

MPRI 2-7-1
Fondements des systèmes de preuves

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Wednesday, November 21st, 2018

1 hour and a half.
All documents can be used.

1

(5 points)

Let P and Q be proposition symbols (that is, two predicate symbols of arity 0).

- (a) Give a proof in natural deduction of the proposition

$$P \Rightarrow P \Rightarrow (P \Rightarrow Q) \Rightarrow Q$$

- (b) Express this proof as a term of simply typed λ -calculus.
(c) What is the irreducible form of this term?
(d) How many closed irreducible terms of type

$$P \Rightarrow P \Rightarrow (P \Rightarrow Q) \Rightarrow Q$$

are there?

2

(5 points) Let A be the proposition

$$a \leq b$$

- (a) In which language (predicate symbols and function symbols) is this proposition expressed?
(b) Give a model where the proposition A is not valid.
(c) Give a model where $\neg A$ is not valid.
(d) Give a model where neither A nor $\neg A$ are valid.
(e) Give a model where $A \vee \neg A$ is not valid.

3

(5 points)

- (a) Let Q be an atomic proposition and A a non atomic one. Prove that a constructive cut free proof of the sequent $Q \vdash A$ ends with an introduction rule.
- (b) Let Q be an atomic proposition. Prove that there is no constructive cut free proof of the sequent $Q \vdash \perp$.
- (c) Let P and Q be two proposition symbols (that is two predicate symbols of arity 0) and consider the theory in Deduction modulo theory, formed with the rewrite rule

$$P \longrightarrow Q \wedge (P \Rightarrow \perp)$$

Prove that there is no constructive cut free proof of the sequent $Q \vdash \perp$ in this theory.

- (d) Give a constructive proof of the sequent $Q \vdash \perp$ in this theory.
- (e) Is this theory super-consistent?

4

(5 points)

Consider a language with three sorts of terms: *point*, *line* and *scalar*, two predicate symbols $=$ with arity $\langle \text{scalar}, \text{scalar} \rangle$ and \in with arity $\langle \text{point}, \text{line} \rangle$ and two function symbols d , *distance*, with arity $\langle \text{point}, \text{point}, \text{scalar} \rangle$ and b , *bisector*, with arity $\langle \text{point}, \text{point}, \text{line} \rangle$.

Consider the theory formed with the rewrite rule

$$x \in b(y, z) \longrightarrow d(x, y) = d(x, z)$$

and the axiom

$$\forall x \forall y \forall z (x = y \Rightarrow y = z \Rightarrow x = z)$$

Let A a proposition stating that if two bisectors of the triangle xyz intersect at a point w , then the three bisectors intersect at this point:

$$\forall w \forall x \forall y \forall z (w \in b(x, y) \Rightarrow w \in b(y, z) \Rightarrow w \in b(x, z))$$

- (a) What is the irreducible form of this proposition?
- (b) Give a proof of this the proposition A in this theory.
- (c) Is this theory super-consistent?
- (d) Express the proof built at question (b) as a term of the $\lambda\Pi$ -calculus modulo theory.