TD 8: Partial-Order Reduction

Reminder:

(C0) $red(s) = \emptyset$ iff $en(s) = \emptyset$.

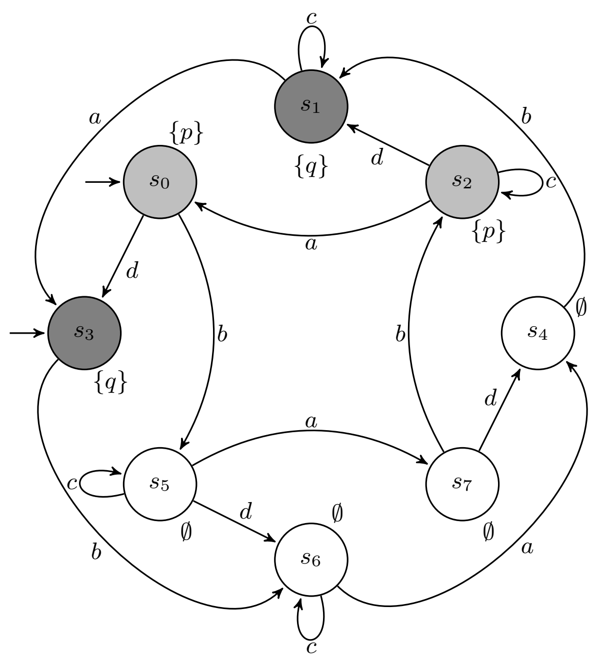
(C1) For every path $s \xrightarrow{a_1} s_1 \xrightarrow{a_2} \cdots \xrightarrow{a_n} s_n \xrightarrow{a} t$ in $\mathcal{K}$ (for any $n \geq 0$), if $a \notin red(s)$ and a depends on some action in $red(s)$ (i.e. there exists $b \in red(s)$ such that $(a, b) \notin I$), then there exists $1 \leq i \leq n$ such that $a_i \in red(s)$.

(C2) If $red(s) \neq en(s)$, then all actions in $red(s)$ are invisible.

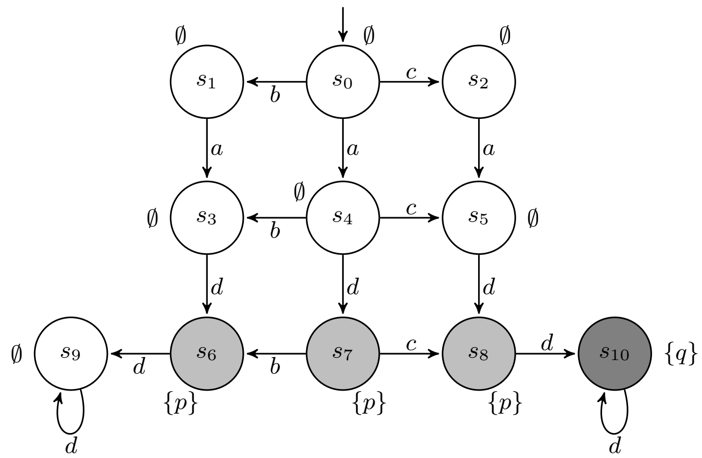
(C3) For all cycles in the reduced system $\mathcal{K}'$, the following holds: if $a \in en(s)$ for some state s in the cycle, then $a \in red(s')$ for some (possibly other) state s' in the cycle.

Exercise 1. Consider the following transition system with state set $S = \{s_0, \ldots, s_7\}$ and transition alphabet $\Delta = \{a, b, c, d\}$:

1. Compute the independance set I and the set of invisible actions U.
2. Propose an assignment $red : S \to 2^{\Delta}$ of ample sets satisfying conditions C_0–C_3 of the lecture notes.



Exercise 2. Consider the following system with $A = \{a, b, c, d\}$:



1. Let $red(s_0) = \{b, c\}$ and $red(s) = en(s)$ for $s \neq s_0$; show that this ample set assignment is compatible with C_0–C_3.

2. Exhibit a CTL(U) formula that distinguishes between the original system and its reduction.

3. Can you propose an assignment that also complies with C_4: if $red(s) \neq en(s)$, then $|red(s)| = 1$?

Exercise 3. Show that (C_0)–(C_2) is not sufficient to ensure stuttering equivalence.

Exercise 4. Let φ be an LTL formula. We define the X-depth $d_{\mathsf{X}}(\varphi)$ and the U-depth $d_{\mathsf{U}}(\varphi)$ of φ as the maximal nesting of X- or U-operators in φ:

$$
\begin{aligned}
d_{\mathsf{X}}(p) &= 0 & d_{\mathsf{U}}(p) &= 0 \\
d_{\mathsf{X}}(\neg\varphi) &= d_{\mathsf{X}}(\varphi) & d_{\mathsf{U}}(\neg\varphi) &= d_{\mathsf{U}}(\varphi) \\
d_{\mathsf{X}}(\varphi \wedge \psi) &= \max(d_{\mathsf{X}}(\varphi), d_{\mathsf{X}}(\psi)) & d_{\mathsf{U}}(\varphi \wedge \psi) &= \max(d_{\mathsf{U}}(\varphi), d_{\mathsf{U}}(\psi)) \\
d_{\mathsf{X}}(\mathsf{X}\,\varphi) &= 1 + d_{\mathsf{X}}(\varphi) & d_{\mathsf{U}}(\mathsf{X}\,\varphi) &= d_{\mathsf{U}}(\varphi) \\
d_{\mathsf{X}}(\varphi \mathbin{\mathsf{U}} \psi) &= \max(d_{\mathsf{X}}(\varphi), d_{\mathsf{X}}(\psi)) & d_{\mathsf{U}}(\varphi \mathbin{\mathsf{U}} \psi) &= 1 + \max(d_{\mathsf{U}}(\varphi), d_{\mathsf{U}}(\psi))
\end{aligned}
$$

We denote by LTL($\mathsf{U}^m, \mathsf{X}^n$) the set of LTL formulas φ with $d_{\mathsf{X}}(\varphi) \leq n$ and $d_{\mathsf{U}}(\varphi) \leq m$, where $n = \infty$ or $m = \infty$ indicates no restriction of the operator in question.

1. We say that two words $w, w' \in \Sigma^\omega$ are *n-stutter-equivalent* if there exists letters $a_0, a_1, \ldots \in \Sigma$ and $f, g : \mathbb{N} \to \mathbb{N}^*$ such that $w = a_0^{f(0)} a_1^{f(1)} \ldots$, $w' = a_0^{g(0)} a_1^{g(1)} \ldots$,

and for all $i \geq 0$, $a_i = a_{i+1}$ implies $a_i = a_j$ for all $j > i$, and $f(i) < n+1$ or $g(i) < n+1$ implies $f(i) = g(i)$.

Show that for all $n \geq 0$ and $\varphi \in \mathrm{LTL}(\mathsf{U}^\infty, \mathsf{X}^n)$, $L(\varphi)$ is closed under n-stutter-equivalence.

2. A similar principle can be formulated when the U-depth is restricted, by considering stuttering of factors instead of letters. Show that for all $m \geq 1$ and $\varphi \in \mathrm{LTL}(\mathsf{U}^m, \mathsf{X}^0)$, for all $u, v \in \Sigma^*$ and $w \in \Sigma^\omega$, we have $uv^m w \in L(\varphi)$ iff $uv^{m+1} w \in L(\varphi)$.