivalence of two systems for a given test semantics.
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Related work

- I. Castellani and M Hennessy: Testing Theories for Asynchronous Languages, *Proc. FSTTCS '98*, Springer Lecture Notes in Computer Science **1530** (1998) 90–101.
- R. de Nicola and M. Hennessy: Testing equivalences for processes, *Theoretical Computer Science*, **34** (1984) 83–133.
- A. Petrenko, N. Yevtushenko and J.L. Huo: Testing Transition Systems with Input and Output Testers, *Proc IFIP TC6/WG6.1 XV International Conference on Testing of Communicating Systems (TestCom 2003)*, Sophia Antipolis, France, (2003) 129–145.
- J. Tretmans: Test Generation with Inputs, Outputs and Repetitive Quiescence, *Software—Concepts and Tools*, **17**(3) (1996) 103–120.

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2 Input/Output semantics

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Queue semantics (Tretman)

Conclusion

The model

Labelled transition system

$TS = (S, \Sigma, I, T)$ where

- S is the set of states
- $I \subseteq S$ is the set of initial states
- $\Sigma = \Sigma_i \uplus \Sigma_o$ is the set of input/output actions
- $T \subseteq S \times \Sigma \times S$ is the set of transitions

$L(TS) = \{w \in \Sigma^* \mid I \xrightarrow{w} \text{ in } TS\}$.

$s \in S$ is quiescent if it refuses all output actions: $s \stackrel{\Sigma_o}{\not\rightarrow}$.

Some further properties

- No infinite output-only behaviour.
- Receptiveness: $\forall s \in S$ quiescent, $\forall a \in \Sigma_i$, $s \xrightarrow{a}$

  If this is not the case, we may

  - discard unexpected inputs
  - enter a dead state accepting all inputs and with no possible outputs.

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Asynchronous IO-Behaviours

Intuition: Provide some test input $u \in \Sigma_i^*$ up front and observe the maximal outcome $v \in \Sigma_o^*$. Corresponds to static test generation.

Definition: IO-Behaviours

Let $TS = (S, \Sigma, I, T)$. $IOBeh(TS)$ is the set of pairs $(u, v) \in \Sigma_i^* \times \Sigma_o^*$ such that there is a run $i \xrightarrow{w} s$ in TS with

- $i \in I$ and s quiescent
- $\pi_o(w) = v$, and
- either $\pi_i(w) = u$ or there exists $a \in \Sigma_i$ such that $\pi_i(w)a \preceq u$ and $s \stackrel{a}{\nrightarrow}$.

Example

$IOBeh(TS_1)$:
$(\varepsilon, \varepsilon)$
(a, x), (a, xy)
(a^2, x), (a^2, xy), (a^2, x^2)
(a^n, x), (a^n, xy), (a^n, x^2) if $n \geq 2$.

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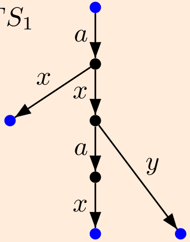
TS_1



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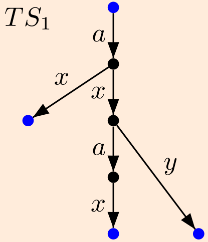

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Asynchronous testing equivalence (1)

IO-equivalence

Two transition systems TS and TS' are IO-equivalent, denoted $TS \sim_{io} TS'$ if

$$IOBeh(TS) = IOBeh(TS')$$

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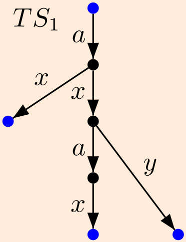

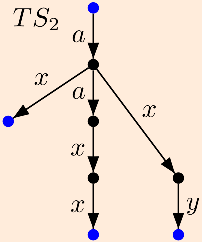

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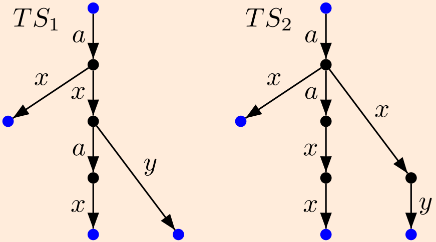

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Rational relations

Definition

Let A, B be two finite (and disjoint) alphabets.
A rational relation over A and B is a rational subset R of the monoid $A^* \times B^*$.

Equivalently, $R \subseteq A^* \times B^*$ is a rational relation if there exists an automaton $\mathcal{A} = (S, A \cup B, I, F, T)$ such that

$$R = \{(u, v) \in A^* \times B^* \mid \exists i \xrightarrow{w} f \text{ in } \mathcal{A} \text{ with } i \in I, f \in F, \pi_A(w) = u, \pi_B(w) = v\}$$

Example

$$\mathcal{R}(\mathcal{A}) = \{(a, x), (a, xy), (a^2, x^2)\}$$

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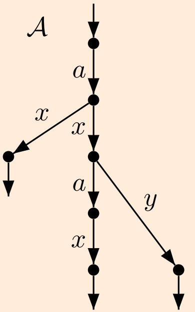

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From IO-behaviours to rational relations

Proposition

From a transition system $TS = (S, \Sigma, I, T)$, we can construct an automaton $\mathcal{A}$ over $\Sigma = \Sigma_i \uplus \Sigma_o$ such that

$$IOBeh(TS) = \mathcal{R}(\mathcal{A})$$

Proof. Intuition: transform quiescent states into final states

Let $D \subseteq S$ be the set of quiescent states of TS. Define $\mathcal{A} = (S', \Sigma, I', F', T')$

- $S' = S \uplus \overline{D} \uplus \{f\}$ where $\overline{D}$ is a copy of D.
- $I' = I \uplus \overline{I \cap D}$ and $F' = \overline{D} \uplus \{f\}$
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Let $(u, v) \in IOBeh(TS)$ and $i \xrightarrow{w} s$ in TS with $i \in I$, $s \in D$, $\pi_o(w) = v$ and $u = \pi_i(w)au'$ with $s \stackrel{a}{\nrightarrow}$.

Then, $i \xrightarrow{w} \bar{s} \xrightarrow{a} f \xrightarrow{u'} f$ in $\mathcal{A}$ and $u = \pi_i(wau')$, $w = \pi_o(wau')$.

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Decidability of IO-equivalence

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Corollary

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From rational relations to IO-behaviours

Several problems:

- Final states may not be quiescent (easy to fix).
- Quiescent states may not be final (harder to fix).

Example

Same rational relation: $\mathcal{R}(\mathcal{A}_1) = \{(a^2, x^3)\} = \mathcal{R}(\mathcal{A}_2)$

But different IO-behaviours:

$$\begin{aligned} IOBeh(\mathcal{A}_1) &= \{(\varepsilon,\varepsilon),(a,x^2)\} \cup \{(a^n,x^3) \mid n \geq 2\} \\ IOBeh(\mathcal{A}_2) &= \{(\varepsilon,\varepsilon),(a,x)\} \cup \{(a^n,x^3) \mid n \geq 2\} \end{aligned}$$

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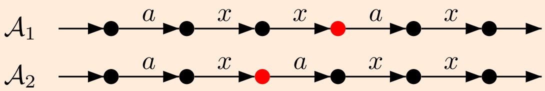

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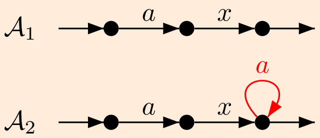

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$\mathrm{Rat}(B^*)$-automata

Definition

A $\mathrm{Rat}(B^*)$-automaton over A is a tuple $\mathcal{A} = (S, A, \lambda, \mu, \gamma)$ where

- S is the finite set of states
- $\lambda : S \to \mathrm{Rat}(B^*)$ $\qquad \xrightarrow{\lambda_s} (s)$

  A word in λ_s is emitted when entering $\mathcal{A}$ in state s.
- $\mu : A \to (S \times S \to \mathrm{Rat}(B^*))$ $\qquad (r) \xrightarrow{a \,/\, \mu(a)_{r,s}} (s)$

  A word in $\mu(a)_{r,s}$ is emitted when taking a transition from r to s labelled a.
- $\gamma : S \to \mathrm{Rat}(B^*)$ $\qquad (s) \xrightarrow{\gamma_s}$

  A word in γ_s is emitted when exiting $\mathcal{A}$ in state s.

Then, $(u, v) \in \mathcal{R}(\mathcal{A})$ if there is a path $P = s_0 \xrightarrow{a_1} s_1 \cdots s_{n-1} \xrightarrow{a_n} s_n$ in $\mathcal{A}$ with

- $u = a_1 \cdots a_n$
- $v \in \lambda_{s_0} \mu(a_1)_{s_0,s_1} \cdots \mu(a_n)_{s_{n-1},s_n} \gamma_{s_n}$.

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  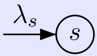

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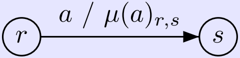
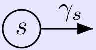
  

  A word in $\mu(a)_{r,s}$ is emitted when taking a transition from r to s labelled a.
- $\gamma : S \to \mathrm{Rat}(B^*)$

  A word in γ_s is emitted when exiting $\mathcal{A}$ in state s.

Then, $(u, v) \in \mathcal{R}(\mathcal{A})$ if there is a path $P = s_0 \xrightarrow{a_1} s_1 \cdots s_{n-1} \xrightarrow{a_n} s_n$ in $\mathcal{A}$ with

- $u = a_1 \cdots a_n$
- $v \in \lambda_{s_0} \mu(a_1)_{s_0,s_1} \cdots \mu(a_n)_{s_{n-1},s_n} \gamma_{s_n}$.

$\mathrm{Rat}(B^*)$-automata and rational relations

Theorem

A relation $R \subseteq A^* \times B^*$ is rational iff there exists a $\mathrm{Rat}(B^*)$-automaton $\mathcal{A}$ with $R = \mathcal{R}(\mathcal{A})$.

Theorem

If $|A| \geq 2$ then equivalence is undecidable for $\mathrm{Rat}(B^*)$-automata over A.
This holds even if

- $|B| = 1$
- We use only finite languages: $\mathcal{P}_{\mathsf{fin}}(B^*)$-automata
- There is no output when entering the automaton: $\lambda_s \neq \emptyset$ implies $\lambda_s = \{\varepsilon\}$
- There is no output when exiting the automaton: $\gamma_s \neq \emptyset$ implies $\gamma_s = \{\varepsilon\}$
- All transitions are visible: $\varepsilon \notin \mu(a)_{r,s}$

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Undecidability of IO-equivalence

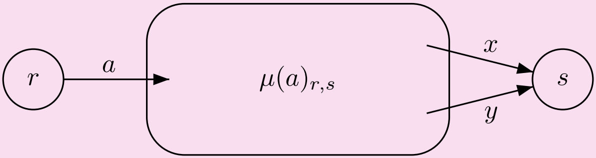
Theorem

IO-equivalence is undecidable if $|\Sigma_i| \geq 2$ and $|\Sigma_o| \geq 2$.

Proof

Let $\mathcal{A} = (S, A, \lambda, \mu, \gamma)$ be a $\mathcal{P}_{\mathsf{fin}}(B^+)$-automaton with $|A| = 2$ and $|B| = 1$.
Define $\mathcal{A}' = (S', \Sigma, I', T')$ by

- $\Sigma_i = A$, $\Sigma_o = B \uplus \{\#\}$ and $I' = \{s \in I \mid \lambda_s \neq \emptyset$ (i.e., $\lambda_s = \{\varepsilon\}$)$\}$
- transitions $r \xrightarrow{a \ / \ \mu(a)_{r,s}} s$ of $\mathcal{A}$ are replaced in $\mathcal{A}'$ by



Note that quiescent states of $\mathcal{A}'$ are exactly the states of $\mathcal{A}$.

Claim: $(u, v) \in IOBeh(\mathcal{A}')$ iff there is a path $s_0 \xrightarrow{a_1} s_1 \cdots s_{n-1} \xrightarrow{a_n} s_n$ in $\mathcal{A}$ with $\lambda_{s_0} = \{\varepsilon\}$, $v \in \mu(a_1)_{s_0,s_1} \cdots \mu(a_n)_{s_{n-1},s_n}$, and
$u = a_1 \cdots a_n$ or $u = a_1 \cdots a_n a u'$ with $\mu(a)_{s_n,s} = \emptyset$ for all $s \in S$.

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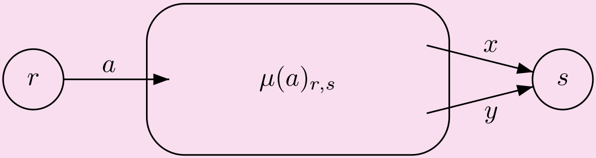

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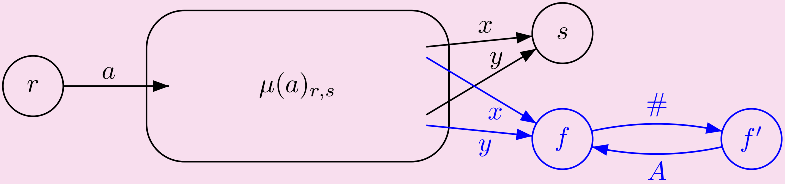

Note that quiescent states of $\mathcal{A}'$ are exactly the states in $S \uplus \{f'\}$.

Lemma $IOBeh(\mathcal{A}'') = IOBeh(\mathcal{A}') \cup \mathcal{R}(\mathcal{A}) \cdot \{(x, \#^{1+|x|}) \mid x \in A^*\}$.

Lemma $\mathcal{A}' \uplus \mathcal{B}'' \sim_{io} \mathcal{A}'' \uplus \mathcal{B}'$ if and only if $\mathcal{R}(\mathcal{A}) = \mathcal{R}(\mathcal{B})$.

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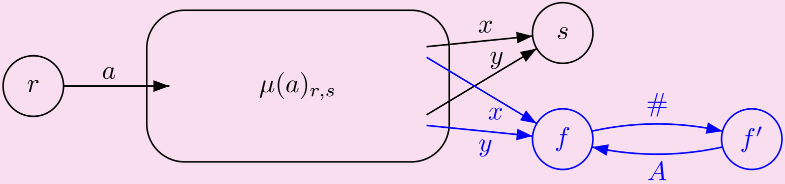

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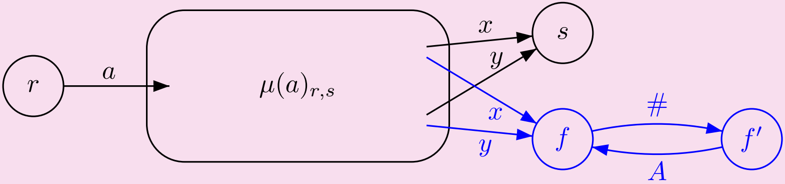

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Outline

Asynchronous IO-blocks semantics

Definition

A *block observation* of $TS = (S, \Sigma, I, T)$ is a sequence $(u_1, v_1) \cdots (u_n, v_n)$ where

- $u_1 \in \Sigma_i^*$ and $u_j \in \Sigma_i^+$ for $1 < j \leq n$,
- $v_k \in \Sigma_o^*$ for $1 \leq k \leq n$

and there is a run $s_0 \xrightarrow{w_1} s_1 \cdots \xrightarrow{w_k} s_k$ with $s_0 \in I$, $1 \leq k \leq n$ and:

- $s_1, s_2, \ldots, s_k$ are quiescent.
- $\pi_o(w_j) = v_j$ for $1 \leq j \leq k$ and $v_j = \varepsilon$ for $k < j \leq n$.
- $\pi_i(w_j) = u_j$ for $0 \leq j < k$.
- Either $k = n$ and $\pi_i(w_n) = u_n$ or there exists $a \in \Sigma_i$ with $\pi_i(w_k)a \preceq u_k$ and $s_k \stackrel{a}{\nrightarrow}$.

Let $IOBlocks(TS)$ denote the set of block observations of TS.

IO-block equivalence

IO-block equivalence

Two transition systems TS and TS' are IO-block equivalent if

$$IOBlocks(TS) = IOBlocks(TS')$$

This equivalence is denoted $TS \sim_{ioblock} TS'$.

Remark

IO-block equivalence corresponds to the queued suspension trace equivalence of

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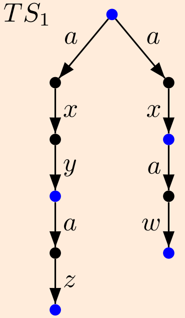

$IOBlocks(TS_1)$:
$(\varepsilon, \varepsilon)$
(a, xy)
(a, x)
$(a, xy)(a^n, z)$ for $n \geq 1$
$(a, x)(a^n, w)$ for $n \geq 1$
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Proposition

If $TS_1 \sim_{ioblock} TS_2$, then $TS_1 \sim_{io} TS_2$.

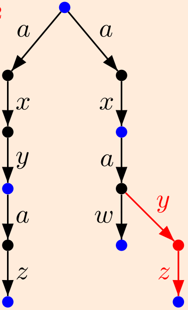
Proof

$IOBeh(TS) = IOBlocks(TS) \cap (\Sigma_i^* \times \Sigma_o^*)$

IO-block equivalence

Example

TS_2



$IOBlocks(TS_2)$:
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(a, x)
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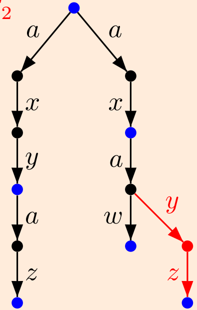

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$IOBeh(TS) = IOBlocks(TS) \cap (\Sigma_i^* \times \Sigma_o^*)$

Decidability of IO-block equivalence

Definition

A transition system is well-structured if every state either refuses Σ_i or refuses Σ_o.

Definition

A block observation $\alpha = (u_1, v_1) \cdots (u_n, v_n)$ is reduced if $u_1 = \varepsilon$ and $u_j \in \Sigma_i$ for $1 < j \leq n$.

$redIOBlocks(TS)$ denotes the set of reduced block observations of TS.

Definition

Let α and β be block-observations. We say that α is *finer than* β, denoted $\alpha \preceq \beta$, if β can be obtained from α by merging consecutive blocks.

Lemma

Let TS be well-structured. Then, $IOBlocks(TS) = \uparrow redIOBlocks(TS)$

where $\uparrow$ denotes the upward closure for $\preceq$.

Decidability of IO-block equivalence

Theorem

For finite well structured transition systems, $\sim_{ioblock}$ is decidable.

Proof

For $w = v_1 a_2 v_2 \cdots a_n v_n \in \Sigma^*$ with $v_j \in \Sigma_o^*$ and $a_j \in \Sigma_i$, we define the reduced block observation $f(w) = (\varepsilon, v_1)(a_2, v_2) \cdots (a_n, v_n)$.

Let $L_\delta(TS)$ be the language accepted by TS with quiescent states as final states.

For $a \in \Sigma_i$, let $L_{\delta,a}(TS)$ be the language accepted by TS with quiescent states that refuse a as final states.

$$redIOBlocks(TS) = f\left(L_\delta(TS) \cup \bigcup_{a \in \Sigma_i} L_{\delta,a}(TS) \cdot a \cdot \Sigma_i^*\right)$$

$$f^{-1}(IOBlocks(TS)) = L_\delta(TS) \cup \bigcup_{a \in \Sigma_i} L_{\delta,a}(TS) \cdot a \cdot \Sigma_i^*$$

Outline

Queue semantics (Tretmans)

Definition

Let $TS = (S, \Sigma, I, T)$ be a transition system. Define $Q(TS) = (S', \Sigma, I', T')$ by

- $S' = S \times \Sigma_i^* \times \Sigma_o^*$: configurations of TS.
- $I' = I \times \{\varepsilon\} \times \{\varepsilon\}$: initial configurations
- Transitions of TS are broken up into two moves, one visible and one invisible (labelled τ):

Input

$$\frac{}{(s, \sigma_i, \sigma_o) \xrightarrow{a} (s, \sigma_i a, \sigma_o)} \qquad \frac{s \xrightarrow{a} s'}{(s, a\sigma_i, \sigma_o) \xrightarrow{\tau} (s', \sigma_i, \sigma_o)}$$

Output

$$\frac{s \xrightarrow{x} s'}{(s, \sigma_i, \sigma_o) \xrightarrow{\tau} (s', \sigma_i, \sigma_o x)} \qquad \frac{}{(s, \sigma_i, x\sigma_o) \xrightarrow{x} (s, \sigma_i, \sigma_o)}$$

- $L(Q(TS))$ is the set of traces of $Q(TS)$.

Queue equivalence (Tretmans)

Definition

$$TS \sim_Q TS' \quad \stackrel{\text{def}}{=} \quad Q(TS) \sim_{syn} Q(TS').$$

Intuitively, synchronous testing equivalence $\sim_{syn}$ corresponds to failure semantics.

Definition

- $w \in L(Q(TS))$ is a quiescent trace if there is a run $(r, \varepsilon, \varepsilon) \xrightarrow{w} (s, \sigma_i, \varepsilon)$ with $r \in I$ and $(s, \sigma_i, \varepsilon)$ quiescent in $Q(TS)$.
- We denote by $L_\delta(Q(TS))$ the set of quiescent traces of $Q(TS)$.

Proposition (Tretmans)

$TS \sim_Q TS' \quad$ iff $\quad L(Q(TS)) = L(Q(TS'))$ and $L_\delta(Q(TS)) = L_\delta(Q(TS'))$

Pb: characterization of $\sim_Q$ on TS instead of $Q(TS)$.

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Ape relation (Tretmans)

Ape relation for the queue semantics

- Output actions may always be postponed: $w_1 x a w_2 \text{ @ } w_1 a x w_2$
  For $x \in \Sigma_o$ and $a \in \Sigma_i$, we have
  $$w_1 x a w_2 \in L(Q(TS)) \text{ implies } w_1 a x w_2 \in L(Q(TS)).$$
- Input actions may always be added: $w \text{ @ } wa$
  For $a \in \Sigma_i$, we have
  $$w \in L(Q(TS)) \text{ implies } wa \in L(Q(TS)).$$
- We denote @ the reflexive and transitive closure of the relations postponing an output action: $w_1 x a w_2 \text{ @ } w_1 a x w_2$
  or adding an input action: $w \text{ @ } wa$.
- $\text{Tracks}(TS)$ is the set of @-minimal words in $L(Q(TS))$.
  @-minimal: no trailing input, outputs as early as possible.
- $L(Q(TS))$ is the @-upward closure of $\text{Tracks}(TS)$.
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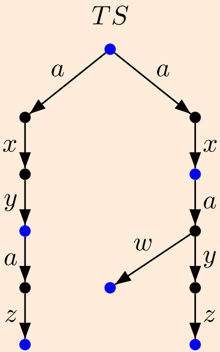
Tracks (Tretmans)

Example



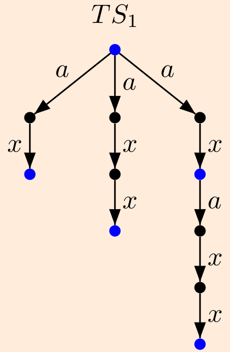
Tracks(TS):	$L(Q(TS))$:
ε	a^*
ax	a^+xa^*
axy	$a^+xa^*ya^*$
$axyaz$	$a^+xa^+ya^*za^*$
$axaw$	$a^+xa^*ya^+za^*$
not $axayz$	$a^+xa^+wa^*$

Comparing the equivalences

Proposition

If $TS_1 \sim_Q TS_2$, then $TS_1 \sim_{io} TS_2$.

The converse does not hold



$\text{Tracks}(TS_1)$:	$IOBeh(TS_1)$:
ε	$(\varepsilon, \varepsilon)$
ax	(a^n, x) for $n \geq 1$
axx	(a^n, x^2) for $n \geq 1$
$axaxx$	(a^n, x^3) for $n \geq 2$

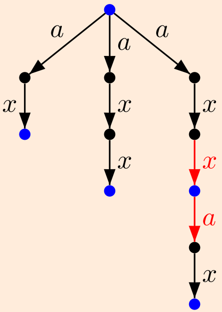
Comparing the equivalences

Proposition

If $TS_1 \sim_Q TS_2$, then $TS_1 \sim_{io} TS_2$.

The converse does not hold

TS_2



$\text{Tracks}(TS_2)$:
ε
ax
axx
$axxax$

$axxax$ @ $axaxx$

$IOBeh(TS_2)$:
$(\varepsilon, \varepsilon)$
(a^n, x) for $n \geq 1$
(a^n, x^2) for $n \geq 1$
(a^n, x^3) for $n \geq 2$

Quiescent traces (Tretmans)

Empty and blocked quiescent traces

- $w \in L(Q(TS))$ is an empty quiescent trace if there is a run $(r, \varepsilon, \varepsilon) \xrightarrow{w} (s, \varepsilon, \varepsilon)$ with $r \in I$ and s quiescent in TS.
  We denote by $L_\delta^{\text{empty}}(Q(TS))$ the empty quiescent traces of $Q(TS)$.
- $w \in L(Q(TS))$ is a blocked quiescent trace if there is a run $(r, \varepsilon, \varepsilon) \xrightarrow{w} (s, a\sigma_i, \varepsilon)$ with $r \in I$ and in TS, s quiescent and $s \stackrel{a}{\not\rightarrow}$.
  We denote by $L_\delta^{\text{block}}(Q(TS))$ the blocked quiescent traces of $Q(TS)$.

Proposition

$$L_\delta(Q(TS)) = L_\delta^{\text{empty}}(Q(TS)) \cup L_\delta^{\text{block}}(Q(TS))$$

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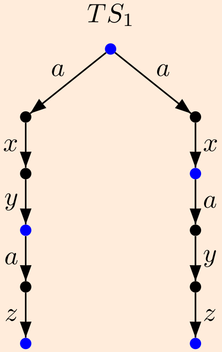
$$L_\delta(Q(TS)) = L_\delta^{\text{empty}}(Q(TS)) \cup L_\delta^{\text{block}}(Q(TS))$$

Quiescent traces (Tretmans)

Example



$L_\delta^{\mathsf{empty}}(TS_1)$:
ε
ax
axy
$axayz$
$axyaz$
$aaxyz$

$L_\delta^{\mathsf{block}}(TS_1)$:
$a^+xya^+za^+$
$a^+xa^+yza^+$
aa^+xyza^+
$a^+xa^+ya^+za^*$
$\ldots$

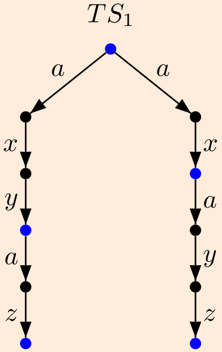
@-upper closure of
$axyaza$

Lemma

$L_\delta^{\mathsf{block}}(Q(TS)) = \{w \in \Sigma^* \mid \exists\, r \xrightarrow{w'} s$ in TS with $r \in I, s$ quiescent, and
$\exists\, a \in \Sigma_i$ such that $s \overset{a}{\nrightarrow}$ and $w'a \sqsubseteq w\}$.

Quiescent traces (Tretmans)

Example



$L_\delta^{\text{empty}}(TS_1)$:
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axy
$axayz$
$axyaz$
$aaxyz$

$L_\delta^{\text{block}}(TS_1)$:
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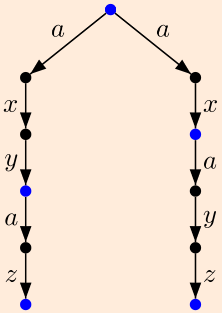
Quiescent traces (Tretmans)

Example

TS_1



$L_\delta^{\mathsf{empty}}(TS_1)$:
ε
ax
axy
$axayz$
$axyaz$
$aaxyz$

$L_\delta^{\mathsf{block}}(TS_1)$:
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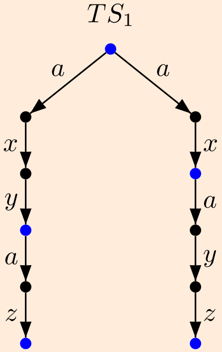
Quiescent traces (Tretmans)

Example



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axy
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$L_\delta^{\text{block}}(TS_1)$:
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@-upper closure of
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Lemma

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$\exists\ a \in \Sigma_i$ such that $s \overset{a}{\not\rightarrow}$ and $w'a \mathbin{@} w\}$.

Strict ape relation (Tretmans)

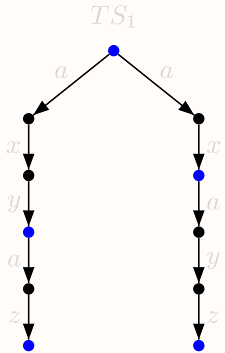
Strict ape relation for the queue semantics

We denote $|@|$ the reflexive and transitive closure of the relation postponing an output action: $w_1 x a w_2 \; @ \; w_1 a x w_2$.

Lemma

$$L_\delta^{\text{empty}}(Q(TS)) = \{w \in \Sigma^* \mid \exists\, r \xrightarrow{w'} s \text{ in } TS \text{ with } r \in I, s \text{ quiescent}, w' \; |@| \; w\}$$

Example



$L_\delta^{\text{empty}}(TS_1)$:
$|@|$-upper closure of
ε
ax
axy
$axyaz$

$axyaz \; |@| \; axayz$
$axyaz \; |@| \; aaxyz$

$L_\delta^{\text{block}}(TS_1)$:
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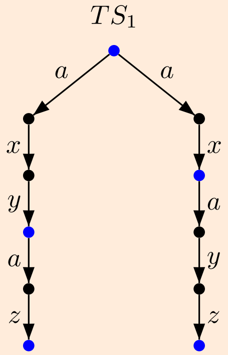
Example



$L_\delta^{\text{empty}}(TS_1)$:
$|@|$-upper closure of
ε
ax
axy
$axyaz$

$axyaz \; |@| \; axayz$
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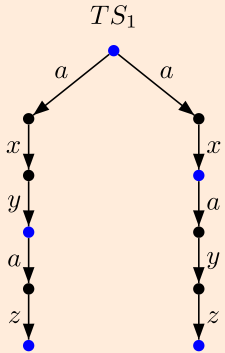
Example



$L_\delta^{\text{empty}}(TS_1)$:
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ε
ax
axy
$axyaz$

$axyaz \;|@|\; axayz$
$axyaz \;|@|\; aaxyz$

$L_\delta^{\text{block}}(TS_1)$:
@-upper closure of
$axyaza$

Undecidability of $\sim_Q$

Theorem

$\sim_Q$ is undecidable

Proof

Reduction from the PCP problem.

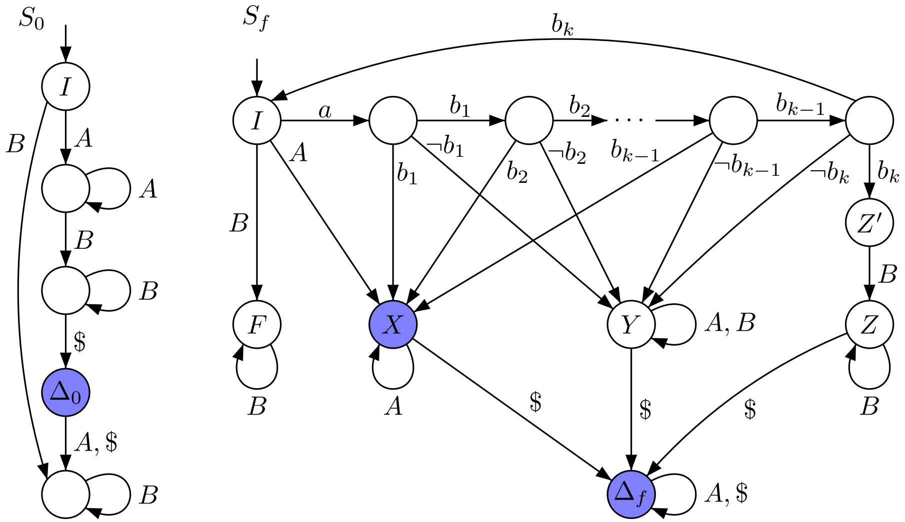
A PCP instance consists in two morphisms $f, g : A^+ \to B^+$ where A, B are finite alphabets.

The PCP instance f, g has a solution if there exists $u \in A^+$ such that $f(u) = g(u)$.

We construct two systems M_1 and M_2 such that the PCP instance (f, g) has no solution iff $M_1 \sim_Q M_2$.

Reduction from the PCP problem

Let $f, g : A^+ \to B^+$ be a PCP instance. We define



Reduction from the PCP problem

We want to compare the following two systems:

- $M_1 = S_0 + S_f + S_g$
- $M_2 = S_f + S_g$

Lemma

$L_\delta^{\text{block}}(M_1) = L_\delta^{\text{block}}(M_2) = \emptyset$.

Lemma

$\text{Tracks}(M_1) = \text{Tracks}(M_2) = \text{Tracks}(S_f) = B^*$.

Lemma

- $L_\delta^{\text{empty}}(S_0)$ is the $|@|$-upper closure of $A^+B^+\$$.
- Let $u \in A^+$ and $v \in B^+$. Then, $uv\$ \in L_\delta^{\text{empty}}(S_f)$ if and only if $v \neq f(u)$.

Theorem

$M_1 \sim_Q M_2$ iff the PCP instance (f, g) has no solution.

Outline

Conclusion

Summary

- We have investigated 3 asynchronous testing equivalences.
- We have shown that $\sim_{io}$ is strictly weaker than $\sim_Q$ and $\sim_{ioblock}$, but $\sim_Q$ and $\sim_{ioblock}$ are incomparable.
- $\sim_{ioblock}$ is decidable, while $\sim_{io}$ and $\sim_Q$ are undecidable.

Open problems

- Construct test suites based on the IO-Blocks semantics.
- Investigate distributed testing.
  See e.g. C. Jard: Synthesis of distributed testers from true-concurrency models of reactive systems, Information & Software Technology, 2003.

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