

# Time has something to tell us about Network Address Translation

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## **Abstract**

In this paper we introduce a new technique to count the number of host behind a NAT. This technique based on TCP timestamp option, work with Linux and BSD system and therefore is complementary to the previous one base on IPID than does not work for those systems. Our implementation demonstrates the practicability of this method.

## **1 Introduction**

Network Address Translation (NAT) also known as IP Masquerading involves re-writing the source and/or destination addresses of IP packets as they pass through a Router or firewall. Most systems using NAT do so in order to enable multiple hosts on a private network to access the Internet using a single public IP address. This technique as became popular for number of reasons, including the shortage of IPv4 address. It has become a standard feature in routers for home and small-office Internet connections. Until now, the only method known to count how many host are hidden behind a NAT rely on the IPID field. However, this method as main limitation: It does not work with systems that does not implement the IPID field as a simple counter. For instance Linux use the constant 0 for the IPID field and FreeBSD and OpenBSD use a pseudo random generator fir the IPID field. This is why a complementary approach is necessary to deal with these implementations. The main contribution of this paper is a technique base on the analysis of the TCP timestamp option to count the hosts behind a NAT. Experimentation shows that this technique works against Linux and OpenBSD NAT mechanism. Additionally, this technique can be use to performs active fingerprint consistency. The remainder of this paper is structured as follows. Section 2 provides an overview of related work concerning NAT detection. Section 3

presents the TCP timestamp option and its current uses in security. Section 4 details the algorithm that use the timestamp for NAT counting. Section 5 introduces several additionally uses of the techniques. Section 6 covers the limitation of the technique and how it can be evaded. Section 7 concludes the paper discussing directions for future work.

## 2 Related Work

The NAT mechanism was introduced in [4], and the TCP timestamp option in [6]. The technique that use IPID field to count NATed hosts was presented in [1]. Several improvement was made to this technique including the use of OS fingerprinting as in [3]. NAT can also been detected because it may alters some packet headers as suggested in [9]. The timestamp option is already used in security for uptime analysis [5]. The first use of IPID simple counter behavior, for port scanning was presented in [2].

## 3 The TCP Timestamp Option

The Timestamps option carries two four-byte timestamp fields. The timestamp Value field (TSval) contains the current value of the timestamp clock of the TCP sending the option. According to RFC1323: *"The timestamps are used for two distinct mechanisms: RTTM (Round Trip Time Measurement) and PAWS (Protect Against Wrapped Sequences)*. Each OS implementation increment this value at a specific rate. So it is possible to infer the system uptime from the timestamp option if the targeted OS is known [7]. For example Linux increment it every 1ms and FreeBSD every 500 ms. This technique is used in the popular scanner Nmap use to infer remotely the computer uptime :

Uptime 26.271 days (since Wed Jun 27 09:31:13 2007)

Once Nmap has determine the OS type by the mean of TCP/IP fingerprinting, it compute the uptime using the following formula:

$$\frac{timestamp}{numinbysec} = uptimeinsec$$

This is the simplified version that does not take into account the TSval warp around. The OS specific increment used in our algorithm are those found in Nmap source code.

## 4 Counting Hosts using the Timestamp Option

The timestamp value can be represented as a linear function:

$$\text{timestamp} = \text{numinbysec} * \text{sec} + \text{in}$$

where *numinbysec* is the slop and is dependent of the OS and *in* is the initial timestamp value. The initial number is zero for every OS type except windows. Windows computer use a random number as *in*. This is why the uptime cannot be computed remotely for windows computer.

To overcome this limitation and have a technique that can be used passively we do not rely on the OS fingerprinting to infer the increments or the initial timestamp value. Instead we use two values of the timestamp, to solve the following equation:

$$ts1 - ts2 = A * (s1 - s2) \Leftrightarrow A = \frac{ts1 - ts2}{s1 - s2}$$

where *ts1* and *ts2* are the two timestamp gathered at time *t1* and *t2*. This allows us to determine passively the increment.

The NAT counting algorithm is based on the idea that two hosts behind a NAT have a distinctive timestamp because they can't have the same uptime unless they have been booted within the same millisecond which is more than unlikely. Intuitively the counting algorithm work as follows: For each host a set of linear function is stored. Every packet timestamp is verified against this set of linear function. If it belongs to one of the already known functions then the host is already counted. If not then a new host is detected.

This algorithm is efficient in term of space complexity because it requires only to keeps two points for each real host. It is also efficient in time complexity because it only works with simple linear function. As additional benefit, knowing the increment for a given host allows to infer its type of OS. This allows to fingerprint computer that are hidden behind a NAT.

Experiment have shown that it works against standard router such as Netgear one but also against OpenBSD scrub NAT function because this function does not modify the timestamp option.

## 5 Additional Uses

An other application of the method is to combine it with other active fingerprinting method. This is useful to defeat anti-fingerprint method, by comparing the computed increment with the OS type detected by the fingerprint active method. Of course the method can be used by it self as a fingerprint method but it is less accurate since it can only infers the type of OS.

PAT (Port Address Translation) is used to access a service which is located on computer behind a NAT. Once again our method can help to determine if a PAT mechanism is used because the linear function associated to the PAT will not be consistent with the router one. Moreover if the PAT is used for round robin purpose then it possible to detect it and infers the number of hosts used in the round robin. There will be one linear function for each host that participate to the round robin. To the best of our knowledge this is the only technique usable to determine how many host are involved in a PAT round robin.

## 6 Limitation and Evasion

The main limitation of the technique is the fact that by default Windows system does not use the timestamp option when they establish a connection [8]. This limitation only affects the passive analysis because if the active probe use the timestamp option, then the Windows system response will also contains the timestamp option. As explained in the introduction, the host counting technique based on the IPID works with Windows system, hence the combination of the two techniques allows to overcome this limitation. To be complete on the subject, it seems that Windows 95, 98, Me TCP/IP stack does not implement the timestamp option. The other limitation is due to the network latency: it may introduce a bias in the measurement of the time interval between two packets. Fortunately the increment value used by the different class of OS are sufficiently different to allow the implementation to have a windows of acceptance large enough to overcome this limitation. Experiment confirms that on a production network, this measure is sufficient. There is mainly three methods to evade this technique:

1. Disabling the TCP timestamp option. This is not a viable option since it is needed on high speed network in particular to prevent sequence warp around problem.

2. Add a rewriting mechanism to the NAT computer in order to normalize packets timestamp.
3. Change the OS implementation so it use a random increment value for its timestamp.

## 7 Conclusion

In this paper we have introduced a new method to count hosts hidden behind NAT. This method uses the TCP timestamp option. We also have presented how this method can be used to infer the OS type of each computer hidden behind a given NAT and how to use it to improve active fingerprinting. A future work is to see how this method can be combined to other passive fingerprint methods to improve their accuracy.

## References

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