Directed homotopy and homology theories for geometric models of true concurrency

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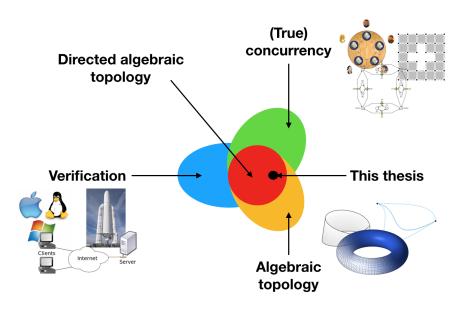
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Applications of directed algebraic topology

Some work done:

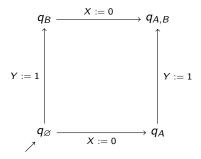
- models for true concurrency and distributed tasks [Pratt, van Glabbeek, Winskel, Nielsen],
- state space reduction techniques [Haucourt, Goubault, Fajstrup, Raussen, Mimram],
- verification (deadlocks/unreachable states detection, correctness via serialisability, decomposition of processes) [Haucourt, Goubault, Fajstrup, Raussen, Mimram, Ninin]

Work in progress:

- fault-tolerance (possibility and impossibility results) [Herlihy, Rajsbaum, Kozlov, Goubault, Mimram, Tasson],
- higher-order rewriting [Malbos],
- Conley-Morse-Forman theory [Mrozek],
- pursuit-evasion games and max flow-min cut duality [Krishnan, Grhist],
- relativity [Dodson, Postdon].

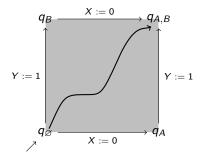
The geometry of true concurrency

Independent actions: interleaving vs true concurrency



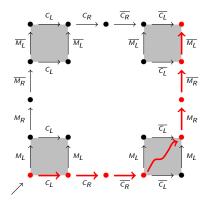
Interleaving behaviors: A then B or B then A

Independent actions: interleaving vs true concurrency



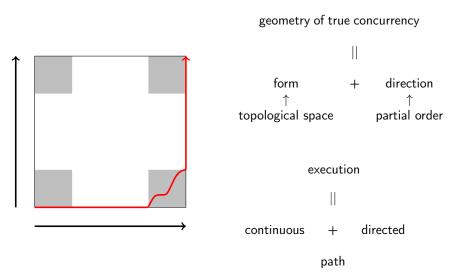
Continuous behaviors: any scheduling of A and B Refinement **[van Glabbeek, Goltz]**: in reality X := 0 and Y := 1 are not atomic

Truly concurrent systems



HDA **[Pratt]** = transition system with higher dimensional data that accounts for true concurrency

The geometry of true concurrency



True concurrency, geometrically

| truly concurrent system | directed space |
|---|-----------------------------|
| states | points |
| executions | directed paths |
| modulo scheduling of independent actions | modulo directed homotopy |

The objective of this thesis

• Understand the geometry of those directed spaces...

• ...with the study of truly concurrent systems through the prism of geometry as a goal.

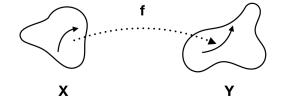
Directed algebraic topology

D-spaces and dipaths [Grandis]

A **d-space** is a topological X with a subset $\overrightarrow{P}(X)$ of paths, called **dipaths**, which contains constant paths, and which is closed under concatenation and non decreasing reparametrizations.

Ex: a partially ordered space with monotonic paths

A **dimap** is a continuous function $f : X \longrightarrow Y$ such that for every $\gamma \in \overrightarrow{P}(X)$, $f \circ \gamma \in \overrightarrow{P}(Y)$. We note **dTop** the category of d-spaces and dimaps.

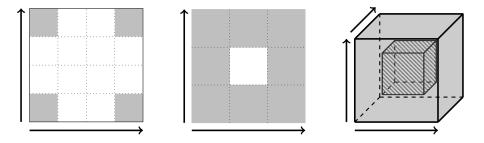


Cubical complexes

Euclidian cubical complex: any subspace of \mathbb{R}^n which is a finite union of cubes of the form

$$[a_1, a_1 + \alpha_1] \times \ldots \times [a_n, a_n + \alpha_n]$$

with $a_i \in \mathbb{Z}$ and $\alpha_i \in \{0, 1\}$.



True concurrency, geometrically

| truly concurrent system | directed space |
|---|-----------------------------|
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True concurrency, geometrically

| truly concurrent system | d-space |
|--|--------------------------------|
| states | points |
| executions | dipaths |
| modulo scheduling of independent actions | modulo directed homotopy ?? |

Dihomotopy of dipaths

A **dihomotopy** from γ to τ , dipaths from *a* to *b*, is a dimap

$$H: \overline{[0,1]} \times \overline{[0,1]} \longrightarrow X$$

such that:

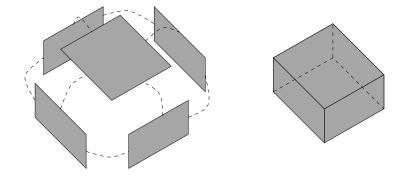
- for every t, H(0, t) = a,
- for every t, H(1, t) = b,
- for every t, $H(t,0) = \gamma(t)$
- for every t, $H(t, 1) = \tau(t)$.

Two dipaths are **dihomotopic** if there is a dihomotopy between them.

dihomotopic

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non-dihomotopic
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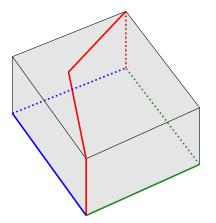
Example: Fahrenberg's matchbox



Example: Fahrenberg's matchbox

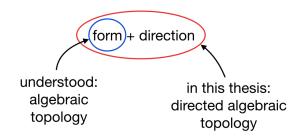
homotopic...

Example: Fahrenberg's matchbox



... but not dihomotopic

Goal of thesis



Extend the work from algebraic topology:

- homology
- homotopy

Homology theories

Homology = counting



hole of dimension 1



not a hole of dimension 1 ... but a hole of dimension 2

 $H_n(X)\simeq \mathbb{R}^{ ext{number}}$ of holes of dimension n $H_0(X)\simeq \mathbb{R}^{ ext{number}}$ of path-connected components

Properties of homology

- **sound**: invariant of homotopy
- precise: do not lose too much information from homotopy [Hurewicz], partially complete [Whitehead]
- **modular**: the homology of a space can be expressed from the homology of simpler spaces [Mayer-Vietoris]
- computable: when the space is finitely presented

What about a directed homology ?

Our contribution: a homology theory for d-spaces with the same kind of properties

Previous proposals of directed homology

- past and future homologies [Goubault 95]
- ordered homology groups [Grandis 04]
- directed homology via ω-categories [Fahrenberg 04]
- homology graph [Kahl 13]
- and some others

The matchbox, a litmus test

ordered homology groups [Grandis 04] = classical homology + partial order

Not precise enough : do not distinguish matchbox from a point



The main ingredients of our directed homology [D.&G.G. ICALP'15]

- 1. Consider trace spaces [Raussen] (\simeq spaces of dipaths).
- 2. Look at how they evolve with time (extensions).
- 3. Apply classical homology on those data.
- 4. Look at the evolution (using bisimulation techniques).

Basic block: space of dipaths ?

$$\overrightarrow{P}(X)(a,b)=$$
 the set of $\gamma\in\overrightarrow{P}(X)$ with $\gamma(0)=a$ and $\gamma(1)=b$

 $\vec{P}(X)(a,b)$ can be equipped with a topology: its path-connected components are the dihomotopy class of dipaths

Problem: concatenation of dipaths is not associative



Basic block: trace spaces [Raussen]

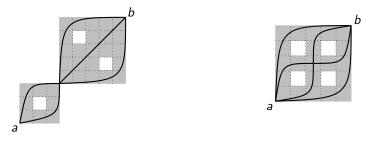
Concatenation is associative modulo reparametrization

Trace = dipath modulo reparametrization

We can also define the **trace space** $\overrightarrow{T}(X)(a, b)$ as the quotient space of $\overrightarrow{P}(X)(a, b)$ modulo reparametrization.

Directed homology = classical homology of a trace space ?

homology of X = classical homology of T(X)(a, b)

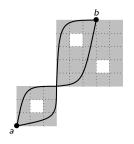


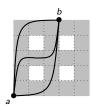
 $(S || U) \bullet (S.U || S.U)$ S.S || U.U

 $T(A)(a, b) \simeq 6 \text{ point space} \simeq T(B)(a, b)$ homology of $A \simeq \mathbb{R}^6 \simeq$ homology of B

Directed homology = evolution of classical homology of trace spaces with time

How to distinguish those two d-spaces ?



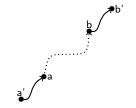




Natural homology

directed homology = collection of modules of homology of the trace spaces indexed by pairs of points, and of linear maps indexed by extensions

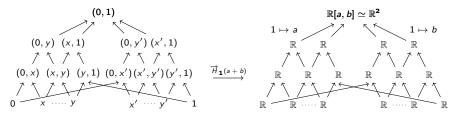
- \mathcal{E}_X = category whose:
 - objects are pairs (a, b) of points such that there is a dipath from a to b
 - morphisms are extensions

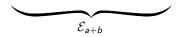


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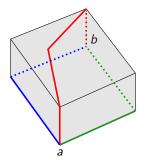
Natural homology [D.&G.G. ICALP'15]: diagram $\overrightarrow{H}_n(X) : \mathcal{E}_X \longrightarrow Mod(\mathcal{R})$ $(a, b) \longmapsto H_{n-1}(\overrightarrow{T}(X)(a, b))$ (H_{n-1} = classical singular homology) Example : first natural homology of a + b







Natural homology on the litmus test



Natural homology detects failure of dihomotopy in the matchbox.

2 dipaths non dihomotopic

$$\Rightarrow \overrightarrow{T}(X)(a,b) \simeq 2 \text{ point space}$$

$$\Rightarrow H_0(\overrightarrow{T}(X)(a,b)) \simeq \mathbb{R}^2$$

$$\Rightarrow \overrightarrow{H}_1(X) \text{ not trivial}$$

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Properties of natural homology

Properties of homology

- **sound**: invariant of homotopy
- precise: do not lose too much information from homotopy [Hurewicz], partially complete [Whitehead]
- **modular**: the homology of a space can be expressed from the homology of simpler spaces [Mayer-Vietoris]
- computable: when the space is finitely presented

Two ways of computing homology

- Direct computation:
 - from a finite presentation of the space, compute a finite presentation of the homology (good for automatization)

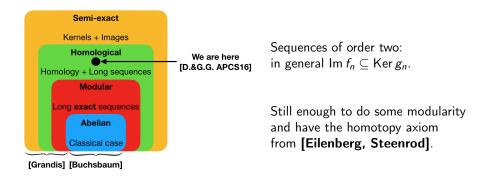
- Modular computation:
 - from a description of a space using simpler spaces, compute the homology of the space from the homology of the spaces involved in the description (ex: spheres)
 - uses the theory of exact sequences

Non-Abelian exactness theory (modularity)

Modularity rely on the possibility to produce long exact sequence:

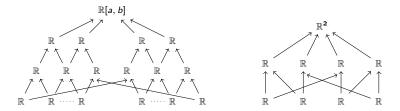
$$\cdots \longrightarrow H_n(V) \xrightarrow{f_n} H_n(W) \xrightarrow{g_n} H_{n-1}(U) \xrightarrow{h_{n-1}} H_{n-1}(V) \longrightarrow \cdots$$

with $\operatorname{Im} f_n = \operatorname{Ker} g_n, \ldots$



How to compare natural homologies ?

Isomorphism of diagrams is the wrong notion to compare diagrams of natural homology (because of the category \mathcal{E}_X)



Key notion [D.&G.G. ICALP'15]:

Compare natural homologies up-to evolution with time. Idea similar to bisimulations in concurrent systems.

Bisimulation of diagrams [D.&G.G. ICALP'15]

Bisimulation between $F : C \longrightarrow Mod(\mathcal{R})$ and $G : D \longrightarrow Mod(\mathcal{R})$ = set R of triples (c, η, d) such that :

- c is an object of C,
- d is an object of D,
- $\eta: F(c) \longrightarrow G(d)$ is an isomorphism of modules

satisfying :

•

• for every object c of C, there exists d and η such that $(c, \eta, d) \in R$

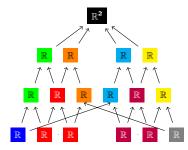
 $\begin{array}{ccc} \mathbf{c} & \mathbf{Fc} & \stackrel{\eta}{\longrightarrow} \mathbf{Gd} & \mathbf{d} \\ i & & \mathbf{Fi} & & & & \\ \mathbf{c'} & \mathbf{Fc'} & \stackrel{\eta'}{\longrightarrow} \mathbf{Gd'} & \mathbf{d'} \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ \end{array}$

 $(c, n, d) \in R$

and symmetrically

Similar to bisimulations of event structures [Rabinovitch, Trakhtenbrot 88].

Examples



the first natural homology of the matchbox is not bisimilar to the one of a point space

Deciding bisimilarity

Theorem [D. submitted]:

It is decidable in EXPSPACE wether two finitary diagrams are bisimilar.

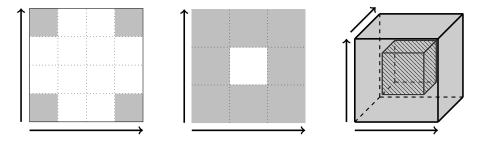
Proof (sketch):

in the finite dimensional case: bisimilarity = problem of matrices in reals. \rightarrow can be encoded in the existential theory of the reals (decidable).

Corollary (soon):

It is decidable wether two finitely presented d-spaces have the same natural homology, up-to bisimilarity.

Finite structure of a Euclidian cubical complex



The trace spaces of a Euclidian cubical complex can be finitely presented, and their homology can be computed [Raussen, Ziemianski]

 \mathcal{E}_X can also be describe finitely using discrete traces = traces which are a glueing of segments joining centers of cubes

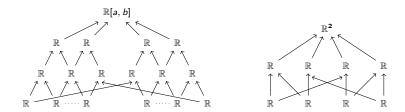


except that everything works only modulo bisimilarity. Jérémy Dubut (LSV, ENS Paris-Saclay) Directed homotopy and homology theories

Computability

Theorem [D.G.G-L. ICALP'15]:

When X is (finite) Euclidian cubical complex, we can compute a finitary diagram bisimilar to $\overrightarrow{H}_n(X)$. It is then decidable wether two such complexes have the same natural homology (up-to bisimilarity).



Homotopy theories and soundness

Soundness in classical algebraic topology

Homotopy equivalence = maps which is a homeomorphism up-to continuous deformations

Classical invariance: homology is an invariant of homotopy equivalence, i.e., a homotopy equivalence induces isomorphisms $H_n(f) : H_n(X) \longrightarrow H_n(Y)$ for all n.

Can we define a notion of dihomotopy equivalence, i.e., of isomorphism of d-spaces up-to continuous deformations that preserve the directed structure, that makes sense and that is compatible with natural homology ?

Dihomotopy of dimaps and dihomotopy equivalence à la Grandis

A dihomotopy from $f : X \longrightarrow Y$ to $g : X \longrightarrow Y$ is a continuous function $H : X \longrightarrow \overrightarrow{P}(Y)$ with $H(_)(0) = f$ and $H(_)(1) = g$, and such that for every t, $H(_)(t)$ is a dimap.

Two dimaps are dihomotopic if there is a zig-zag of dihomotopies between them.

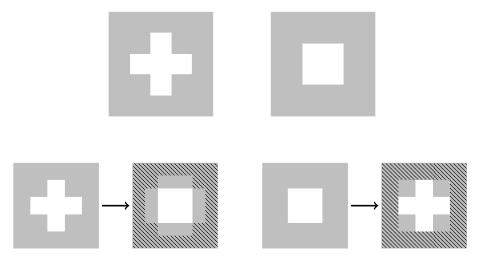
A **dihomotopy equivalence** à la Grandis is a dimap $f : X \longrightarrow Y$ such that there is $g : Y \longrightarrow X$ with $f \circ g$ and $g \circ f$ dihomotopic to identities.

Problem $n^{\circ}1$: the matchbox

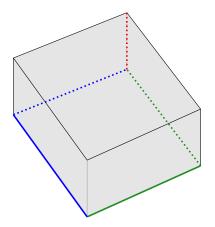


The matchbox is equivalent à la Grandis to a point

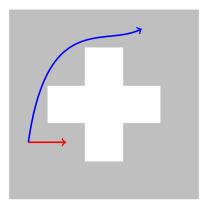
Problem n°2: deadlocks/unreachable states



Solution $n^{\circ}1$: do not follow this path !



Solution $n^{\circ}2$: do not follow this path either !



Inessential dipaths [D.&G.G. CSL'16]

The set $\Im(X)$ of inessential dipaths of X is the largest set of dipaths such that :

- it is closed under concatenation and dihomotopy;
- for every γ ∈ ℑ(X) from x to y, for every z ∈ X such that P(X)(z,x), the map γ ★ _: P(X)(z,x) → P(X)(z,y) δ ↦ γ ★ δ is a homotopy equivalence;
- symmetrically for $_\star\gamma$;
- $\Im(X)$ has the right and left Ore condition modulo dihomotopy:

$$f' \in \mathfrak{I}(X) \qquad \begin{array}{c} g' \\ w & \longrightarrow X \\ w & \text{mod. dihomot.} \end{array} \qquad f \in \mathfrak{I}(X) \qquad \begin{array}{c} z & \longrightarrow Y \\ f \in \mathfrak{I}(X) \\ z & \longrightarrow y \end{array} \qquad \begin{array}{c} f \in \mathfrak{I}(X) \\ f \in \mathfrak{I}(X) \\ x & \longrightarrow W \end{array} \qquad \begin{array}{c} f' \in \mathfrak{I}(X) \\ g' \end{array}$$

Idea similar to the construction of the category of components of a d-space **[Goubault, Haucourt]**

Inessential dihomotopy equivalence [D.&G.G. CSL'16]

A **future inessential deformation retract** (FIDR) of X on a sub-d-space A is a continuous map

H : X - (J(X)

such that :

- for every $x \in X$, H(x)(0) = x;
- for every $a \in A$, $t \in [0,1]$, H(a)(t) = a;
- for every $x \in X$, $H(x)(1) \in A$;
- for every $t \in [0,1]$, the map $H_t : x \mapsto H(x)(t)$ is a dmap;
- + technical conditions

Definition :

Two d-spaces are inessentially homotopically equivalent iff there is a zigzag of FIDR and PIDR between them.

Soundness

Theorem [D.&G.G. CSL'16]:

Natural homology is an invariant of inessential equivalence: if two Euclidian cubical complexes are inessentially equivalent, then their natural homologies are bisimilar.

Conclusion

We have produced a homology theory of d-spaces:

- using trace spaces and their evolution as a basic block,
- using a notion of **bisimulation** to compare the objects in a smoother way,
- and which have directed analogues of properties of the classical homology theory:
 - precise enough to detect non-cancellative default of dihomotopy (matchbox),
 - is compatible with a non-Abelian theory of exactness,
 - is computable in the case of Euclidian cubical complexes,
 - is an invariant of a dihomotopy equivalence (inessential equivalence).



• Is the computation of natural homology can be related to techniques from persistency ?

• What is the precise complexity of the bisimilarity problem ?

• Do we have model structures (or similar structures) for d-spaces up-to dihomotopy equivalences ?