

Directed homology theories and Eilenberg-Steenrod axioms

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Abstract In this paper, we define and study a homology theory, that we call “natural homology”, which associates a natural system of abelian groups to every space in a large class of directed spaces and precubical sets. We show that this homology theory enjoys many important properties, as an invariant for directed homotopy. Among its properties, we show that subdivided precubical sets have the same homology type as the original ones ; similarly, the natural homology of a precubical set is of the same type as the natural homology of its geometric realization. By same type we mean equivalent up to some form of bisimulation, that we define using the notion of open map. Last but not least, natural homology, for the class of spaces we consider, exhibits very important properties such as Hurewicz theorems, and most of Eilenberg-Steenrod axioms, in particular the dimension, homotopy, additivity and exactness axioms. This last axiom is studied in a general framework of (generalized) exact sequences.

Keywords Algebraic topology, directed topology, homology theory, exact sequences

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1 Introduction

We aim at defining and studying thoroughly a new homology theory, well adapted to directed algebraic topology. Directed algebraic topology [20, 14] is a variant of algebraic topology in which the object of study is directed topological spaces, i.e. topological spaces with a preferred direction (“time”), which cannot be reversed. The concepts were introduced originally in concurrency theory [7, 42], but they are pervasive to many areas: categorically speaking, directed algebraic topology is deeply linked to the study of $(\infty, 1)$ -categories.

Although the theory has now matured, there is still no satisfactory notion of *directed* homology. The first attempt in that direction was made by Goubault [21, 22]. Unfortunately, the abelian character of the homology which he defined made the invariants too weak to be useful for characterizing directed spaces modulo directed homotopy. In particular, there is no hope for having Hurewicz-like theorems, linking (directed) homology to (directed) homotopy. There are in particular simple directed spaces which are not dihomotopic, and which have the same (directed) homology, in all further (directed) homology theories defined since [21], among which [19, 10, 31].

In order to tackle these problems, while keeping the computationally tractable character of homology, we need to define some form of non-abelian homology theory, as put forward in recent work by Krishnan [33]. Our approach is based on a form of Baues-Wirshing homology theory [38, 39]: we define a natural system of abelian groups, which accounts for the variation of the homology of the so-called trace spaces [43], the topological space of directed paths modulo reparametrization from one point to another, when the start and end points change. This was suggested first in [43] and studied in [9].

In this paper, we give an in-depth account of the first notions defined in [9] and carry on the study of this “natural homology”, with a particular focus on Eilenberg-Steenrod axioms, in particular exactness properties, and also Hurewicz types of theorems. This paves the way for practical (directed) homological calculations of concrete directed spaces. Exactness properties are dealt with following the approach for defining generalized homology theories put forward by Grandis [16, 17].

Contents We first recap the central notion of trace space, originally introduced by Martin Raussen [43], and give some simple examples, in Section 2.1. Trace spaces contain the essential (directed) homotopical and homological information, but only when a start and end point are fixed. In Section 2.2, we introduce an algebraic structure, of diagrams, which accounts for the change of start and end points, and collect all trace spaces with varying end points. We then compose these diagrams of spaces with the homotopy, respectively homology, functor, to derive our notion of “natural homotopy” and “natural homology”. The name comes from the fact that these actually form natural systems (of abelian groups for natural homology) in the sense of Baues and Wirshing [38, 39].

In Section 2.3 we recap the notion of precubical set, and of the classical geometric realization of precubical sets (in topological spaces). In order to have a geometric realization which maps onto partially ordered spaces, which are the only types of directed spaces we have been considering in Section 2, we restrict ourselves to particular “geometric” complexes, see Section 2.3. We can then define a notion of natural homology, on these “discrete” spaces, in Section 2.4.

Still, geometric realization does not quite map isomorphically discrete (or combinatorial) natural homology onto the continuous one, at least, not on the nose. In order to relate them, we have to use a notion of bisimulation, which comes originally from concurrency theory [41,37]. We begin by recapping the notion of bisimilarity via open maps in Section 3.1, and define a notion of open map for natural homology. We give other characterizations of bisimilarity in Section 3.2, and show in particular that this is equivalent with the notion of bisimulation of Definition 17, which is simpler to apprehend. We prove in Section 3.3 that indeed, with this notion of bisimilarity, combinatorial and topological homology theories are bisimilar.

Bisimilarity also relates our natural homology, which is a Baues-Wirsing homology theory, with a simpler Hochschild-Mitchell homology theory. This interesting result, both for theoretical and computational purposes, is proved in Section 3.4.

The last part of the paper is dedicated to proving that most of the properties of classical homology theories, such as Hurewicz theorems, or Eilenberg-Steenrod axioms, which are so convenient for practical calculations in many cases, hold, see Section 4. In this section, we have to use a refined notion of exactness, see Section 4.2.2, following the theory of Grandis [16]. We show that our natural homology theory is homological but not modular in general, which is the main difference between our notion of directed homology with standard homology theories. Still, the (relative homology) exactness axiom of Eilenberg-Steenrod is valid, as shown in Section 4.2.5. The dimension, homotopy and additivity axioms are simpler to prove, this is done in Section 4.2.1.

We conclude by sketching some possible future work.

2 Homology of pospaces

Our aim is to define a notion of homology of so-called directed topological spaces, and more precisely, for the purpose of this paper, of partially-ordered topological spaces, or pospaces [40] (see Definition 1). The main ingredient of our homological theories will be the notion of the trace space of a pospace, between two points [43] i.e. the topological space of directed paths modulo reparametrization, see Definition 4.

2.1 Trace spaces

In full generality, directed topology is about d-spaces [20], (pre-)cubical sets [11] or streams [32]. Most of the homological theories we are going to construct can be generalized to such spaces, but we will only prove important properties about these homological theories in some interesting particular cases. In the topological setting, we will restrict ourselves to partially-ordered spaces, whereas in the combinatorial setting, we will restrict our study to particular pre-cubical sets. We indicate in all constructions and theorems whether these are only valid in these particular cases or if this can be generalized.

Partially-ordered spaces, or pospaces, are topological spaces with a simple notion of flow of time:

Definition 1 ([40]) Let \leq be a partial order (reflexive, transitive, antisymmetric relation) on X , a topological space. We say that (X, \leq) is a *pospace* if \leq is a closed subspace of $X \times X$. We note **PoTop** the category of pospace and monotonic continuous functions.

Of course, as for ordinary topological spaces, we are interested in determining whether two pospaces are isomorphic, i.e. when there exists an homeomorphism $f : X \rightarrow Y$, such that f and f^{-1} are monotonic maps. This is the classification problem of pospaces. And as for ordinary topological spaces, this is a very hard problem, this actually includes it as a particular case, when \leq is taken to be equality.

We are looking for algebraic invariants of pospaces modulo directed homeomorphisms, in the form of directed homotopical and homological structures. For defining these structures, we need first to construct paths, and spaces of such. In pospaces, we have the classical notion of (continuous) path, but also the new notion of directed path, i.e. a path which agrees with the flow of time, in a pospace:

Definition 2 Let X be a topological space, $a, b \in X$. A *path from a to b* is a continuous function $f : [0, 1] \rightarrow X$ such that $f(0) = a$ and $f(1) = b$. If moreover X is a pospace, a *dipath from a to b* is a path from a to b which is monotonic (when $[0, 1]$ is equipped with the usual order).

Let f be a (di)path from a to b and g a (di)path from b to c . We define the *concatenation* $f \star g : [0, 1] \rightarrow X$ by: for all $t \in [0, 1/2]$, $(f \star g)(t) = f(2t)$ and for all $t \in [1/2, 1]$, $(f \star g)(t) = g(2t - 1)$. It is a (di)path from a to c .

As usual in algebraic topology, paths are a bit crude - for instance, when it comes to composing them and getting an interesting algebraic structure, i.e. at least associativity of composition, it is necessary to consider them up to reparametrization, at the very least¹:

Definition 3 ([43]) A *monotonic reparametrization* r is a monotonic continuous surjection from $[0, 1]$ to $[0, 1]$.

¹ Or to consider another class of paths, such as Moore paths.

Let X be a pospace and p and q two dipaths from a to b in X . We say that p is *reparametrized in q* if there exists a monotonic reparametrization γ such that $p \circ \gamma = q$. The *trace* of p , written $\langle p \rangle$ is the equivalence class modulo monotonic reparametrization.

Now we can put together all dipaths from point a to point b , modulo monotonic reparametrization in a topological space:

Definition 4 ([43]) Let X be a pospace and a and $b \in X$. We topologize the set of traces of dipaths from a to b , with the compact-open topology. Its quotient $\vec{\mathfrak{X}}(X)(a, b)$ by reparametrization, with the quotient topology is called the *trace space in X from a to b* .

Trace spaces have good theoretical properties (see for instance [43]) but also computational ones. For instance, it is possible to compute a finite representation of the trace space of the hypercube I^n in \mathbb{R}^n minus a set of isothetic hypercubes [44].

This construction can be further generalized to non-looping pre-cubical sets, see e.g. [45], with applications to static analysis of concurrent programs [13].

Example 1 Let us start with simple examples of trace spaces, pictured in Figure 1: these are subspaces $X_1 = [0, 1] \times [0, 1] \setminus [\frac{1}{3}, \frac{2}{3}] \times [\frac{1}{3}, \frac{2}{3}]$ of \mathbb{R}^2 , respectively $X_2 = [0, 1] \times [0, 1] \times [0, 1] \setminus [\frac{1}{3}, \frac{2}{3}] \times [\frac{1}{3}, \frac{2}{3}] \times [\frac{1}{3}, \frac{2}{3}]$ of \mathbb{R}^3 , with componentwise ordering, i.e. $(x, y) \leq (x', y')$, respectively $(x, y, z) \leq (x', y', z')$ if $x \leq x'$ and $y \leq y'$ with the standard ordering in \mathbb{R} , respectively if $x \leq x'$, $y \leq y'$ and $z \leq z'$ in \mathbb{R} .

As shown in [43], $\vec{\mathfrak{X}}(X_1)((0, 0), (1, 1))$ is a topological space with two connected components, one is composed of the traces which have dihomotopy type of the path going up to the left of, then above the hole, the other component is composed of the traces which have the dihomotopy type of the path going along the bottom of the hole then up on its right. Moreover, the two connected components of $\vec{\mathfrak{X}}(X_1)((0, 0), (1, 1))$ are contractible ; it is thus homotopy equivalent to two points.

As shown again in [43], $\vec{\mathfrak{X}}(X_2)((0, 0, 0), (1, 1, 1))$ is homotopy equivalent to the circle S^1 : there is a unique path up to dihomotopy, hence the trace space $\vec{\mathfrak{X}}(X_2)((0, 0, 0), (1, 1, 1))$ is connected, but there is a finer structure of dihomotopies which accounts for the non simple-connectness character of $\vec{\mathfrak{X}}(X_2)((0, 0, 0), (1, 1, 1))$.

2.2 Homological and homotopical diagrams

The study of the shape of such spaces of dipaths, when moving a and b , will be the essential ingredient of invariants under directed homeomorphisms.

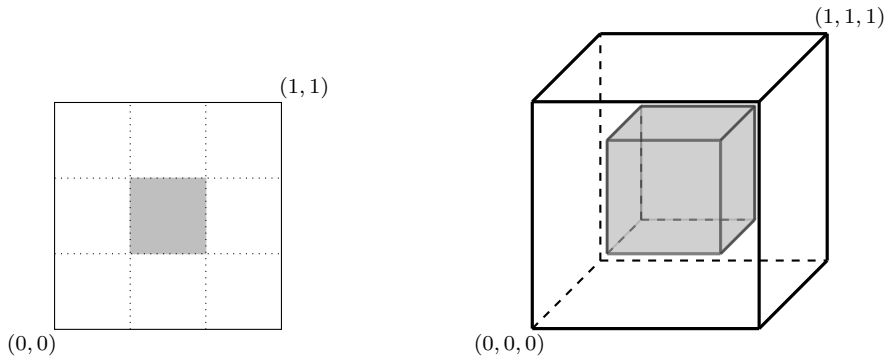


Fig. 1 Two simple partially-ordered subspaces of \mathbb{R}^2 and \mathbb{R}^3 with componentwise ordering which have non-homotopic trace spaces.

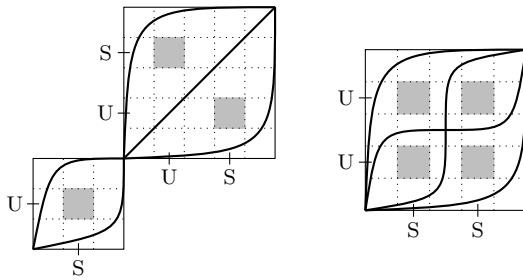


Fig. 2 Examples of pospaces coming from non-equivalent concurrent programs.

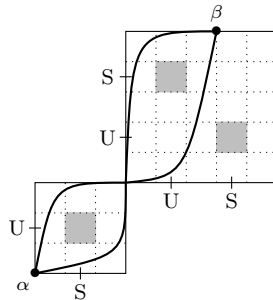


Fig. 3 Changing the base points for exhibiting a particular trace space.

Example 2 In Figure 2, we have depicted two spaces, which are built as the gluing of squares (the white ones), each of which is equipped with the product ordering of \mathbb{R}^2 . They will come as the geometric realization of some precubical sets, see Section 2.3. They are not dihomeomorphic spaces since they are already non homotopy equivalent: the fundamental group of the leftmost one, that we call X , is the free abelian group on three generators, whereas the fundamental group of the rightmost one, that we call Y , is the free abelian group on four generators. Consider now our construction of trace space, for X and

Y , from the lowest point of X (resp. Y), that we call α_X (resp. α_Y), to the highest point of X (resp. Y), that we call β_X (resp. β_Y). We can see easily that $\vec{\mathfrak{X}}(X)(\alpha_X, \beta_X)$ is homotopy equivalent to a six point space (corresponding to the six dihomotopy classes of dipaths pictured in 2), and that $\vec{\mathfrak{X}}(Y)(\alpha_Y, \beta_Y)$ is also homotopy equivalent to a six point space (corresponding again to the six dihomotopy classes of dipaths pictured in Figure 2).

It can be shown that those two programs are not equivalent in the sense that they do not act the same way. But their trace spaces have the homotopy type, the one of a discrete space with six points.

Our main idea is then to see how homotopy type of trace spaces vary when we move the points between which we take the traces.

With the possibility to consider all the trace spaces, we can distinguish the two former pospaces because if we consider the trace space between α and β , as in Figure 3, it has the homotopy type of a discrete space with four points but we can show that, in Y , there is no pair of points between which we have a trace space with the same homotopy type.

Let X be a pospace, we first define the “base” category on which to vary the end points of traces, and on top of which we will observe the variation of topology of trace spaces.

Definition 5 We define \mathcal{T}_X to be the category whose:

- objects are traces of X
- morphisms (also called extensions) from $\langle p \rangle$ to $\langle q \rangle$ with p , a dipath from x to y and q , one from x' to y' are pairs of traces $(\langle \alpha \rangle, \langle \beta \rangle)$ such that $\langle q \rangle = \langle \alpha \star p \star \beta \rangle$

We then define $\vec{T}_*(X) : \mathcal{T}_X \longrightarrow \mathbf{Top}_*$ which maps:

- every trace $\langle p \rangle$ with p from x to y to the pointed space $(\vec{\mathfrak{X}}(X)(x, y), \langle p \rangle)$
- every extension $(\langle \alpha \rangle, \langle \beta \rangle)$ with α dipath from x' to x and β dipath from y to y' to the continuous map $\langle \alpha \star _ \star \beta \rangle : \vec{\mathfrak{X}}(X)(x, y) \longrightarrow \vec{\mathfrak{X}}(X)(x', y')$ which maps $\langle p \rangle$ to $\langle \alpha \star p \star \beta \rangle$.

Diagrams in a category themselves form a category, which we make explicit below:

Definition 6 (The category of diagrams) Let \mathcal{M} be a category. We define $\mathbf{Diag}(\mathcal{M})$ the category whose:

- objects are diagrams i.e. functors from any small category \mathcal{C} to \mathcal{M}
- morphisms from $F : \mathcal{C} \longrightarrow \mathcal{M}$ to $G : \mathcal{D} \longrightarrow \mathcal{M}$ are pairs (Φ, σ) where
 - Φ is a functor from \mathcal{C} to \mathcal{D}
 - σ is a natural transformation from F to $G \circ \Phi$
- the identity on $F : \mathcal{C} \longrightarrow \mathcal{M}$ is $id_F = (id_{\mathcal{C}}, 1_F)$ where
 - $id_{\mathcal{C}}$ is the identity functor on \mathcal{C}
 - 1_F is the identity natural transformation on F

- the composition is defined as follows: $(\Psi, \tau) \circ (\Phi, \sigma)$ where $(\Phi, \sigma) : (F : \mathcal{C} \rightarrow \mathcal{M}) \rightarrow (G : \mathcal{D} \rightarrow \mathcal{M})$ and $(\Psi, \tau) : (G : \mathcal{D} \rightarrow \mathcal{M}) \rightarrow (H : \mathcal{E} \rightarrow \mathcal{M})$ is $(\Psi \circ \Phi, (\tau_{\Phi(c)} \circ \sigma_c)_{c \in \text{Ob}(\mathcal{M})})$

Now, trace diagrams can be seen as a functor into diagrams in \mathbf{Top}_* :

Proposition 1 \vec{T} extends to a functor from \mathbf{PoTop} to $\mathbf{Diag}(\mathbf{Top}_*)$.

Proof

If $f : X \rightarrow Y$ is a dimap, we define $\vec{T}(f) = (\Phi, \sigma) : \vec{T}(X) \rightarrow \vec{T}(Y)$ as follows:

- $\Phi : \mathcal{T}_X \rightarrow \mathcal{T}_Y$ such that $\Phi(\langle p \rangle) = \langle f \circ p \rangle$ and $\Phi(\langle \alpha \rangle, \langle \beta \rangle) = (\langle f \circ \alpha \rangle, \langle f \circ \beta \rangle)$
- if p is a dipath from x to y , $\sigma_{(p)} : \vec{\mathfrak{T}}(X)(x, y) \rightarrow \vec{\mathfrak{T}}(Y)(f(x), f(y))$
 $\langle p \rangle \mapsto \langle f \circ p \rangle$

□

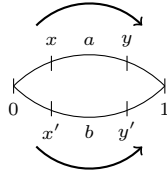
From trace diagrams in \mathbf{Top}_* we can derive relevant homological and homotopical information by composing with the classical singular homology (resp. homotopy) functor in the category \mathbf{Ab} of abelian groups (resp. \mathbf{Grp})

Definition 7 (Natural homotopy) We define for $n \geq 1$, $\vec{H}_n(X) : \mathcal{T}_X \rightarrow \mathcal{M}$ (where \mathcal{M} is either \mathbf{Set} , \mathbf{Grp} or \mathbf{Ab}) composing $\vec{T}_*(X)$ with the $(n-1)^{\text{th}}$ homotopy group (set if $n = 1$) functor π_{n-1} .

Definition 8 (Natural homology) We define for $n \geq 1$, $\vec{H}_n(X) : \mathcal{T}_X \rightarrow \mathcal{M}$ (where \mathcal{M} is \mathbf{Ab}) composing $\vec{T}_*(X)$ with the $(n-1)^{\text{th}}$ homology group functor H_{n-1} .

Remark 1 \mathcal{T}_X is actually the category of factorization (or twisted arrow category [34]) of the category whose objects are points of X and morphisms are traces and this makes $\vec{T}_*(X)$ into a natural system in the sense of [1].

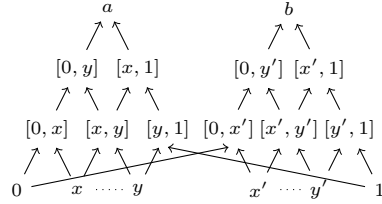
Example 3 We consider the following pospace, denoted $a + b$, which is made up of two directed segments a and b where their initial points are identified, and their final points are identified too. In the following picture, we distinguish two particular points x and y on a , with $x < y$ (respectively x' and y' on b , with $x' < y'$), which we will use to describe the category of factorization \mathcal{T}_{a+b} as well as the natural homology $\vec{H}_n(a + b)$.



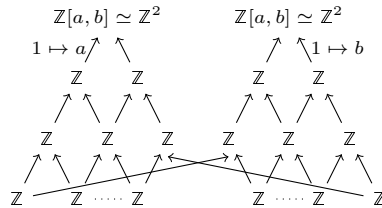
The description of \mathcal{T}_{a+b} is now as follows. Objects of \mathcal{T}_{a+b} are dipaths, which can be either:

- constant dipaths, $0, x, y, x', y', 1$, for all points x, y, x', y' that we chose to distinguish in the picture of $a + b$.
- non constant and non maximal dipaths of the form $[0, x], [x, y], [y, 1]$ etc.
- maximal dipaths a and b

We chose below to draw a picture of a subcategory of \mathcal{T}_{a+b} , where x, y, x' and y' are any distinguished points of a and b as discussed before. The extension morphisms in \mathcal{T}_{a+b} are pictured below as arrows ; for instance, there is an extension morphism from dipath $[x, y]$ to $[0, y]$ and to $[x, 1]$, among other extension morphisms:



Now, we can picture a subdiagram of $\vec{H}_1(a+b)$, by applying the homology functor on the trace spaces from the starting point to the end point of the dipaths, objects of \mathcal{T}_{a+b} . For instance, the trace space $\vec{\mathfrak{Z}}(a+b)(x, y)$ (respectively $\vec{\mathfrak{Z}}(a+b)(0, y)$) corresponding to dipath $[x, y]$ (respectively $[0, y]$) in the diagram above, is just a point, hence has zeroth homology group equal to \mathbb{Z} (respectively \mathbb{Z}). All other zeroth homology groups are trivial with the exception of the ones corresponding to the two maximal dipaths (up to reparametrization) a and b , going from 0 to 1 . In that case, $\vec{\mathfrak{Z}}(a+b)(0, 1)$ is composed of two points, that we can identify with a and b , and has \mathbb{Z}^2 (or $\mathbb{Z}[a, b]$ with the identification we just made) as zeroth homology. Now the extension morphism from $[0, y]$ to a induces a map in homology which maps the only generator of $H_0(\vec{\mathfrak{Z}}(a+b)(0, y))$ to generator a in $\mathbb{Z}[a, b]$ as indicated in the picture below:



2.3 Pre-cubical sets, cubical complexes

Now, we turn to an alternative definition of natural homology and homotopy, on a combinatorial version of pospaces. As in the classical case, we want to express that our combinatorial and topological homology and homotopy theories are strongly related. This will be Theorem 1.

Definition 9 A *pre-cubical set* K is a graded set $(K_n)_{n \in \mathbb{N}}$ together with (face) maps $\partial_i^\alpha : K_n \rightarrow K_{n-1}$ with $1 \leq i \leq n$ and $\alpha \in \{0, 1\}$ satisfying:

$$\partial_i^\alpha \circ \partial_j^\beta = \partial_{j-1}^\beta \circ \partial_i^\alpha$$

In order to relate (some) pre-cubical sets with pospaces, we need to define a geometric realization, hence first, topological “models” for abstract n -cubes :

Definition 10 Let $\square_n = \{(t_1, \dots, t_n) \in \mathbb{R}^n \mid \forall 1 \leq i \leq n, 0 \leq t_i \leq 1\}$ the standard n -cube in \mathbb{R}^n .

We define for $\alpha \in \{0, 1\}$, $n \in \mathbb{N}$, $1 \leq i \leq n+1$, $\rho_i^\alpha : \square_n \rightarrow \square_{n+1}$ by:

$$\rho_i^\alpha(t_1, \dots, t_n) = (t_1, \dots, t_{i-1}, \alpha, t_i, \dots, t_n)$$

Let K be a pre-cubical set. We note $R(K) = \bigsqcup_{n \in \mathbb{N}} K_n \times \square_n$, the topological space constructed as the disjoint union (with disjoint topology) of the product (with the product topology) of K_n (with the discrete topology) and \square_n (with the topology of \mathbb{R}^n).

The elements of $R(K)$ are pairs (e, \mathbf{a}) where e is an n -dimensional cube in K and $\mathbf{a} \in [0, 1]^n$, for some n .

The *geometric realization* $Geom(K)$ of K is the quotient of $R(K)$ by the least equivalence relation \equiv such that:

$$\forall \alpha \in \{0, 1\}, n \in \mathbb{N}, 1 \leq i \leq n, x \in K_n, t \in \square_{n-1}, (\partial_i^\alpha(x), t) \equiv (x, \rho_i^\alpha(t))$$

with the quotient topology.

We shall write $[e, \mathbf{a}]$ for the point obtained as the equivalence class of (e, \mathbf{a}) .

Geometric realizations of pre-cubical sets are directed spaces of a much more general type than the pospaces we have been considering up to now. We restrict our study to particular “geometric” complexes, which are realizations of specific pre-cubical sets and form pospaces:

Definition 11 A (d -dimensional) *cubical complex* K is a finite set of *cubes* (D, \mathbf{x}) , where $D \subseteq \{1, 2, \dots, d\}$ and $\mathbf{x} \in \mathbb{Z}^d$, which is closed under taking past and future faces (to be defined below). The cardinality $|D|$ of D is the *dimension* of the cube (D, \mathbf{x}) . Let $\mathbf{1}_k$ be the d -tuple whose k th component is 1, all others being 0. Each cube (D, \mathbf{x}) is *realized* as the geometric cube $\iota(D, \mathbf{x}) = I_1 \times I_2 \times \dots \times I_d$ where $I_k = [x_k, x_k + 1]$ if $k \in D$, $I_k = [x_k, x_k]$ otherwise, matching the definition of [30].

When $|D| = n$, we write $D[i]$ for the i th element of D . For example, if $D = \{3, 4, 7\}$, then $D[1] = 3$, $D[2] = 4$, $D[3] = 7$. We also write $\partial_i D$ for D minus $D[i]$. Every n -dimensional cube (D, \mathbf{x}) has n *past faces* $\partial_i^0(D, \mathbf{x})$, defined as $(\partial_i D, x)$, and n *future faces* $\partial_i^1(D, \mathbf{x})$, defined as $(\partial_i D, x + \mathbf{1}_{D[i]})$, $1 \leq i \leq n$.

Proposition 2 A *cubical complex* is a *pre-cubical set*. Moreover,

$$\begin{aligned} \epsilon : Geom(K) &\rightarrow \bigsqcup_{(D, \mathbf{x}) \in K} \iota(D, \mathbf{x}) \\ [(D, \mathbf{x}), a] &\mapsto \mathbf{x} + \sum_{i=1}^n a_i \mathbf{1}_{D[i]} \end{aligned}$$

is a homeomorphism and thus induces a pospace structure on $\overrightarrow{Geom}(K)$ by $[(D, \mathbf{x}), a] \leq [(D', \mathbf{x}'), a']$ iff $\epsilon([(D, \mathbf{x}), a]) \leq \epsilon([(D', \mathbf{x}'), a'])$. We note $\overrightarrow{Geom}(K)$ this pospace.

2.4 Combinatorial natural homology

Let K be a pre-cubical set. In this section, we define the analogue of (continuous) traces in the (discrete) setting of pre-cubical set, so that to build a factorization category (Definition 14) and then a natural homology (Definition 15).

Definition 12 Let x and $y \in K$. We say that x is a *future boundary* (resp. a *past boundary*) of y if there exist $k \geq 0$ and i_0, \dots, i_k such that $x = \partial_{i_k}^1 \circ \dots \circ \partial_{i_0}^1(y)$ (resp. $x = \partial_{i_0}^0 \circ \dots \circ \partial_{i_k}^0(y)$). We note $x \preceq y$ when:

- either x is a past boundary of y
- either y is a future boundary of x

Definition 13 A *discrete trace* from x to y in K is a sequence c_0, \dots, c_n of K (with $n \geq 0$) such that $c_0 = x$, $c_n = y$ and for all $i \in \{1, \dots, n\}$ $c_{i-1} \preceq c_i$. An extension of a discrete trace c_0, \dots, c_n is a pair $((d_0, \dots, d_m), (e_0, \dots, e_p))$ of discrete traces such that $d_m = c_0$ and $c_n = e_0$.

Similarly as for pospaces (Definition 5), discrete traces of a pre-cubical set K form a category \mathcal{T}_K^d whose objects are discrete traces, and whose morphisms are extensions of discrete traces which are discrete traces.

We note first that when K is a cubical complex, a discrete trace can be realized as a trace in the geometric realization $\overrightarrow{Geom}(K)$. For x in K , write \hat{x} the point $[x, \bullet]$ of $\overrightarrow{Geom}(K)$ where $\bullet = (\frac{1}{2}, \dots, \frac{1}{2})$. When x is a past boundary of y , i.e. $x = \partial_{i_k}^0 \circ \dots \circ \partial_{i_0}^0(y)$, write \widehat{xy} the path in $\overrightarrow{Geom}(K)$ from \hat{x} to \hat{y} defined by

$$\widehat{xy}(t) = [y, t\bullet + (1-t)\rho_{i_0}^0 \circ \dots \circ \rho_{i_k}^0(\bullet)]$$

When y is a future boundary of x , i.e. $y = \partial_{i_k}^1 \circ \dots \circ \partial_{i_0}^1(x)$, write \widehat{xy} the path in $\overrightarrow{Geom}(K)$ from \hat{x} to \hat{y} defined by

$$\widehat{xy}(t) = [x, (1-t)\bullet + t\rho_{i_0}^1 \circ \dots \circ \rho_{i_k}^1(\bullet)]$$

\widehat{xy} is uniquely defined when K is a non-self linked pre-cubical set [12] and is a dipath when K is a cubical complex.

Then, a discrete trace c_0, \dots, c_n can be realized as the trace $\langle \widehat{c_0 c_1} \star \dots \star \widehat{c_{n-1} c_n} \rangle$ in $\overrightarrow{Geom}(K)$, noted $\widehat{c_0, \dots, c_n}$.

Definition 14 We define \mathcal{T}_K^d the category whose:

- objects are discrete traces of K

- morphisms from c_0, \dots, c_n (discrete trace from x to y) to d_0, \dots, d_m (discrete trace from x' to y') are pairs of discrete traces $((e_0, \dots, e_p), (f_0, \dots, f_k))$ from (x', y) to (x, y') (so $c_0 = e_p = x$ and $c_n = f_0 = y$) such that $d_0, \dots, d_m = e_0, \dots, e_{p-1}, c_0, \dots, c_n, f_1, \dots, f_k$.

Then $\widehat{}$ extends to a functor from \mathcal{T}_K^d to $\overrightarrow{\mathcal{T}}_{\text{Geom}(K)}$.

Definition 15 Let K be a cubical complex. For $n \geq 1$, we define the n -th natural homology of K or discrete natural homology $\overrightarrow{h}_n(K) : \mathcal{T}_K^d \rightarrow \mathbf{Ab}$ as

$$\overrightarrow{h}_n(K) = \overrightarrow{H}_n(\overrightarrow{\text{Geom}}(K)) \circ \widehat{}$$

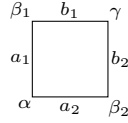
But natural homology in the combinatorial version does not quite correspond to natural homology in the topological case, as we exemplify now:

Example 4 We continue with Example 3. The space we called $a + b$ is dihedral-omorph to the digeometric realization of the following 2-dimensional cubical complex K_{a+b} :

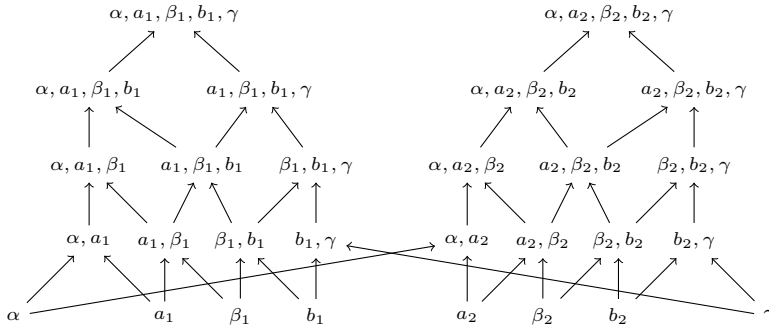
$$\{\alpha = (\emptyset, (0, 0)), \beta_1 = (\emptyset, (0, 1)), \beta_2 = (\emptyset, (1, 0)), \gamma = (\emptyset, (1, 1)),$$

$$a_1 = (\{2\}, (0, 0)), b_1 = (\{1\}, (0, 1)), a_2 = (\{1\}, (0, 0)), b_2 = (\{2\}, (1, 0))\}$$

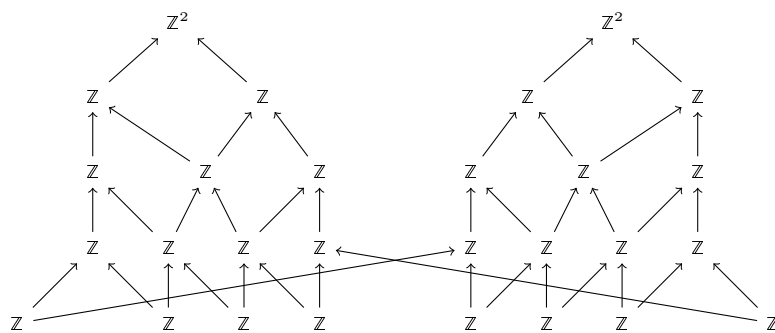
which can be geometrically depicted like this:



The category $\mathcal{T}_{K_{a+b}}^d$ is:



The image of $\overrightarrow{h}_1(K_{a+b})$ is:



Still, the difference between this (discrete) natural homology with what we had in the continuous case, in Example 3, is somehow inessential. We now define an weaker equivalence than isomorphism of diagrams that abstract away this difference. This is based on the notion of bisimulation, which has been introduced in concurrency theory [41,37]. We will see in Section 3.1 that this equivalence can be seen as a bisimulation equivalence defined by a notion of open maps [29].

3 Bisimulation of diagrams

3.1 \mathcal{P} -bisimilarity of diagrams

First, we recap the general construction of bisimulations with open maps [29], we then apply it to our category of diagrams.

Definition 16 Given a category \mathcal{N} and a subcategory \mathcal{P} (called the category of paths), we say that a morphism $f : X \rightarrow Y$ of \mathcal{N} is \mathcal{P} -open if for every diagram:

$$\begin{array}{ccc}
 P & \xrightarrow{p} & X \\
 g \downarrow & & \downarrow f \\
 Q & \xrightarrow{q} & Y
 \end{array}$$

where $g : P \rightarrow Q$ is a morphism of \mathcal{P} , then there exists $r : Q \rightarrow X$ making the following diagram commutative:

$$\begin{array}{ccc}
 P & \xrightarrow{p} & X \\
 g \downarrow & \nearrow r & \downarrow f \\
 Q & \xrightarrow{q} & Y
 \end{array}$$

In other words, f has the right lifting property with respect to morphisms of \mathcal{P} .

We then say that two objects X and Z of \mathcal{N} are \mathcal{P} -bisimilar if there exists a span of open maps between them:

$$\begin{array}{ccc} & Z & \\ f \swarrow & & \searrow g \\ X & & Y \end{array}$$

We consider now the category $\mathcal{N} = \mathbf{Diag}_{iso}(\mathcal{M})$, which is the subcategory of $\mathbf{Diag}(\mathcal{M})$ whose morphisms are of the form (Φ, σ) where σ is a natural isomorphism. We consider \mathcal{P} the following subcategory of \mathcal{N} : its objects are the functors of \mathbf{n} in \mathcal{M} for any $n \in \mathbb{N}$ (we will call them paths of length n) where \mathbf{n} is the category induced by the natural total order on $\{1, \dots, n\}$, i.e.:

$$1 \longrightarrow 2 \longrightarrow \dots \longrightarrow n$$

and its morphisms are the morphisms of $\mathbf{Diag}_{iso}(\mathcal{M})$ of the form $(ext_{n,m}, id)$ with $n \leq m$ where $ext_{n,m} : \mathbf{n} \rightarrow \mathbf{m}$ is the functor induced by the increasing function $(i \mapsto i)$. We will also note those morphisms $ext_{n,m}$.

Lemma 1 *A morphism $(\Phi, \sigma) : (F : \mathcal{C} \rightarrow \mathcal{M}) \rightarrow (G : \mathcal{D} \rightarrow \mathcal{M}) \in \mathbf{Diag}_{iso}(\mathcal{M})$ is \mathcal{P} -open iff it has the right lifting property for $ext_{n,n+1}$ for all $n \in \mathbb{N}$.*

Proof

\Rightarrow by definition of open maps

\Leftarrow we show that it has the right lifting property with respect to $ext_{n,m}$, using the fact that $ext_{n,m} = ext_{m-1,m} \circ \dots \circ ext_{n,n+1}$. □

3.2 Equivalent definitions of \mathcal{P} -bisimilarity of diagrams

3.2.1 Via a simpler notion of open maps

Proposition 3 *A morphism $(\Phi, \sigma) : (F : \mathcal{C} \rightarrow \mathcal{M}) \rightarrow (G : \mathcal{D} \rightarrow \mathcal{M}) \in \mathbf{Diag}_{iso}(\mathcal{M})$ is \mathcal{P} -open iff it has the right lifting property for $ext_{0,1}$ and $ext_{1,2}$ i.e. iff:*

- Φ is surjective on objects
- for every morphism $j : d \rightarrow d'$ of \mathcal{D} and every c of \mathcal{C} such that $\Phi(c) = d$ there exists $i : c \rightarrow c'$ such that $\Phi(i) = j$.

Proof

\Rightarrow by definition of open maps

\Leftarrow Let us show that (Φ, σ) as the right lifting property with respect to $ext_{n,n+1}$. The case 0 is by hypothesis. Assume $n \geq 1$ and given a commutative diagram:

$$\begin{array}{ccccc} P : \mathbf{n} & \longrightarrow & \mathcal{M} & \xrightarrow{(p, \rho)} & F : \mathcal{C} \longrightarrow \mathcal{M} \\ (ext_{n,n+1}, id) \downarrow & & & & \downarrow (\Phi, \sigma) \\ Q : \mathbf{n+1} & \longrightarrow & \mathcal{M} & \xrightarrow{(q, \eta)} & G : \mathcal{D} \longrightarrow \mathcal{M} \end{array}$$

We want (r, θ) such that:

$$\begin{array}{ccccc} P : \mathbf{n} & \longrightarrow & \mathcal{M} & \xrightarrow{(p, \rho)} & F : \mathcal{C} \longrightarrow \mathcal{M} \\ (ext_{n,n+1}, id) \downarrow & & & \nearrow (r, \theta) & \downarrow (\Phi, \sigma) \\ Q : \mathbf{n+1} & \longrightarrow & \mathcal{M} & \xrightarrow{(q, \eta)} & G : \mathcal{D} \longrightarrow \mathcal{M} \end{array}$$

commutes.

$r : \mathbf{n+1} \rightarrow \mathcal{C}$ is constructed this way:

- for $1 \leq i \leq n$, $r(i) = p(i)$ and for all $j \leq i$, $r(j \leq i) = p(j \leq i)$
- it remains to construct $r(n \leq n+1)$ (which will determine $r(n+1)$).

We know that $q(n \leq n+1) : q(n) \rightarrow q(n+1)$ and that $q(n) = \Phi(p(n))$.

Thus by the second property of Φ , there exists a morphism $i : p(n) = r(n) \rightarrow c$ of \mathcal{C} such that $\Phi(i) = q(n \leq n+1)$. We pose $r(n \leq n+1) = i$.

First, it is easy to see that $p = r \circ ext_{n,n+1}$ and $\Phi \circ r = q$.

$\theta = (\theta_i)_{i \leq n+1} : Q \rightarrow F \circ r$ is define this way:

- for all $i \leq n$, $\theta_i = \rho_i = \sigma_{q(i)} \circ \eta_i$ which is an isomorphism
- $\theta_{n+1} = \sigma_{q(n+1)} \circ \eta_{n+1}$ which is an isomorphism

θ is natural because it is the composition of two natural transformations.

By construction, it makes the diagram commutes. \square

By abuse of notation, we will say that a morphism (Φ, σ) of $\mathbf{Diag}(\mathcal{M})$ is open if σ is a natural isomorphism and Φ satisfies the conditions of the previous proposition.

3.2.2 A relational characterization of bisimilarity

We now turn to a more classical characterization of bisimulation for our diagrams, that relates to theoretical computer science and concurrency theory :

Definition 17 A bisimulation R between two diagrams $F : \mathcal{C} \rightarrow \mathcal{M}$ and $G : \mathcal{D} \rightarrow \mathcal{M}$ is a set of triples (c, f, d) where c is an object of \mathcal{C} , d is an object of \mathcal{D} and $f : F(c) \rightarrow G(d)$ is an isomorphism of \mathcal{M} such that for all (c, f, d) in R :

- if there exists $i : c \rightarrow c' \in \mathcal{C}$ then there exists $j : d \rightarrow d' \in \mathcal{D}$ and $g : F(c') \rightarrow G(d') \in \mathcal{M}$ such that $g \circ F(i) = G(j) \circ f$ and $(c', g, d') \in R$
- if there exists $j : d \rightarrow d' \in \mathcal{D}$ then there exists $i : c \rightarrow c' \in \mathcal{C}$ and $g : F(c') \rightarrow G(d') \in \mathcal{M}$ such that $g \circ F(i) = G(j) \circ f$ and $(c', g, d') \in R$

$$\begin{array}{ccc} c & Fc & \xrightarrow{f} Gd & d \\ i \downarrow & Fi \downarrow & & \downarrow Gj \downarrow j \\ c' & Fc' & \xrightarrow{g} Gd' & d' \end{array}$$

and such that:

- for all $c \in \mathcal{C}$, there exists d and f such that $(c, f, d) \in R$
- for all $d \in \mathcal{D}$, there exists c and f such that $(c, f, s) \in R$

Proposition 4 *Two functors $F : \mathcal{C} \rightarrow \mathcal{M}$ and $G : \mathcal{D} \rightarrow \mathcal{M}$ are \mathcal{P} -bisimilar iff there exists a bisimulation between them.*

Proof

\Rightarrow Assume that there is a span:

$$\begin{array}{ccc} & H : \mathcal{E} \rightarrow \mathcal{M} & \\ (\Phi, \sigma) \swarrow & & \searrow (\Psi, \tau) \\ F : \mathcal{C} \rightarrow \mathcal{M} & & G : \mathcal{D} \rightarrow \mathcal{M} \end{array}$$

of open maps (here in the sense of the previous proposition). We define $R = \{(\Phi(e), \tau_e \circ \sigma_{\Phi(e)}^{-1}, \Psi_e) \mid e \in \mathcal{E}\}$ and show that this is a bisimulation. First, it is well defined because τ and σ are isomorphisms. The third condition of a bisimulation comes from the surjectivity of Φ . Idem for the forth and the surjectivity of Ψ . The first condition comes from the second condition on Φ of an open map: let $(\Phi(e), \tau_e \circ \sigma_{\Phi(e)}^{-1}, \Psi_e)$ in R and $i : \Phi(e) \rightarrow c' \in \mathcal{C}$. By the condition on Φ there exists $k : e \rightarrow e'$ in \mathcal{E} such that $\Phi(k) = i$. Then define $j = \Psi(k)$, $d' = \Psi(e')$ and $g = \tau_{e'} \circ \sigma_{\Phi(e')}^{-1}$. $(\Phi(e'), g, d')$ belongs to R by construction and $g \circ F(i) = G(j) \circ \tau_e \circ \sigma_{\Phi(e)}^{-1}$ by naturality of σ and τ . Idem for the second condition of a bisimulation.

\Leftarrow Assume now that there is a bisimulation R between F and G . We will construct a span of open maps. Let \mathcal{E} be the small category whose objects are elements of R , and whose morphisms from (c, f, d) to (c', f', d') are pairs (i, j) of a morphism $i : c \rightarrow c'$ in \mathcal{C} and of a morphism $j : d \rightarrow d'$ in \mathcal{D} , such that the following diagram commutes:

$$\begin{array}{ccc} F(c) & \xrightarrow{f} & G(d) \\ F(i) \downarrow & & \downarrow G(j) \\ F(c') & \xrightarrow{f'} & G(d') \end{array}$$

Define the tip H of the span between F and G as the functor $H : E \rightarrow \mathcal{M}$ that maps every object $(c, f, d) \in R$ to $F(c)$, and every morphism $(i, j) : (c, f, d) \rightarrow (c', f', d')$ to $F(i) : F(c) \rightarrow F(c')$.

We now build a morphism (Φ, σ) from H to F . We start by building $\Phi : \mathcal{E} \rightarrow \mathcal{C}$. We define Φ as the functor that maps every object (c, f, d) to c and every morphism $(i, j) : (c, f, d) \rightarrow (c', f', d')$ to $i : c \rightarrow c'$. We verify that Φ satisfies the condition of the previous proposition:

1. Φ is surjective on objects: this is third condition of the definition of R as a bisimulation.
2. Let $i : \Phi(e) \rightarrow c'$ be a morphism of X . The object e must be a triple $(c, f, d) \in R$, and i is a morphism from c to c' in \mathcal{C} . By the first condition of the definition of R as a bisimulation, there is a triple $(c', f', d') \in R$ and a morphism $j : d \rightarrow d'$ of \mathcal{D} such that the following diagram commutes:

$$\begin{array}{ccc} F(c) & \xrightarrow{f} & G(d) \\ F(i) \downarrow & & \downarrow G(j) \\ F(c') & \xrightarrow{f'} & G(d') \end{array}$$

In particular, (i, j) is a morphism of E , from (c, f, d) to (c', f', d') . Moreover, $H(i, j) = i$.

For every $(c, f, d) \in R$, let $\sigma_{(c, f, d)} = id_{F(c)} : H(c, f, d) = F(c) \rightarrow F \circ \Phi(c, f, d) = F(c)$. Those are isomorphisms, and define a natural transformation $\sigma : H \rightarrow F \circ \Phi$. It follows that (Φ, σ) is an open map from H to F .

We define the open map (Ψ, τ) from H to G similarly. □

3.3 Combinatorial and topological homology theories are bisimilar

We are now in a position to prove that, at least for cubical complexes, the homology theories we have defined combinatorially, and topologically, are equivalent. This is to be regarded as similar to the equivalence between simplicial homology and singular homology composed with geometric realization:

Theorem 1 *If K is a cubical complex, then there exists an open map*

$$(\mathcal{C}, \sigma) : \vec{H}_n(\overrightarrow{Geom}(K)) \rightarrow \vec{h}_n(K)$$

In particular $\vec{H}_n(\overrightarrow{Geom}(K))$ and $\vec{h}_n(K)$ are bisimilar.

First, let us define $\mathcal{C} : \mathcal{T}_{\overrightarrow{Geom}(K)} \rightarrow \mathcal{T}_K$. If b is a point of $\overrightarrow{Geom}(K)$, it can uniquely be written as $[e, \mathbf{a}]$ where $a \in]0, 1[^k$ for a certain k . We note $\mathcal{C}(b) = e \in K$ and call it *carrier of b* . For a dipath, we can intuitively define the sequence of all the cubes that are the carriers of points of this. Formally,

Proposition 5 ([11]) *Given a dipath π of $\overrightarrow{\text{Geom}}(K)$, there is a unique sequence c_0, c_1, \dots, c_k of elements of K and a unique sequence of real numbers $0 = t_0 \leq t_1 \leq \dots \leq t_k \leq t_{k+1} = 1$ (call them the times of change) such that:*

- for every $1 \leq i \leq k$, $c_{i-1} \neq c_i$,
- for every $0 \leq i \leq k$, for every $t \in [t_i, t_{i+1}]$, $\pi(t)$ is a point of the form $[c, \mathbf{a}]$ with $c = c_i$,
- for every $0 \leq i \leq k$, for every $t \in (t_i, t_{i+1})$, $\mathcal{C}(\pi(t)) = c_i$,
- $\mathcal{C}(\pi(0)) = c_0$ and $\mathcal{C}(\pi(1)) = c_k$,
- for every $1 \leq i \leq k$, $\mathcal{C}(\pi(t_i)) \in \{c_{i-1}, c_i\}$ and if furthermore $t_i = t_{i+1}$ then $\mathcal{C}(\pi(t_i)) = c_i$.

The sequence c_0, c_1, \dots, c_k is a discrete trace and is called the carrier sequence of π . Two dipaths that are equivalent modulo reparametrization have the same carrier sequence, so it is legitimate to call carrier sequence of a trace $\langle \pi \rangle$ the carrier sequence $\mathcal{C}(\pi)$ of π .

This induces a functor $\mathcal{C} : \mathcal{T}_{\overrightarrow{\text{Geom}}(K)} \longrightarrow \mathcal{T}_K$ which has the property of the functorial part of an open map. For the lifting property, this is a consequence of the following lemma:

Lemma 2 *Let $\langle \pi \rangle$ be a trace in $\overrightarrow{\text{Geom}}(K)$ with carrier sequence $c_0 \preceq c_1 \preceq \dots \preceq c_k$.*

- For every cube $c_{-1} \preceq c_0$, there is a dipath α in X such that $\mathcal{C}(\langle \alpha \star \pi \rangle) = c_{-1} \preceq c_0 \preceq c_1 \preceq \dots \preceq c_k$.
- For every cube c_{k+1} such that $c_k \preceq c_{k+1}$, there is a dipath β in X such that $\mathcal{C}(\langle \pi \star \beta \rangle) = c_0 \preceq c_1 \preceq \dots \preceq c_k \preceq c_{k+1}$.

Proof We examine the second case only: the other case is symmetric. Since $c_k \preceq c_{k+1}$, c_k can be a past boundary of c_{k+1} , or c_{k+1} can be a future boundary of c_k . We examine both cases:

- If c_k is a past boundary of c_{k+1} , say $c_k = \partial_{i_p}^0 \dots \partial_{i_0}^0 c_{k+1}$, then by using the precubical equations we may require $i_0 > \dots > i_p$. Writing $\pi(1)$ as $[c_k, \mathbf{a}]$, we also have $\pi(1) = [c_{k+1}, \delta_{i_0}^0 \dots \delta_{i_p}^0 \mathbf{a}]$ by the definition of the geometric realization. Since $\mathcal{C}(\pi(1)) = c_k$, no component a_i of \mathbf{a} is equal to 0 or 1. Let $\mathbf{b} = \delta_{i_0}^0 \dots \delta_{i_p}^0 \mathbf{a}$: it follows that the components b_i of \mathbf{b} that are equal to 0 are exactly those such that $i \in \{i_0, \dots, i_p\}$. Let \mathbf{a}' be the tuple whose i th component a'_i is $1/2$ if $b_i = 0$, and b_i otherwise. We define the dipath β by $\beta(t) = [c_{k+1}, (1-t)\mathbf{b} + t\mathbf{a}']$, $t \in [0, 1]$. Note that β is indeed monotonic, because $b_i \leq a'_i$ for every i . One easily checks that $\beta(0) = \pi(1)$, and that the carrier sequence of $\langle \beta \rangle$ is $c_k \preceq c_{k+1}$: for $t = 0$, $\mathcal{C}(\beta(0)) = (\pi(1)) = c_k$, and, for $t \neq 0$, $\beta(t) = [c_{k+1}, (1-t)\mathbf{b} + t\mathbf{a}']$ where no component of $(1-t)\mathbf{b} + t\mathbf{a}'$ is equal to 0 or 1, so its carrier $\mathcal{C}(\beta(t))$ is c_{k+1} . It follows that $\mathcal{C}(\langle \pi \star \beta \rangle) = c_0 \preceq c_1 \preceq \dots \preceq c_k \preceq c_{k+1}$.
- If c_{k+1} is a future boundary of c_k , then c_{k+1} is of the form $\partial_{i_p}^1 \dots \partial_{i_0}^1 c_k$ with $i_0 > \dots > i_p$, and $\pi(1) = [c_k, \mathbf{a}]$ for some tuple \mathbf{a} whose components a_i are

all different from 0 or 1 (because $\mathcal{C}(\pi(1)) = c_k$). Let \mathbf{b} be the tuple obtained from \mathbf{a} by changing the i th component into 1 if and only if $i \in \{i_0, \dots, i_p\}$. In other words, let $b_i = 1$ if $i \in \{i_0, \dots, i_p\}$, $b_i = a_i$ otherwise. One can therefore write \mathbf{b} as $\delta_{i_0}^1 \cdots \delta_{i_p}^1 \mathbf{b}'$, where \mathbf{b}' is the tuple obtained from \mathbf{b} by removing its components of indices i_0, \dots, i_p . Define the dipath β by $\beta(t) = [c_k, (1-t)\mathbf{a} + t\mathbf{b}]$. This is monotonic because $a_i \leq b_i$ for every i . For $t \neq 1$, no component of $(1-t)\mathbf{a} + t\mathbf{b}$ is equal to 0 or 1, so $\mathcal{C}(\beta(t)) = c_k$, and for $t = 1$, $\beta(1) = [c_k, \mathbf{b}] = [c_{k+1}, \mathbf{b}']$, which shows that $\mathcal{C}(\beta(1)) = c_{k+1}$ since no component of \mathbf{b}' is equal to 0 or 1. Again, it follows that $\mathcal{C}(\langle \pi \star \beta \rangle) = c_0 \preceq c_1 \preceq \cdots \preceq c_k \preceq c_{k+1}$. \square

Proof (\mathcal{C} is the functorial part of an open map)

- \mathcal{C} is surjective on objects: given a discrete trace c_0, \dots, c_k of K , $\mathcal{C}(\widehat{c_0, \dots, c_k}) = c_0, \dots, c_k$
- \mathcal{C} has the lifting property: let (α, β) be a morphism of \mathcal{T}_K i.e. a pair of a discrete trace α from x' to x and of a discrete trace β from y to y' such that there is a discrete trace from x to y . Given a pair (s, t) of points of $\overrightarrow{\text{Geom}}(K)$ such that $\mathcal{C}(s) = x$ and $\mathcal{C}(t) = y$. Then using the previous lemma and an induction, we can find traces α' from s' to s and β' from t to t' such that $\mathcal{C}(\alpha') = \alpha$ and $\mathcal{C}(\beta') = \beta$, which is what is required. \square

We now need to build a natural isomorphism

$$\sigma : \overrightarrow{H}_n(\overrightarrow{\text{Geom}}(K)) \longrightarrow \overrightarrow{h}_n(K) \circ \mathcal{C}$$

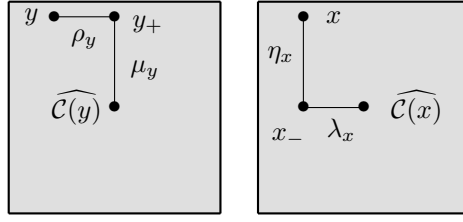
In other words, we need to build for every a, b in $\overrightarrow{\text{Geom}}(K)$ such that there is a dipath from a to b a group isomorphism

$$\sigma_{a,b} : \overrightarrow{H}_n(\overrightarrow{\text{Geom}}(K))(a, b) \longrightarrow \overrightarrow{h}_n(K)(\mathcal{C}(a), \mathcal{C}(b))$$

that is natural, in the sense that, for every extension $(\langle \alpha \rangle, \langle \beta \rangle)$ from (a, b) to (a', b') , the following square commutes:

$$\begin{array}{ccc} \overrightarrow{H}_n(\overrightarrow{\text{Geom}}(K))(a, b) & \xrightarrow{\sigma_{a,b}} & \overrightarrow{h}_n(K)(\mathcal{C}(a), \mathcal{C}(b)) \\ \downarrow \overrightarrow{H}_n(\overrightarrow{\text{Geom}}(K))(\langle \alpha \rangle, \langle \beta \rangle) & & \downarrow \overrightarrow{h}_n(K)(\mathcal{C}(\langle \alpha \rangle), \mathcal{C}(\langle \beta \rangle)) \\ \overrightarrow{H}_n(\overrightarrow{\text{Geom}}(K))(a', b') & \xrightarrow{\sigma_{a',b'}} & \overrightarrow{h}_n(K)(\mathcal{C}(a'), \mathcal{C}(b')) \end{array}$$

Let x and y such that there is a dipath from x to y . Every cube I^k has a lattice structure whose meet \wedge is pointwise min and whose join \vee is pointwise max. Write x as $[\mathcal{C}(x), \mathbf{a}]$, and let $x_- = [\mathcal{C}(x), \mathbf{a} \wedge \bullet]$. Recall that $\bullet = (\frac{1}{2}, \dots, \frac{1}{2})$, and that $\widehat{\mathcal{C}(x)} = [\mathcal{C}(x), \bullet]$. Similarly, let $\widehat{\mathcal{C}(y)} = [\mathcal{C}(y), \bullet]$, and we define $y_+ = [\mathcal{C}(y), \mathbf{b} \vee \bullet]$, where $y = [\mathcal{C}(y), \mathbf{b}]$. The situation is illustrated here:



There are obvious dipaths η_x , λ_x , μ_y , ρ_y as displayed there, too. Let us make formal the construction of the dipath η_x . The other three are similar. This is a dipath from $x_- = [\mathcal{C}(x), \mathbf{a} \wedge \bullet]$ to $x = [\mathcal{C}(x), \mathbf{a}]$, and so we just let $\eta_x(t) = [\mathcal{C}(x), (1-t)(\mathbf{a} \wedge \bullet) + t\mathbf{a}]$. Those dipaths induce continuous maps between trace spaces by concatenation. For example, there is a continuous map $\eta_x^* : \overrightarrow{\mathfrak{X}}(\overrightarrow{Geom}(K))(x, y) \longrightarrow \overrightarrow{\mathfrak{X}}(\overrightarrow{Geom}(K))(x_-, y)$ that sends each trace $\langle \pi \rangle$ to $\langle \eta_x \star \pi \rangle$. Similarly, for the other three.

Lemma 3 *The map η_x^* is a homotopy equivalence.*

Proof By abuse of language, write $\eta_x^*(\pi)$ for the dipath $\eta_x \star \pi$ as well—we reason on spaces of dipaths first, then take a reparametrization quotient. Accordingly, let $P(X; x, y)$ denote the space of dipaths from x to y in $X = \overrightarrow{Geom}(K)$, with the usual compact-open topology. (The space $\overrightarrow{\mathfrak{X}}(X)(x, y)$ is a quotient of this space.)

Observe that η_x^* maps $P(X; x, y)$ to $P(X; x_-, y)$. We need to build a map $\nu : P(X; x_-, y) \longrightarrow P(X; x, y)$ such that $\eta_x^* \circ \nu$ and $\nu \circ \eta_x^*$ are homotopic to the identity.

For every dipath π from x to y , the carrier sequence c_0, c_1, \dots, c_k of $\eta_x^*(\pi)$ is equal to that of π . In the other direction, we shall define ν so that it also preserves the carrier sequence. This will turn out to be the crucial property that will allow us to conclude.

For every dipath π from x_- to y , with carrier sequence c_0, c_1, \dots, c_k , and with times of change $0 = t_0 \leq t_1 \leq \dots \leq t_k \leq t_{k+1} = 1$, we define $\nu(\pi)$ as follows. We abuse the notation \vee , and write $[c, \mathbf{a}] \vee_c [c, \mathbf{b}]$ for $[c, \mathbf{a} \vee \mathbf{b}]$. The three occurrences of c must be the same for this notation to make sense, but our intuition is best served by ignoring the c subscript to \vee , and to understand this as taking maxes, componentwise, in a local cube c . We then define $\nu(\pi)(u)$ for increasing values of u , inductively, as $x \vee_{c_0} \pi(u)$ for $u \in [t_0, t_1]$, as $\nu(\pi)(t_1) \vee_{c_1} \pi(u)$ for $u \in [t_1, t_2]$, \dots , and finally as $\nu(\pi)(t_k) \vee_{c_k} \pi(u)$ for $u \in [t_k, t_{k+1}]$.

On $[t_0, t_1]$, $\nu(\pi)$ is a continuous monotonic map, with value $\nu(\pi)(0) = x \vee_{c_0} x_- = x$ at $u = t_0 = 0$, and with value $\nu(\pi)(t_1) = x \vee_{c_0} \pi(t_1)$ at $u = t_1$.

Let us show by induction on j that for every u with $0 \leq u \leq t_j$, $\mathcal{C}(\nu(\pi)(u)) = \mathcal{C}(\pi(u))$. For $j = 0$, this says that $\mathcal{C}(x) = \mathcal{C}(x_-)$, which is by construction of x_- . Otherwise, by induction hypothesis, for every u with $0 \leq u \leq t_j$, $\mathcal{C}(\nu(\pi)(u)) = \mathcal{C}(\pi(u))$. Let $t_j < u \leq t_{j+1}$. We can write $\pi(t_j)$ as $[c_j, (b_1, \dots, b_m)]$ and $\pi(u)$ as $[c_j, (a_1, \dots, a_m)]$, where $b_i \leq a_i$ for every i .

- If $u < t_{j+1}$, by the properties of the carrier sequence, $\mathcal{C}(\pi(u)) = c_j$, so with $0 < a_i < 1$ for every i . Since $b_i \leq a_i$, $b_i < 1$ for every i . Let us

write $\nu(\pi)(t_j)$ as $[c_j, (b'_1, \dots, b'_m)]$. Since $\mathcal{C}(\nu(\pi)(t_j)) = \mathcal{C}(\pi(t_j))$, $b_i = 1$ iff $b'_i = 1$. It follows that $b'_i < 1$ for every i . Therefore $0 < \max(a_i, b'_i) < 1$, so $\mathcal{C}(\nu(\pi)(u)) = c_k$.

- If $u = t_{j+1}$, we observe that $\max(a_i, b'_i)$ is equal to 1, resp. to 0, resp. in $(0, 1)$, if and only if a_i is. This observation is enough to conclude that $\mathcal{C}(\nu(\pi)(t_{j+1})) = \mathcal{C}(\pi(t_{j+1}))$, and is proved as follows. If $a_i = 1$, then $\max(a_i, b'_i) = 1$. If $a_i = 0$ then $b_i = 0$; moreover, since $\mathcal{C}(\nu(\pi)(t_j)) = \mathcal{C}(\pi(t_j))$, $b_i = 0$ iff $b'_i = 0$, so $b'_i = 0$, from which we obtain $\max(a_i, b'_i) = 0$. Finally, if $0 < a_i < 1$ then $b_i < 1$, and $b'_i < 1$ (since $\mathcal{C}(\nu(\pi)(t_j)) = \mathcal{C}(\pi(t_j))$), $b_i = 1$ iff $b'_i = 1$, so $0 < \max(a_i, b'_i) < 1$.

This finishes our argument that c_0, \dots, c_k is the carrier sequence of $\nu(\pi)$, with times of change $0 = t_0 \leq \dots \leq t_{k+1} = 1$.

It remains to show that $\nu(\pi)(1) = y$. This is the only place where we need the ϵ mapping. The above argument works in general precubical sets, not just cubical complexes. Here, however, we need the specific features of cubical complexes to show that $\nu(\pi)(1) = y$.

We know that $\mathcal{C}(y) = \mathcal{C}(\nu(\pi)(1)) = c_k$. Moreover, y is below $\nu(\pi)(1)$ in the ordering \leq of the pospace $\overrightarrow{\text{Geom}}(K)$, because $\nu(\pi)(1) = \nu(\pi)(t_k) \vee_{c_k} \pi(1) = \nu(\pi)(t_k) \vee_{c_k} y$. Suppose that $\nu(\pi)(1) \not\leq y$. Because K is a cubical complex, we can make use of the ϵ isomorphism. From $\nu(\pi)(1) \not\leq y$, we obtain $\epsilon(\nu(\pi)(1)) \not\leq \epsilon(y)$. Let us write $\epsilon(\nu(\pi)(t_j))$ as (x_1^j, \dots, x_d^j) and $\epsilon(\pi(t_j))$ as (y_1^j, \dots, y_d^j) . We show that $\epsilon(\nu(\pi)(t_j)) \not\leq \epsilon(y)$ by decreasing induction on j . The case $j = k+1$ is by assumption. Suppose $\epsilon(\nu(\pi)(t_{j+1})) \not\leq \epsilon(t)$. There must be an index $m \in \{1, 2, \dots, d\}$ such that $x_m^{j+1} > y_m^{k+1}$. It is easy to see that the identity $\epsilon([c, \mathbf{a}] \vee_c [c, \mathbf{b}]) = \epsilon([c, \mathbf{a}]) \vee \epsilon([c, \mathbf{b}])$ holds, where the right-hand \vee is componentwise max in \mathbb{R}^d (a property that is not usually implied by the mere fact that ϵ is an isomorphism). From that and $\nu(\pi)(t_{j+1}) = \nu(\pi)(t_j) \vee_{c_j} \pi(t_{j+1})$, we infer that $x_m^{j+1} = \max(x_m^j, y_m^{j+1})$, hence $y_m^{j+1} \leq x_m^{j+1}$. But π restricts to a dipath from t_j to y , so $\epsilon(\pi(t_j)) \leq \epsilon(t)$, and therefore $y_m^{j+1} \leq y_m^{k+1} < x_m^{j+1}$. From $y_m^{j+1} < x_m^{j+1}$ and $x_m^{j+1} = \max(x_m^j, y_m^{j+1})$, we obtain $x_m^{j+1} = x_m^j$, whence $x_m^j > y_m^{k+1}$. In particular, $\epsilon(\nu(\pi)(t_j)) \not\leq \epsilon(y)$.

Taking $j = 0$, this implies that $\epsilon(x) \not\leq \epsilon(y)$. This is impossible, since there is a dipath from x to y .

We have constructed a map ν such that π and $\nu(\pi)$ have same carrier sequence. We can now conclude by the following lemma:

Lemma 4 *Let $F, G : P(X; s, t) \longrightarrow P(X; s', t')$ such that:*

- *for every pair of dipaths p, q that are equivalent modulo reparametrization, $F(p)$ and $F(q)$ are equivalent modulo reparametrization—so F induces $\tilde{F} : \overrightarrow{\mathfrak{Z}}(X)(s, t) \longrightarrow \overrightarrow{\mathfrak{Z}}(X)(s', t')$, and similarly for G .*
- *for every π , $F(\pi)$ and $G(\pi)$ have the same carrier sequence.*

Then \tilde{F} and \tilde{G} are homotopic.

Proof Let $C(X; s', t')$ be the subspace of $P(X; s', t') \times P(X; s', t')$ that consists of pairs of dipaths that have the same carrier sequence. The key ingredient

consists in constructing a continuous map $\Gamma : I \times C(X; s', t') \longrightarrow P(X; s', t')$ in such a way that $\Gamma(0, (p, q)) = p$ and $\Gamma(1, (p, q)) = q$. Let c_0, c_1, \dots, c_k be the common carrier sequence to p and q , let $t_0 \leq t_1 \leq \dots \leq t_{k+1}$ be the times of change for p , and $s_0 \leq s_1 \leq \dots \leq s_{k+1}$ be the times of change for q . Define $u_i(t) = ts_i + (1-t)t_i$ for $t \in [0, 1]$, $0 \leq i \leq k+1$. For every $u \in [u_i(t), u_{i+1}(t)]$, define v as $\frac{u-u_i(t)}{u_{i+1}(t)-u_i(t)}$. (This is defined provided $u_i(t) \neq u_{i+1}(t)$; if this is not the case, let $v = 0$.) Then $p(v(t_{i+1} - t_i) + t_i)$ is of the form $[c_i, (a_1^u, \dots, a_m^u)]$ and $q(v(s_{i+1} - s_i) + s_i)$ is of the form $[c_i, (b_1^u, \dots, b_m^u)]$. We then define $\Gamma(t, p, q)(u) = [c_i, (1-t)a_j^u + tb_j^u]$.

We have to define a homotopy $H : I \times \overrightarrow{\mathfrak{X}}(X)(s, t) \longrightarrow \overrightarrow{\mathfrak{X}}(X)(s', t')$. It will be defined as the composition of:

- $id \times \kappa : I \times \overrightarrow{\mathfrak{X}}(X)(s, t) \longrightarrow I \times P(X; s, t)$, where κ is a continuous map from $\overrightarrow{\mathfrak{X}}(X)(s, t)$ to $P(X; s, t)$, defined in such a way that $\langle \kappa(\langle \pi \rangle) \rangle = \langle \pi \rangle$ for every trace $\langle \pi \rangle$, therefore defining a canonical dipath representing a given trace. The existence of such a map is shown by Raussen in [43], as the composition $norm \circ \overrightarrow{s}$ of two more elementary maps.
- $id \times (F, G) : I \times P(X; s, t) \longrightarrow I \times C(X; s', t')$, where (F, G) maps π to $(F(\pi), G(\pi))$.
- $\Gamma : I \times C(X; s', t') \longrightarrow P(X; s', t')$, as defined above.
- and $\langle - \rangle : P(X; s', t') \longrightarrow \overrightarrow{\mathfrak{X}}(X)(s', t')$, which maps each dipath to its trace.

We compute: $H(0, \langle \pi \rangle) = \langle \Gamma(0, (F(\kappa(\langle \pi \rangle)), G(\kappa(\langle \pi \rangle))) \rangle = \langle F(\kappa(\langle \pi \rangle)) \rangle = \tilde{F}(\langle \pi \rangle)$. Similarly, $H(1, -) = \tilde{G}$ and therefore H is an homotopy from \tilde{F} to \tilde{G} . \square

It only remains to prove that the construction is natural. The following diagram:

$$\begin{array}{ccc} \overrightarrow{\mathfrak{X}}(X)(x, y) & \xrightarrow{(*\mu_y)^{-1} \circ *\rho_y \circ (\lambda_x^*)^{-1} \circ \eta_x^*} & \overrightarrow{\mathfrak{X}}(X)(\mathcal{C}(x), \mathcal{C}(y)) \\ \langle \alpha \star - \star \beta \rangle \Big\downarrow & & \Big\downarrow \langle \hat{\gamma} \star - \star \hat{\delta} \rangle \\ \overrightarrow{\mathfrak{X}}(X)(x', y') & \xrightarrow{(*\mu_{y'})^{-1} \circ *\rho_{y'} \circ (\lambda_{x'}^*)^{-1} \circ \rho_{x'}^*} & \overrightarrow{\mathfrak{X}}(X)(\mathcal{C}(x'), \mathcal{C}(y')) \end{array}$$

where $\gamma = \mathcal{C}(\alpha)$ and $\delta = \mathcal{C}(\beta)$ is commutative modulo homotopy because of the previous lemma and so the same diagram is commutative in homology, which proves the naturality. \square

Corollary 1 *Let K be a cubical complex, and K' be a subdivision of K . Then $\overrightarrow{h}_n(K)$, and $\overrightarrow{h}_n(K')$ are bisimilar.*

Proof K and K' have isomorphic (in the category of pospaces) geometric realization, then by functoriality of natural homology, they have isomorphic (then bisimilar) natural homology then their discrete natural homology are bisimilar. \square

3.4 Hochschild-Mitchell homology theories and natural systems

Instead of taking as a base space the category of factorization of a trace category, i.e. the category \mathcal{T}_X of Definition 5, we can take the “smaller” category of pairs of base points (initial, final, see Definition 18), together with extensions (in the past, in the future). We will see that the homological information we can construct on this base space is akin to a Hochschild-Mitchell homology theory [38,39] and that the comparison functor from Baues-Wirshing homology (our homological diagrams, i.e. natural systems of Definition 8) actually induces an equivalence under bisimilarity.

Definition 18 We define \preceq the preorder on X defined as $x \preceq y$ iff there exists a dipath from x to y . We define the category \mathcal{P}_X whose:

- objects are pairs of points (x, y) of X such that $x \preceq y$
- morphisms (called extensions) from (x, y) to (x', y') are pairs $(\langle \alpha \rangle, \langle \beta \rangle)$ of trace of X with α dipath from x' to x and β dipath from y to y'

We have a functor $\kappa_X : \mathcal{T}_X \rightarrow \mathcal{P}_X$ which maps each trace to its end points.

Definition 19 We define the trace diagram as the functor $\overrightarrow{T}(X) : \mathcal{P}_X \rightarrow \mathbf{Top}$ which maps:

- every pair (x, y) with $x \preceq y$ to $\overrightarrow{\mathfrak{Z}}(X)(x, y)$
- every extension $(\langle \alpha \rangle, \langle \beta \rangle)$ from (x, y) to (x', y') to the continuous function $\langle \alpha \star _ \star \beta \rangle : \overrightarrow{\mathfrak{Z}}(X)(x, y) \rightarrow \overrightarrow{\mathfrak{Z}}(X)(x', y') \quad \langle p \rangle \mapsto \langle \alpha \star p \star \beta \rangle$

As in definition 8, we can define a diagram of abelian groups:

Definition 20 We define for $n \geq 1$, $\overrightarrow{H}_n^{HM}(X) : \mathcal{P}_X \rightarrow \mathbf{Ab}$ composing $\overrightarrow{T}(X)$ with the functor $(n-1)^{th}$ homology group functor. In particular, $\overrightarrow{H}_n(X) = \overrightarrow{H}_n^{HM}(X) \circ \kappa_X$.

As in pospaces, we can define the category \mathcal{P}_K whose objects are pairs (x, y) of elements of K such that there exists a discrete trace from x to y and morphisms from (x, y) to (x', y') are pairs (α, β) of a discrete trace α from x' to x and another, β , from y to y' . We have a functor $\kappa_K : \mathcal{T}_K \rightarrow \mathcal{P}_K$ which maps each discrete trace to its end points. When K is a cubical complex, previous remarks define a functor $\widehat{H}^{HM} : \mathcal{P}_K \rightarrow \mathcal{P}_{\overrightarrow{Geom}(K)}$. As previously, we can then define a notion of natural homology of a cubical complex:

Definition 21 Let K be a cubical complex. For $n \geq 1$, we define the n -th natural homology of K or discrete natural homology $\overrightarrow{h}_n^{HM}(K) : \mathcal{T}_K^d \rightarrow \mathbf{Ab}$ as

$$\overrightarrow{h}_n^{HM}(K) = \overrightarrow{H}_n^{HM}(\overrightarrow{Geom}(K)) \circ \widehat{H}^{HM}$$

In particular, $\overrightarrow{h}_n(X) = \overrightarrow{h}_n^{HM}(X) \circ \kappa_X$.

All those definitions coincide up-to bisimulation:

Proposition 6

- for every pospace X , $(\kappa_X, \mathbf{1}) : \overrightarrow{H}_n^{HM}(X) \longrightarrow \overrightarrow{H}_n(X)$ is an open map
- for every cubical complex K , $(\kappa_K, \mathbf{1}) : \overrightarrow{H}_n^{HM}(K) \longrightarrow \overrightarrow{H}_n(K)$ is an open map
- for every cubical complex K , like in Theorem 1, there is an open map $(\mathcal{C}, \sigma) : \overrightarrow{H}_n^{HM}(\overrightarrow{Geom}(K)) \longrightarrow \overrightarrow{h}_n^{HM}(K)$

4 Diagrammatic properties of natural homology

We are now in a position to study the diagrammatic properties of our natural homology. By diagrammatic properties, we mean Eilenberg-Steenrod [5] axioms for homology theories, and in particular, exactness properties. We first begin by classical Hurewicz theorems in Section 4.1 and then proceed by looking at the dimension, homotopy and additivity axioms. Finally, we define and study exactness of diagrams in natural systems, in Sections 4.2.2, 4.2.3 and 4.2.4. Many properties of classical homology hold true, but not quite the exactness axiom of Eilenberg-Steenrod. Looking for some form of an excision axiom is also very intricate for natural homology, and is left for future work.

4.1 Hurewicz theorems

Definition 22

We say that X is 1-connected if for every trace $\langle p \rangle$, $\overrightarrow{\Pi}_1(X)(\langle p \rangle)$ is a singleton. We say that X is n -connected with $n \geq 2$ if it is $n - 1$ -connected and if for every trace $\langle p \rangle$, $\overrightarrow{\Pi}_n(X)(\langle p \rangle)$ is a trivial group.

Now, we have for free a Hurewicz theorem between natural homotopy and natural homology:

Theorem 2

$Free \circ \overrightarrow{\Pi}_1(X)$ is isomorphic in $\mathbf{Diag}(\mathbf{Ab})$ to $\overrightarrow{H}_1(X)$ where $Free : \mathbf{Set} \longrightarrow \mathbf{Ab}$ is the functor which gives the free group.

If X is 1-connected, $Ab \circ \overrightarrow{\Pi}_2(X)$ is isomorphic in $\mathbf{Diag}(\mathbf{Ab})$ to $\overrightarrow{H}_2(X)$ where $Ab : \mathbf{Grp} \longrightarrow \mathbf{Ab}$ is the functor which gives the abelianization.

If X is $n - 1$ connected with $n \geq 3$, $\overrightarrow{\Pi}_n(X)$ is isomorphic in $\mathbf{Diag}(\mathbf{Ab})$ to $\overrightarrow{H}_n(X)$.

Proof

This is just a consequence of the naturality in the classical Hurewicz theorem. \square

We can refine this result to our Hochschild-Mitchell [38, 39] form of homology below:

Theorem 3 $Free \circ \overrightarrow{\Pi}_1(X)$ is bisimilar to $\overrightarrow{H}_1^{HM}(X)$ where $Free : \mathbf{Set} \rightarrow \mathbf{Ab}$ is the functor which gives the free group.

If X is 1-connected, $Ab \circ \overrightarrow{\Pi}_2(X)$ is bisimilar to $\overrightarrow{H}_2^{HM}(X)$ where $Ab : \mathbf{Grp} \rightarrow \mathbf{Ab}$ is the functor which gives the abelianization.

If X is $n - 1$ connected with $n \geq 3$, $\overrightarrow{\Pi}_n(X)$ is bisimilar to $\overrightarrow{H}_n^{HM}(X)$.

Proof Consequence of theorem 2 and proposition 6.

4.2 The homological category of diagrams

Another important ingredient of classical homology theories is that they all satisfy the Eilenberg-Steenrod axioms [5].

We first examine the simpler axioms: the dimension, homotopy and additivity axioms. The dimension axiom, in abelian homology theories, states that the homology H_n of a point is the object zero for all $n > 0$. The homotopy axiom states that a weak homotopy equivalence induces an isomorphism in homology. Finally, the additivity axiom states that the homology of a coproduct is the coproduct of the homologies. We will see in Section 4.2.1 that all three are trivial in our setting, if of course we interpret everything up to bisimulation.

The exactness axiom states that the so-called relative homology long sequence is exact. We will first ask ourselves whether natural homology functors transform short exact sequences (as e.g. the ones defining quotients and relative pairs of spaces) into long exact sequences, as standard homology functors do, as this is particularly useful for making actual homological calculations.

Exactness of diagrams as well as exact functors are generally defined in abelian categories [34]. Unfortunately, $\mathbf{Diag}(\mathbf{Ab})$ is not abelian: for example, it does not have a zero object.

In order to be able to consider some sort of exactness property, we will consider the theory of homology in a non-abelian setting of [16, 17]. In Section 4.2.2 we recap the notions of semi-exact and exact categories from Grandis' work, and show that natural homology lives in a semi-exact category ($\mathbf{Diag}(\mathbf{Ab})$). This allows for defining exact sequences and exact functors, as well as do a certain amount of diagram chasing. Then in Section 4.2.3, we recap the notion of a homological category, in which we can do more, we can define quotients, hence homology functors, abstractly, and still carry on diagram chasing. We show that $\mathbf{Diag}(\mathbf{Ab})$ is homological. In Section 4.2.4 we go one step further and introduce modular categories. Homological categories only map short exact sequences in a sequence of order 2, whereas modular categories map them onto long exact sequences. We show that unfortunately, $\mathbf{Diag}(\mathbf{Ab})$ is not modular. Still, we can prove that the relative homology long sequence is exact, in Section 4.2.5.

4.2.1 Dimension, homotopy and additivity axioms

We first have, trivially:

Proposition 7 (Dimension axiom) *The n th natural homology (for $n \neq 0$) of a point is trivial, i.e. it is the functor from the category with one object and one morphism which maps this object to the trivial group.*

Then the homotopy axiom is easy to prove, with natural homology:

Proposition 8 (Homotopy axiom) *Let $f : X \rightarrow Y$ be a monotonic continuous function such that for all $n \leq 1$, $\vec{\Pi}_n(f) : \vec{\Pi}_n(X) \rightarrow \vec{\Pi}_n(Y)$ is an open map then for all $n \leq 1$, $\vec{H}_n(f) : \vec{H}_n(X) \rightarrow \vec{H}_n(Y)$ is an open map.*

Proof The functorial part of $\vec{H}_n(f)$ is the same as of $\vec{\Pi}_n(f)$ and the isomorphism in homology comes from the homotopy axiom in singular homology.

This trivially implies that if f induces a bisimulation equivalence between the homotopy functors for X and Y , which is the natural weak-equivalence we have in our directed homotopy theory, then f induces a bisimulation equivalence between the homology functors for X and Y .

Finally, we consider the additivity axiom. First, we make explicit the coproduct construction in $\mathbf{Diag}(\mathcal{M})$ (although much more can be said about limits and colimits, see Appendix 5):

Lemma 5 (Coproducts in $\mathbf{Diag}(\mathcal{M})$) *The coproducts of diagrams is the disjoint union i.e. let $(F_i : \mathcal{C}_i \rightarrow \mathcal{M})_{i \in I}$ be a family of diagrams then the its coproduct is $\coprod_{i \in I} F_i : \coprod_{i \in I} \mathcal{C}_i \rightarrow \mathcal{M}$ which maps $c \in \mathcal{C}_i$ to $F_i(c)$.*

Proposition 9 (Additivity axiom) *If X is the disjoint union of X_i , $i \in I$, then for all $n \in \mathbb{N}$, $\vec{H}_n(X) = \coprod_{i \in I} \vec{H}_n(X_i)$.*

4.2.2 Semi-exact categories

We now want to be able to talk about exact sequences. Let \mathcal{A} be a category. An *ideal* of \mathcal{A} is a set of morphisms stable under left and right compositions by any (composable) morphism of \mathcal{A} .

Let N be an ideal of \mathcal{A} . We call the morphisms in N , the *null morphisms*. A *null object* is an object of \mathcal{A} whose identity is null.

We say that N is *closed* if every null morphism factorises through a null object i.e. for every $f : A \rightarrow B \in N$, there exists a null object C and two morphisms $g : A \rightarrow C$ and $h : C \rightarrow B$ such that $f = h \circ g$.

The *kernel* (with respect to N) of a morphism $f : A \rightarrow B$ of \mathcal{A} is characterized (if it exists) up to isomorphism by the following property:

- $\ker f : \text{Ker } f \rightarrow A$ such that $f \circ \ker f \in N$
- for all $g : C \rightarrow A$ such that $f \circ g \in N$, there exists a unique $h : C \rightarrow \ker f$ such that $g = \text{Ker } f \circ h$

We define dually, the *cokernel*.

A *semi-exact category* is a pair (\mathcal{A}, N) where N is a closed ideal of the category \mathcal{A} such that every morphism of \mathcal{A} has a kernel and a cokernel with respect to N .

We call *normal mono* (resp. *normal epi*), a morphism which is the kernel (resp. the cokernel) of a morphism.

We call *image* of a morphism f , $im f = ker\ cok f$ and *coimage*, $coim f = cok\ ker f$.

Definition 23 ([16]) The sequence

$$A \xrightarrow{f} B \xrightarrow{g} C$$

is said to be of *order two* (resp. *short exact*, *exact*) if $g \circ f$ is null (resp. if $f = ker\ g$ and $g = cok\ f$, if $im f = ker\ g$).

Lemma 6 *If (\mathcal{M}, N) is semi-exact then $L = \{(\Phi, \tau) \mid \forall f, \tau_f \in N\}$ is a closed ideal of $\mathbf{Diag}(\mathcal{M})$ and $\mathbf{Diag}(\mathcal{M})$ has kernels with respect to L .*

Proof

- L is an ideal because N is
- L closed: we know by 3.7 of [16] that in \mathcal{M} , $f : A \rightarrow B$ is null iff f factorizes through $ker\ 1_B : Ker\ 1_B \rightarrow B$ and that $ker\ 1_B$ is a null object. So, if $(\Phi, \sigma) : (F : \mathcal{C} \rightarrow \mathcal{M}) \rightarrow (G : \mathcal{D} \rightarrow \mathcal{M})$ is null then it factorizes through 0_G where $0_G : \mathcal{D} \rightarrow \mathcal{M}$ with $0_G(d) = ker\ 1_{G(d)}$ and $0_G(f) : 0_G(d) \rightarrow 0_G(d')$ the unique morphism which makes this square commutative:

$$\begin{array}{ccc} ker\ 1_{G(d)} & \xrightarrow{Ker\ 1_{G(d)}} & G(d) \\ \downarrow 0_G(f) & & \downarrow G(f) \\ ker\ 1_{G(d')} & \xrightarrow{Ker\ 1_{G(d')}} & G(d') \end{array}$$

coming from the universal property of $Ker\ 1_{G(d')}$.

- kernels: if $(\Phi, \sigma) : (F : \mathcal{C} \rightarrow \mathcal{M}) \rightarrow (G : \mathcal{D} \rightarrow \mathcal{M})$, we construct $Ker(\Phi, \sigma) = (\Phi_{ker}, \sigma_{ker}) : ker(\Phi, \sigma) = (F_{ker} : \mathcal{C}_{ker} \rightarrow \mathcal{M}) \rightarrow (F : \mathcal{C} \rightarrow \mathcal{M})$ as follows :
 - $\mathcal{C}_{ker} = \mathcal{C}$
 - $\Phi_{ker} = id_{\mathcal{C}}$
 - $F_{ker} : \mathcal{C} \rightarrow \mathcal{M}$ with $F_{ker}(c) = ker\ \sigma_c$ and $F_{ker}(f)$ the unique morphism which makes the left square commutative:

$$\begin{array}{ccccc}
\ker \sigma_c & \xrightarrow{Ker \sigma_c} & F(c) & \xrightarrow{\sigma_c} & G(\Phi(c)) \\
\downarrow F_{\ker(f)} & & \downarrow F(f) & & \downarrow G(\Phi(f)) \\
\ker \sigma_{c'} & \xrightarrow{Ker \sigma_{c'}} & F(c') & \xrightarrow{\sigma_{c'}} & G(\Phi(c'))
\end{array}$$

coming from the universal property of $Ker \sigma_{c'}$

- $(\sigma_{\ker})_c = Ker \sigma_c$

□

Corollary 2 *Normal monos in $(\mathbf{Diag}(\mathcal{M}), L)$ are the (Φ, σ) where Φ is an isofunctor and every σ_c is a normal mono in (\mathcal{M}, N) .*

Lemma 7 *$\mathbf{Diag}(\mathbf{Ab})$ has cokernels with respect to L .*

Proof If $(\Phi, \sigma) : (F : \mathcal{C} \rightarrow \mathbf{Ab}) \rightarrow (G : \mathcal{D} \rightarrow \mathbf{Ab})$, we construct $Cok(\Phi, \sigma) = (\Phi_{cok}, \sigma_{cok}) : (G : \mathcal{D} \rightarrow \mathbf{Ab}) \rightarrow cok(\Phi, \sigma) = (G_{cok} : \mathcal{D}_{cok} \rightarrow \mathbf{Ab})$ as follow:

- $\mathcal{D}_{cok} = \mathcal{D}$
- $\Phi_{cok} = id_{\mathcal{D}}$
- Let $\Gamma = \{(R_d)_{d \in Ob(\mathcal{D})} \mid R_d \text{ subgroup of } G(d) \text{ containing all the elements of } Im \sigma_c \text{ with } \Phi(c) = d \text{ and such that if } f : d \rightarrow d' \text{ then } G(f)(R_d) \subseteq R_{d'}\}$. Γ contains $(G(d))_{d \in Ob(\mathcal{D})}$ and is stable under intersection. Define then $(H_d)_{d \in Ob(\mathcal{D})}$ as the intersection of all the elements of Γ . Then $(H_d)_{d \in Ob(\mathcal{D})} \in \Gamma$.
- We also define $G_{cok} : \mathcal{D} \rightarrow \mathbf{Ab}$ with $G_{cok}(d) = G(d)/H_d$ and $G_{cok}(f)([x]) = [G(f)(x)]$. This is well defined because $(H_d)_{d \in Ob(\mathcal{D})} \in \Gamma$.
- $(\sigma_{cok})_d(x) = [x]$

You can observe that this is the coequalizer of (Φ, σ) and $(\Phi, (0 : F(c) \rightarrow G(\Phi(c))))_{c \in Ob(\mathcal{C})}$. □

Corollary 3 *Normal epis in $(\mathbf{Diag}(\mathbf{Ab}), L)$ are (Φ, σ) where Φ is an isofunctor and every σ_c is surjective.*

Theorem 4 *$(\mathbf{Diag}(\mathbf{Ab}), L)$ is semi-exact.*

4.2.3 Homological categories

Now, to define homology, we have to be able to talk about sub-quotient as in the case of abelian groups i.e. if $K \subseteq H \subseteq G$ are abelian groups, we can define H/K . This is not the case in general groups even if H and K are normal sub-groups of G .

Definition 24 ([16]) We say that a morphism is *exact* if it factories as $n \circ q$ with q , a normal epi and n , a normal mono.

A semi-exact category (\mathcal{A}, N) is said to be *homological* if:

- normal monos and normal epis are stable under composition
- if $m : B \rightarrow A$ is a normal mono and $q : A \rightarrow C$ is a normal epi with $m \geq \ker q$ in $\text{Sub}(A)$ (i.e. there exists k , which is unique and monic, such that $\ker q = m \circ k$) then $q \circ m$ is exact.

If $m : M \rightarrow A$ and $n : N \rightarrow A$ are two normal monos with $m \geq n$, and if $q = \text{cok } n$, the object $\text{coim } q \circ m$ (isomorph to $\text{im } q \circ m$) defined up to isomorphism will be called a *sub-quotient* of A induced by $m \geq n$ and written M/N .

Proposition 10 *Diag(Ab) is homological.*

Proof

- the normal monos (resp. epis) are stable under composition by Corollaries 2 and 3.
- let $(\Phi, \sigma) : (F : \mathcal{C} \rightarrow \mathbf{Ab}) \rightarrow (G : \mathcal{D} \rightarrow \mathbf{Ab})$ be a normal mono. Thus Φ is an isofunctor. As normal mono/epi are stable under composition by an isomorphism, we can suppose, without loss of generality, that $\Phi = \text{id}_{\mathcal{C}}$ and $\mathcal{D} = \mathcal{C}$. The same way we can take $(\Psi, \tau) : (G : \mathcal{C} \rightarrow \mathbf{Ab}) \rightarrow (H : \mathcal{E} \rightarrow \mathbf{Ab})$ a normal epi with $\Psi = \text{id}_{\mathcal{C}}$ and $\mathcal{C} = \mathcal{E}$ and so, $(\Psi, \tau) \circ (\Phi, \sigma) = (\text{id}_{\mathcal{C}}, (\tau_c \circ \sigma_c)_{c \in \text{Ob}(\mathcal{C})})$. In \mathbf{Ab} , $\tau_c \circ \sigma_c = \iota_c \circ \eta_c$ where $\eta_c = \tau_c \circ \sigma_c : F(c) \rightarrow \text{Im } \tau_c \circ \sigma_c$ which is surjective and $\iota_c : \text{Im } \tau_c \circ \sigma_c \rightarrow H(c)$ the inclusion, which is injective. Since $(\eta_c)_{c \in \text{Ob}(\mathcal{C})}$ and $(\iota_c)_{c \in \text{Ob}(\mathcal{C})}$ are natural and so $(\text{id}_{\mathcal{C}}, \eta)$ is a normal epi, $(\text{id}_{\mathcal{C}}, \iota)$ is a normal mono and $(\Psi, \tau) \circ (\Phi, \sigma) = (\text{id}_{\mathcal{C}}, \iota) \circ (\text{id}_{\mathcal{C}}, \eta)$. \square

Definition 25 ([17]) Chain complexes in a semi-exact category \mathcal{M} are defined the same way as in abelian groups requiring that the sequence of morphisms $C_{n+1} \xrightarrow{\partial_{n+1}} C_n \xrightarrow{\partial_n} C_{n-1}$ be of order two. A morphism of chain complexes $(f_n)_n : (C_n, \partial_n) \rightarrow (C'_n, \partial'_n)$ is the data of morphisms $f_n : C_n \rightarrow C'_n$ such that $\partial'_n \circ f_{n+1} = f_n \circ \partial_n$. We define this way a semi-exact category $C_{\bullet}(\mathcal{M})$ which is homological when \mathcal{M} is.

We can define for a chain complex $A_{\bullet} = (A_n, \partial_n)$:

- $Z_n(A_{\bullet}) = \ker \partial_{n-1}$
- $B_n(A_{\bullet}) = \text{im } \partial_n$
- $H_n(A_{\bullet}) = Z_n(A_{\bullet})/B_n(A_{\bullet})$

and those constructions are functorial.

Definition 26 We define $\vec{C} : \mathbf{PoTop} \rightarrow C_{\bullet}(\mathbf{Diag}(\mathbf{Ab}))$ composing \vec{T} with the functor which maps a topological space to its singular chain complex. We can define its homology as in Section 2.2.

Proposition 11 *For every $n \leq 0$, $H_n \circ \vec{C} = \vec{H}_{n+1}$.*

Proof A chain complex of diagram is the same as a diagram of chain complexes.

4.2.4 Modularity and long exact sequences in homology

In the theory of Grandis, we have the following theorem:

Theorem 5 ([17]) *Let \mathcal{M} be a homological category. For each short exact sequence in $C_\bullet(\mathcal{M})$:*

$$U \xrightarrow{m} V \xrightarrow{p} W$$

there exists a sequence of order two in \mathcal{M} :

$$\cdots \longrightarrow H_n(V) \xrightarrow{H_n(p)} H_n(W) \xrightarrow{\partial_n} H_{n-1}(U) \xrightarrow{H_{n-1}(m)} H_{n-1}(V) \longrightarrow \cdots$$

which is natural in the short exact sequence.

Moreover, [17] gives some conditions for this sequence to be exact. In particular, those conditions are always satisfied iff \mathcal{M} is modular.

A normal subobject of $F : \mathcal{C} \rightarrow \mathbf{Ab} \in \mathbf{Diag}(\mathbf{Ab})$ is a morphism of the form $(id_{\mathcal{C}}, \sigma)$ where every σ_c is an inclusion into $F(c)$. The set of all normal subobjects of F is a lattice whose order is inclusion, meet is intersection, join is union, \perp is $\sigma_c = 0$ and \top is $id_{F(c)}$. Moreover, it is a modular lattice that is if $X \leq B$ then $X \vee (A \wedge B) = (X \vee A) \wedge B$. We denote this lattice by $Nsb(F)$. If $f : F \rightarrow G$ is a morphism in $\mathbf{Diag}(\mathbf{Ab})$, we can define a Galois connection (f_*, f^*) where:

- $f_* : Nsb(F) \rightarrow Nsb(G)$ with $f_*(m) = im(f \circ m) = kercok(f \circ m)$
- $f^* : Nsb(G) \rightarrow Nsb(F)$ with $f^*(n) = ker((cokn) \circ f)$

The condition of modularity can be expressed as every morphism $f : F \rightarrow G$ satisfies:

- 1) for every $x \in Nsb(F)$, $f^* \circ f_*(x) = x \vee f^*(\perp)$
- 2) for every $y \in Nsb(G)$, $f_* \circ f^*(y) = y \wedge f_*(\top)$

Theorem 6 *In $\mathbf{Diag}(\mathbf{Ab})$, 1) and 2) fail and so $\mathbf{Diag}(\mathbf{Ab})$ is not modular.*

Proof

1) Let:

- F the functor from the discrete category $\mathbf{2}$ with two objects to \mathbf{Ab} which sends each object to \mathbb{Z}
- G the functor from the discrete category $\mathbf{1}$ with one object to \mathbf{Ab} which sends this object to \mathbb{Z}
- $f : F \rightarrow G$ the morphism (Φ, σ) where Φ sends each object of $\mathbf{2}$ to the unique object of $\mathbf{1}$ and σ is the natural transformation from F to $G \circ \Phi$ which is only composed of identities
- x the normal subobject of F which is the functor from $\mathbf{2}$ to \mathbf{Ab} which sends each object to 0 (identifying the object with the inclusion in F , which is formally the normal subobject)

In this case,

$$x \vee f^*(\perp) = x \neq F = f^* \circ f_*(x)$$

2) Let:

- F the functor from $\mathbf{1}$ to \mathbf{Ab} which this object to \mathbb{Z}
- G the functor from the category $\mathbf{2}'$ with two objects $\{a, b\}$ and one non-identity morphism which goes from a to b to \mathbf{Ab} which sends each object to \mathbb{Z} and the morphism to identity
- $f : F \rightarrow G$ the morphism (Φ, σ) where Φ sends the unique object of $\mathbf{1}$ to a and σ is the natural transformation from F to $G \circ \Phi$ which is only composed of identities
- y the normal subobject of G which is the functor from $\mathbf{2}'$ to \mathbf{Ab} which sends a to 0, b to \mathbb{Z} and the morphism to 0

In this case,

$$y \wedge f_*(\top) = y \neq 0_G = f_* \circ f^*(y)$$

where $0_G : \mathbf{2}' \rightarrow \mathbf{Ab}$ which maps each object to 0 and the morphism to 0.

□

4.2.5 Relative homology and the exactness axiom

We denote by $PoTop_2$ the category of pairs (X, A) with X a pospace and $A \subseteq X$ considered as a pospace by restricting the order of X (we call them potopological pairs). A morphism of potopological pairs from (X, A) to (Y, B) is a map $f : X \rightarrow Y$ with $f(A) \subseteq B$.

Let (X, A) be a potopological pair. As the category $\mathbf{Diag}(\mathbf{Ab})$ is homological, so is $C_\bullet(\mathbf{Diag}(\mathbf{Ab}))$, we can define $\vec{C}_\bullet(X, A) = \text{cok}(inj : \vec{C}_\bullet(A) \rightarrow \vec{C}_\bullet(X))$. We could also have defined it as the complex obtained by taking $\text{cok}(inj : \vec{C}_n(A) \rightarrow \vec{C}_n(X))$ in $\mathbf{Diag}(\mathbf{Ab})$ in each dimension. We have naturally a sequence:

$$\vec{C}_\bullet(A) \xrightarrow{inj} \vec{C}_\bullet(X) \xrightarrow{Cok(inj)} \vec{C}_\bullet(X, A)$$

which is exact but not short exact. However, if we define

$$\vec{C}_\bullet^X(A) = \ker(Cok(inj))$$

we have that $\text{cok}(Ker(Cok(inj)) : \vec{C}_\bullet^X(A) \rightarrow \vec{C}_\bullet(X)) = \vec{C}_\bullet(X, A)$ because Ker and Cok are inverse to each other on normal sub-objects and normal quotients and the sequence:

$$\vec{C}_\bullet^X(A) \xrightarrow{Ker(Cok(inj))} \vec{C}_\bullet(X) \xrightarrow{Cok(inj)} \vec{C}_\bullet(X, A)$$

is short exact in $C_\bullet(\mathbf{Diag}(\mathbf{Ab}))$ and so in $C_\bullet(Cat(\mathcal{T}_X, \mathbf{Ab}))$, where $Cat(\mathcal{T}_X, \mathbf{Ab})$ is abelian. We thus have a long exact sequence:

$$\dots \rightarrow H_n(\vec{C}_\bullet(X)) \rightarrow H_n(\vec{C}_\bullet(X, A)) \rightarrow H_{n-1}(\vec{C}_\bullet^X(A)) \rightarrow H_{n-1}(\vec{C}_\bullet(X)) \rightarrow \dots$$

in $Cat(\mathcal{T}_X, \mathbf{Ab})$ and so in $\mathbf{Diag}(\mathbf{Ab})$.

If $f : (X, A) \rightarrow (Y, B)$ is a morphism of potopological pairs, we can define $f_\bullet : \vec{C}_\bullet(X, A) \rightarrow \vec{C}_\bullet(Y, B)$ (resp. $f_\bullet : \vec{C}_\bullet^X(A) \rightarrow \vec{C}_\bullet^Y(B)$) as the unique morphism which makes commutative the right square of the first diagram (resp. the left square of the second):

$$\begin{array}{ccccc} \vec{C}_\bullet(A) & \xrightarrow{\text{inj}} & \vec{C}_\bullet(X) & \xrightarrow{\text{Cok}(\text{inj})} & \vec{C}_\bullet(X, A) \\ \downarrow f_\bullet & & \downarrow f_\bullet & & \downarrow \text{dotted } f_\bullet \\ \vec{C}_\bullet(B) & \xrightarrow{\text{inj}} & \vec{C}_\bullet(Y) & \xrightarrow{\text{Cok}(\text{inj})} & \vec{C}_\bullet(Y, B) \\ \\ \vec{C}_\bullet^X(A) & \xrightarrow{\text{Ker}(\text{Cok}(\text{inj}))} & \vec{C}_\bullet(X) & \xrightarrow{\text{Cok}(\text{inj})} & \vec{C}_\bullet(X, A) \\ \downarrow \text{dotted } f_\bullet & & \downarrow f_\bullet & & \downarrow f_\bullet \\ \vec{C}_\bullet^Y(B) & \xrightarrow{\text{Ker}(\text{Cok}(\text{inj}))} & \vec{C}_\bullet(Y) & \xrightarrow{\text{Cok}(\text{inj})} & \vec{C}_\bullet(Y, B) \end{array}$$

coming from the universal property of $\vec{C}_\bullet(X, A)$ (resp. $\vec{C}_\bullet^Y(B)$). This defines two functors from $PoTop_2$ to $\mathbf{Diag}(\mathbf{Ab})$. From this, and from Theorem 5, we deduce that the long exact sequence of relative homology is natural in (X, A) .

5 Conclusion and future work

We defined in this paper a notion of directed homology which enjoys many good properties: there are natural homology theories for a large class of pre-cubical sets and for a large class of directed spaces, which agree through geometric realization “modulo bisimulation”, and natural homology distinguishes between non-equivalent directed spaces which would not be distinguished using ordinary homology theories (see also [9]). Also, the natural homology of a subdivided pre-cubical set is bisimilar to the natural homology of the original one. We do have Hurewicz theorems linking directed homotopy with natural homology closely, and natural homology enjoys most of Eilenberg-Steenrod axioms for homology theories: dimension, homotopy, additivity and exactness

axioms in particular. Generally speaking, natural homology maps short exact sequences into some sequences. Unfortunately, they are only order 2 sequences and not long exact sequences, which make calculations slightly more complicated than with ordinary topology. There is nevertheless the hope that some variations of natural homology may agree with modularity, which implies that short exact sequences are mapped onto long exact sequences, for instance some form of natural homology *modulo bisimulation* may be a good basis for that. This is left for future work, as are some form of the excision axiom, the potentially interesting relationships with persistence homology [6] and the Baues-Wirshing homology theories defined in [27] for studying term rewriting systems, and formalism for describing computations which has some links with the concurrency models which were originally at the basis of this work.

Appendix : Limits and colimits in the category of diagrams

As is well known [34], \mathbf{Cat} is complete.

If \mathcal{C} is a small category and $F : \mathcal{C} \rightarrow \mathbf{Cat}$ a diagram in \mathbf{Cat} , the limit of F is the category \mathcal{L}_F whose:

- objects are the families $(x_c)_{c \in \text{Ob}(\mathcal{C})}$ where $x_c \in \text{Ob}(F(c))$ and for every morphism $f : c \rightarrow c'$ of \mathcal{C} , $F(f)(x_c) = x_{c'}$
- morphisms from $(x_c)_{c \in \text{Ob}(\mathcal{C})}$ to $(y_c)_{c \in \text{Ob}(\mathcal{C})}$ are the families $(g_c)_{c \in \text{Ob}(\mathcal{C})}$ where $g_c : x_c \rightarrow y_c$ and for every morphism $f : c \rightarrow c'$ of \mathcal{C} , $F(f)(g_c) = g_{c'}$
- composition $(h_c)_{c \in \text{Ob}(\mathcal{C})} \circ (g_c)_{c \in \text{Ob}(\mathcal{C})}$ is $(h_c \circ g_c)_{c \in \text{Ob}(\mathcal{C})}$
- identity of $(x_c)_{c \in \text{Ob}(\mathcal{C})}$ is $(id_{x_c})_{c \in \text{Ob}(\mathcal{C})}$

together with the projection maps.

We are now ready to define limits in $\mathbf{Diag}(\mathcal{M})$.

Definition 27 (\mathcal{C} -diagram, \mathcal{C} -limit) A \mathcal{C} -diagram of \mathcal{M} is a functor from \mathcal{C} to \mathcal{M} . We say that a category \mathcal{M} has \mathcal{C} -limits if every \mathcal{C} -diagram of \mathcal{M} has a limit.

Now, \mathcal{C} -limits of diagrams in \mathcal{M} are fully determined by \mathcal{C} -limits in \mathcal{M} :

Theorem 7 (Preservation of limits) *Let \mathcal{C} be a small category. $\mathbf{Diag}(\mathcal{M})$ has \mathcal{C} -limits iff \mathcal{M} has \mathcal{C} -limits.*

Proof First, suppose that $\mathbf{Diag}(\mathcal{M})$ has \mathcal{C} -limits. Let $F : \mathcal{C} \rightarrow \mathcal{M}$. We consider $G : \mathcal{C} \rightarrow \mathbf{Diag}(\mathcal{M})$ the following \mathcal{C} -diagram:

- for every object c of \mathcal{C} , $G(c)$ is the diagram from the category $\mathbf{1}$ with one object and one morphism to \mathcal{M} which maps this object to $F(c)$
- for every morphism $f : c \rightarrow c'$ of \mathcal{C} , $G(f)$ is the morphism

$$(id_{\mathbf{1}}, (F(f))_{* \in \text{Ob}(\mathbf{1})})$$

As $\mathbf{Diag}(\mathcal{M})$ has \mathcal{C} -limits, G has a limit $(L, \Pi_c : L \rightarrow G(c))$. Necessarily, L is from $\mathbf{1}$ to \mathcal{M} and so is just an object K of \mathcal{M} and the Π_c are just morphism $\pi_c : K \rightarrow F(c)$. It is easy to see that (K, π_c) is the limit of F .

Now, suppose that \mathcal{M} has \mathcal{C} -limits. Let $F : \mathcal{C} \rightarrow \mathbf{Diag}(\mathcal{M})$. F induces a \mathcal{C} -diagram in \mathbf{Cat} $G_F : \mathcal{C} \rightarrow \mathbf{Cat}$ in the following way:

- $G_F(c)$ is the domain of $F(c)$
- $G_F(f)$ is the first component of $F(f)$

As \mathbf{Cat} is complete, G_F has a limit $(\mathcal{L}_{G_F}, \pi_c : \mathcal{L}_{G_F} \rightarrow G(c))$. Now, the limit of F is the functor $\mathcal{L}_F : \mathcal{C} \rightarrow \mathcal{M}$ such that:

- $\mathcal{L}_F((x_c)_{c \in \text{Ob}(\mathcal{C})})$ is the limit of the functor which maps
 - * every $c \in \text{Ob}(\mathcal{C})$ to $F(c)(x_c)$
 - * every $f : c \rightarrow c'$ of \mathcal{C} to the component of x_c in the second component of $F(f)$ (which is a natural transformation).

Such a limit exists because \mathcal{M} has \mathcal{C} -limits.

- $\mathcal{L}_F((g_c)_{c \in \text{Ob}(\mathcal{C})})$ (where $g_c : x_c \rightarrow y_c$) is the unique morphism of \mathcal{M} defined below. As $\mathcal{L}_F((y_c)_{c \in \text{Ob}(\mathcal{C})})$ is a limit in \mathcal{M} , it makes every such diagram:

$$\begin{array}{ccc} \mathcal{L}_F((x_d)_d) & \overset{\exists!}{\dashrightarrow} & \mathcal{L}_F((y_d)_d) \\ \downarrow & & \downarrow \\ F(c)(x_c) & \xrightarrow{F(c)(g_c)} & F(c)(y_c) \end{array}$$

commutative, together with the projection maps

$$(\pi_c, (\sigma_{c, (x_d)_d})_{(x_d)_d \in \text{Ob}(\mathcal{L}_{G_F})}) : \mathcal{L}_F \rightarrow F(c)$$

where $\sigma_{c, (x_d)_d} : \mathcal{L}_F((x_d)_d) \rightarrow F(c)(x_c)$ is the projection map coming from the fact that $\mathcal{L}_F((x_d)_d)$ is a limit. \square

Therefore, a direct consequence is that $\mathbf{Diag}(\mathcal{M})$ is complete iff \mathcal{M} is complete.

Similarly, we know that \mathbf{Cat} is cocomplete [34].

We briefly recap the construction below. We need first to define congruences on categories.

If $f : c \rightarrow c'$, we note $\text{dom}(f) = c$ and $\text{cod}(f) = c'$. If \sim is an equivalence relation on the set $\text{Ob}(\mathcal{C})$ of a small category \mathcal{C} , we denote by $\text{Mor}_+(\mathcal{C})$, the set of non-empty finite sequences of morphisms of \mathcal{C} (f_1, \dots, f_n) .

A congruence on the small category \mathcal{C} is a pair (\sim_o, \sim_m) of an equivalence relation \sim_o on $\text{Ob}(\mathcal{C})$ and a partial equivalence relation \sim_m on $\text{Mor}_+(\mathcal{C})$ satisfying:

- if $(f_0, \dots, f_n) \sim_m (g_1, \dots, g_p)$ then for all $i \leq n$, $\text{dom}(f_i) \sim_o \text{dom}(f_{i-1})$

- $(f_0, \dots, f_n) \sim_m (g_0, \dots, g_p)$ implies $\text{dom}(f_0) \sim_o \text{dom}(g_0)$ and $\text{cod}(f_n) \sim_o \text{cod}(g_p)$
- $c \sim_0 c'$ implies $\text{id}_c \sim_m \text{id}_{c'}$
- $(f_0, \dots, f_n) \sim_m (g_0, \dots, g_p)$, $(f_{n+1}, \dots, f_{n'}) \sim_m (g_{p+1}, \dots, g_{p'})$ and $\text{cod}(f_n) \sim_o \text{dom}(f_{n+1})$ then $(f_0, \dots, f_{n'}) \sim_m (g_0, \dots, g_{p'})$
- if $\text{cod}(f) = \text{dom}(g)$ then $(f, g) \sim_m (g \circ f)$

Proposition 12 ([2]) *If (\sim_0, \sim_m) is a congruence on \mathcal{C} , the category $\mathcal{C}/(\sim_0, \sim_m)$ whose:*

- *objects are equivalence class of objects of \mathcal{C} modulo \sim_o , written $[_]$*
- *morphisms from $[c]$ to $[c']$ are equivalence class of elements (f_1, \dots, f_n) of $\text{Mor}_+(\mathcal{C})$ such that $\text{dom}(f_1) \sim_o c$ and $\text{cod}(f_n) \sim_o c'$, written $[f_1, \dots, f_n]$*
- *composition $[g_1, \dots, g_p] \circ [f_1, \dots, f_n]$ is $[f_1, \dots, f_n, g_1, \dots, g_p]$*
- *identity on $[c]$ is $[\text{id}_c]$*

is well defined and is called the quotient of \mathcal{C} by (\sim_o, \sim_m)

Proof (of Proposition 5) First, it is known [2] that if R_o is a relation on $\text{Ob}(\mathcal{C})$ and R_m is a relation on $\text{Mor}_+(\mathcal{C})$ then there exists a least congruence on \mathcal{C} that contains (R_o, R_m) .

Now, let \mathcal{C} be a small category and $F : \mathcal{C} \rightarrow \mathbf{Cat}$ a diagram in \mathbf{Cat} . The colimit Γ_F of F is the quotient of the disjoint union of all the $F(c)$ by the least congruence (\sim_o, \sim_m) such that:

- if $f : c \rightarrow c'$ is a morphism of \mathcal{C} , then for every object x_c of $F(c)$, $F(f)(x_c) \sim_o x_c$
- if $f : c \rightarrow c'$ is a morphism of \mathcal{C} , then for every morphism g_c of $F(c)$, $(F(f)(g_c)) \sim_m (g_c)$

□

As in the case of limits, a colimit of a diagram in $\mathbf{Diag}(M)$ is defined as a functor from the colimit of the diagram induced in \mathbf{Cat} . The functor maps each object of this colimit to the colimit in \mathcal{M} of a diagram which depends of this object and is much more complicated than the initial diagram. We can then derive similarly as in the case of limits:

Theorem 8 *If \mathcal{M} is cocomplete then $\mathbf{Diag}(M)$ is cocomplete.*

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