BAN modified version of CCITT X.509 (3)

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Summary: Modified version of the three messages protocol in the recommendations of the CCITT for the CCITT.X.509 standard (CCITT X.509 (3)).

Protocol specification (in common syntax)

```
A, B:
         principal
Na, Nb:
         nonce
Ya, Yb:
         userdata
Xa, Xb:
         userdata
PK, SK:
         principal -> key (keypair)
                      A, \{Na, B, Xa, \{Ya\}PK(B)\}SK(A)
             В
                      B, \{Nb, A, Na, Xb, \{Yb\}PK(A)\}SK(B)
2.
     В
        ->
3.
             В
                      A, \{B, Nb\}SK(A)
        ->
```

Description of the protocol rules

Compared to CCITT X.509 (3), the identity of B has been added to the signature in message 3. This prevents the [BAN89] attack on the CCITT X.509 (3) protocol, which can occur when B does not check the timestamps. With this modification, the timestamps become redundant and can be removed.

Requirements

```
See CCITT X.509 (3).
```

References

[BAN89].

See also

CCITT X.509 (1), CCITT X.509 (1c), CCITT X.509 (3).

Citations

[BAN89] Michael Burrows, Martin Abadi, and Roger Needham. A logic of authentication. Technical Report 39, Digital Systems Research Center, february 1989.